The Mockingbird system
Order-theoretic and enumerative properties

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Outline

1. Combinatory logic

2. Mockingbird lattices

3. Enumerative properties

1. Combinatory logic

Applicative terms

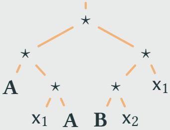
Let $\mathfrak{G}$ be a set, called alphabet.

A $\mathfrak{G}$-term is either

- a variable x_i from the set $\mathbb{X}_n := \{\mathsf{x}_1, \dots, \mathsf{x}_n\}$ for an $n \geqslant 0$;
- a basic combinator $\mathbf{X}$ where $\mathbf{X} \in \mathfrak{G}$;
- a pair $(\mathfrak{t}_1, \mathfrak{t}_2)$ where $\mathfrak{t}_1$ and $\mathfrak{t}_2$ are $\mathfrak{G}$-terms, denoted by $\mathfrak{t}_1 \star \mathfrak{t}_2$.

Let $\mathfrak{T}(\mathfrak{G}) := \bigsqcup_{n \geqslant 0} \mathfrak{T}(\mathfrak{G})(n)$ where $\mathfrak{T}(\mathfrak{G})(n)$ is the set of the $\mathfrak{G}$-terms having all variables in $\mathbb{X}_n$.

– Example –



The tree of the left is the **tree representation** of the $\mathfrak{G}$-term

$$(\mathbf{A} \star (\mathsf{x}_1 \star \mathbf{A})) \star (((\mathbf{B} \star \mathsf{x}_2) \star \mathsf{x}_1))$$

where $\mathfrak{G} := \{\mathbf{A}, \mathbf{B}, \mathbf{C}\}$.
The **short representation** is obtained by considering that $\star$ associates to the left and by removing the superfluous parentheses:

$$\mathbf{A}(\mathsf{x}_1\, \mathbf{A})(\mathbf{B}\, \mathsf{x}_2 \mathsf{x}_1).$$

Combinatory logic systems

A combinatory logic system (CLS) is a pair $(\mathfrak{G}, \to)$ where $\to$ is a binary relation on $\mathfrak{T}(\mathfrak{G})$ such that for each $\mathbf{X} \in \mathfrak{G}$, there is $n \geqslant 1$ and $\mathfrak{t}_{\mathbf{X}} \in \mathfrak{T}(\emptyset)(n)$ such that

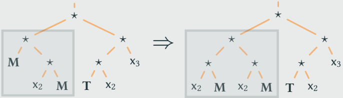
$$\mathbf{X}\, \mathsf{x}_1 \dots \mathsf{x}_n \to \mathfrak{t}_{\mathbf{X}}.$$

The context closure of $\to$ is the binary relation $\Rightarrow$ on $\mathfrak{T}(\mathfrak{G})$ such that $\mathfrak{t} \Rightarrow \mathfrak{t}'$ if $\mathfrak{t}'$ can be obtained from $\mathfrak{t}$ by replacing a pattern $\mathbf{X}\, \mathsf{x}_1 \dots \mathsf{x}_n$ by $\mathfrak{t}_{\mathbf{X}}$.

– Example –

Let the CLS $(\mathfrak{G}, \to)$ such that $\mathfrak{G} := \{\mathbf{M}, \mathbf{T}\}$ where $\mathbf{M}\mathsf{x}_1 \to \mathsf{x}_1\mathsf{x}_1$ and $\mathbf{T}\,\mathsf{x}_1\mathsf{x}_2 \to \mathsf{x}_2\mathsf{x}_1$. We have

$$\big(\underline{\mathbf{M}(\mathsf{x}_2\mathbf{M})}\big)(\mathbf{T}\,\mathsf{x}_2\mathsf{x}_3) \Rightarrow \big(\underline{(\mathsf{x}_2\mathbf{M})(\mathsf{x}_2\mathbf{M})}\big)(\mathbf{T}\,\mathsf{x}_2\mathsf{x}_3)$$



$$(\mathbf{M}(\mathsf{x}_2\mathbf{M}))\big(\underline{\mathbf{T}\,\mathsf{x}_2\mathsf{x}_3}\big) \Rightarrow (\mathbf{M}(\mathsf{x}_2\mathbf{M}))\big(\underline{\mathsf{x}_3\mathsf{x}_2}\big)$$

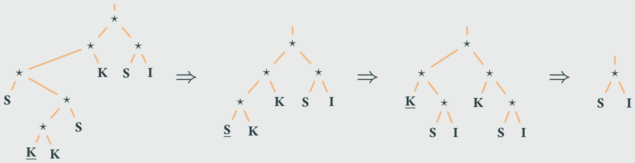
The S, K, I-system

Let the system [**Curry**, 1930] made on the three basic combinators **S**, **K**, and **I**, satisfying

$$\mathbf{S}\,x_1x_2x_3 \to x_1x_3(x_2x_3), \quad \mathbf{K}\,x_1x_2 \to x_1, \quad \mathbf{I}\,x_1 \to x_1.$$

– Example –

Here is a sequence of computation:



This CLS is Turing-complete: there are algorithms to emulate any λ-term by a term of this CLS. These algorithms are called **abstraction algorithms** [**Rosser**, 1955], [**Curry, Feys**, 1958].

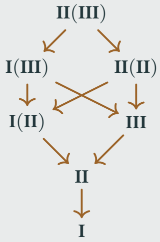
Rewrite graphs

Given a CLS $\mathcal{C} := (\mathfrak{G}, \rightarrow)$, let

- $\preccurlyeq$ be the **reflexive and transitive closure** of $\Rightarrow$;
- $\equiv$ be the **reflexive, symmetric, and transitive closure** of $\Rightarrow$;
- for any $\mathfrak{t} \in \mathfrak{T}(\mathfrak{G})$, let $\mathfrak{t}^* := \{\mathfrak{t}' \in \mathfrak{T}(\mathfrak{G}) : \mathfrak{t} \preccurlyeq \mathfrak{t}'\}$. The graph $(\mathfrak{t}^*, \Rightarrow)$ is the rewrite graph of $\mathfrak{t}$.

– Example –

Let the CLS $(\mathfrak{G}, \rightarrow)$ such that $\mathfrak{G} := \{\mathbf{I}\}$ and $\mathbf{I}\mathrm{x}_1 \rightarrow \mathrm{x}_1$.



This is the rewrite graph of $\mathbf{II}(\mathbf{III})$.

We have $\mathbf{I}(\mathbf{III}) \preccurlyeq \mathbf{I}$ and $\mathbf{I}(\mathbf{III}) \not\preccurlyeq \mathbf{II}(\mathbf{II})$.

It is possible to prove that for any $\mathfrak{t}, \mathfrak{t}' \in \mathfrak{T}(\mathfrak{G})$, $\mathfrak{t} \equiv \mathfrak{t}'$.

The Enchanted Forest of combinator birds

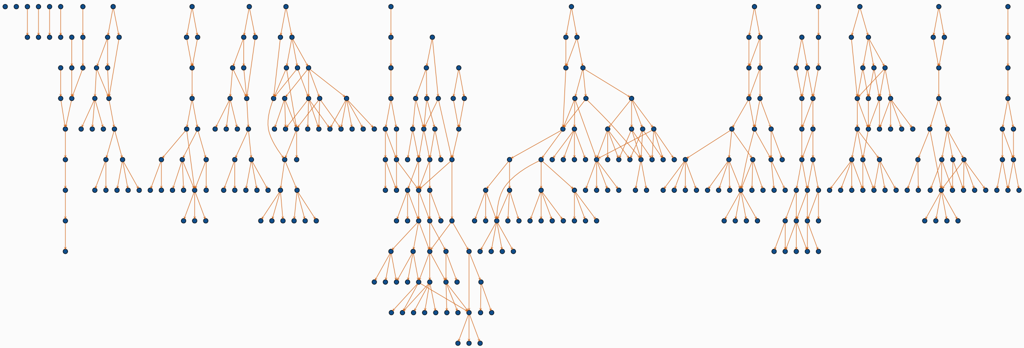
In *To Mock a Mockingbird: and Other Logic Puzzles* [Smullyan, 1985], a great number of basic combinators with their rules are listed, forming the **Enchanted forest of combinator birds**.

Here is a sublist:

- **Identity bird**: $\mathbf{I}\,x_1 \to x_1$
- **Mockingbird**: $\mathbf{M}\,x_1 \to x_1 x_1$
- **Kestrel**: $\mathbf{K}\,x_1 x_2 \to x_1$
- **Thrush**: $\mathbf{T}\,x_1 x_2 \to x_2 x_1$
- **Mockingbird 1**: $\mathbf{M_1}\,x_1 x_2 \to x_1 x_1 x_2$
- **Warbler**: $\mathbf{W}\,x_1 x_2 \to x_1 x_2 x_2$
- **Lark**: $\mathbf{L}\,x_1 x_2 \to x_1 (x_2 x_2)$
- **Owl**: $\mathbf{O}\,x_1 x_2 \to x_2 (x_1 x_2)$
- **Turing bird**: $\mathbf{U}\,x_1 x_2 \to x_2 (x_1 x_1 x_2)$
- **Cardinal**: $\mathbf{C}\,x_1 x_2 x_3 \to x_1 x_3 x_2$
- **Vireo**: $\mathbf{V}\,x_1 x_2 x_3 \to x_3 x_1 x_2$
- **Bluebird**: $\mathbf{B}\,x_1 x_2 x_3 \to x_1 (x_2 x_3)$
- **Starling**: $\mathbf{S}\,x_1 x_2 x_3 \to x_1 x_3 (x_2 x_3)$
- **Jay**: $\mathbf{J}\,x_1 x_2 x_3 x_4 \to x_1 x_2 (x_1 x_4 x_3)$

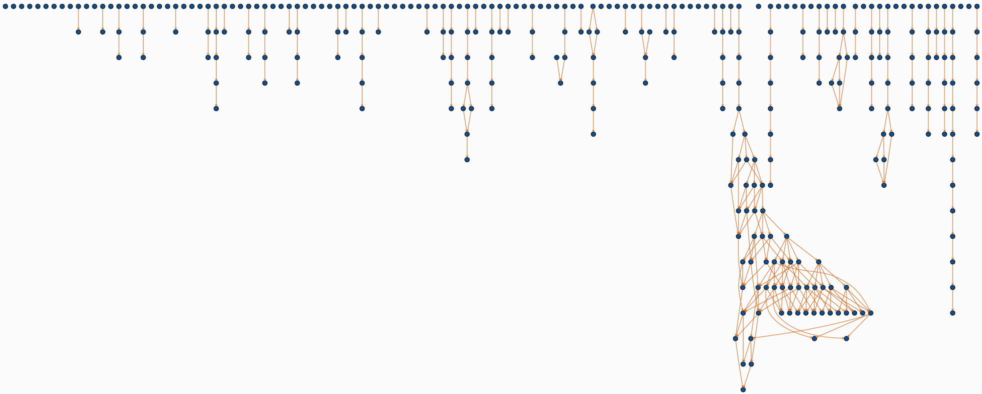
Rewrite graph of L

In the CLS containing only the **Lark L**, the rewrite graphs of closed terms of degrees up to 5 and up to 4 rewrite steps have the shape



Rewrite graph of S

In the CLS containing only the **Starling S**, the rewrite graphs of closed terms of degrees up to 6 and up to 11 rewrite steps have the shape



Usual questions

Let $\mathcal{C}$ be a CLS.

– Word problem –

Is there an algorithm to decide, given two terms t and t' of $\mathcal{C}$, if $t \equiv t'$?
See [**Baader, Nipkow**, 1998], [**Statman**, 2000].

- Yes for the CLS on **L** [**Statman**, 1989], [**Sprenger, Wymann-Böni**, 1993].
- Yes for the CLS on **W** [**Sprenger, Wymann-Böni**, 1993].
- Yes for the CLS on $\mathbf{M_1}$ [**Sprenger, Wymann-Böni**, 1993].
- Open for the CLS on **S** [**RTA Problem #97**, 1975].

– Strong normalization problem –

Is there an algorithm to decide, given a term t of $\mathcal{C}$, if all rewrite sequences from t are finite?

- Yes for the CLS on **S** [**Waldmann**, 2000].
- Yes for the CLS on **J** [**Probst, Studer**, 2000].

Order theoretical questions

A lattice is a partial order (poset) wherein each pair $\{x, x'\}$ of elements has a greatest lower bound $x \wedge x'$ and a least upper bound $x \vee x'$.

Let $\mathcal{C}$ be a CLS. A $\mathfrak{G}$-term $\mathfrak{t}$ has

1. the poset property if $(\mathfrak{t}^*, \preccurlyeq)$ is a **poset**;
2. the lattice property if $(\mathfrak{t}^*, \preccurlyeq)$, is a **lattice**.

This CLS has the poset (resp. lattice) property if all terms of $\mathcal{C}$ have the poset (resp. lattice) property.

– Poset and lattice properties –

Is there an algorithm to decide, given a term $\mathfrak{t}$ of $\mathcal{C}$, if $\mathfrak{t}$ has the poset (resp. lattice) property?

Given a term $\mathfrak{t}$ of $\mathcal{C}$, perform a **combinatorial study** of $(\mathfrak{t}^*, \preccurlyeq)$ as the enumeration of its elements and intervals.

– A new source of posets –

Use combinatory logic as a source to **build original posets**.

Outline

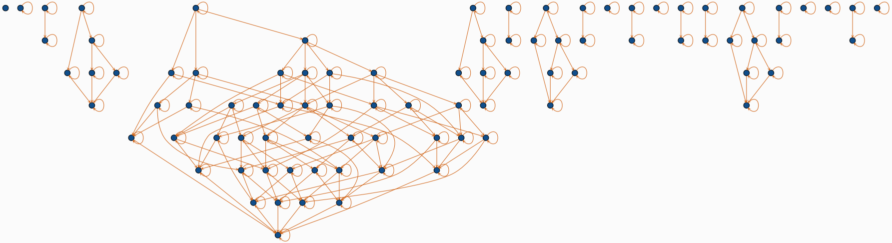
2. Mockingbird lattices

The Mockingbird system

The Mockingbird system is the CLS $\mathcal{C}$ containing only the Mockingbird **M**.

Recall that **M** satisfies $\mathbf{M}x_1 \to x_1 x_1$.

The **rewrite graphs** of closed terms of $\mathcal{C}$ of degrees up to 4 have the shape



– Proposition [G., 2022] –

The CLS $\mathcal{C}$ has the **poset property** and each $\equiv$-equivalence class of $\mathcal{C}$ is **finite** and contains a **greatest** and a **least** element.

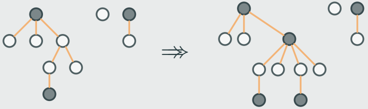
Duplicative forests

A duplicative forest is a **forest of planar rooted trees** where nodes are either black ● or white ○.

Let $\mathcal{D}^{\mathrm{F}}$ be the set of the duplicative forests and $\mathcal{D}^{\mathrm{T}}$ be the set of the duplicative trees.

Let $\Rrightarrow$ be the binary relation on $\mathcal{D}^{\mathrm{F}}$ such that for any $\mathfrak{f}, \mathfrak{f}' \in \mathcal{D}^{\mathrm{F}}$, we have $\mathfrak{f} \Rrightarrow \mathfrak{f}'$ if $\mathfrak{f}'$ is obtained by blackening a white node of $\mathfrak{f}$ and then by **duplicating** its sequence of descendants.

– Example –



The reflexive and transitive closure $\ll$ of $\Rrightarrow$ is a **partial order relation**.

For any $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$, let $\mathfrak{f}^* := \left\{\mathfrak{f}' \in \mathcal{D}^{\mathrm{F}} : \mathfrak{f} \ll \mathfrak{f}'\right\}$.

Lattice of duplicative forests

Let $\wedge$ and $\vee$ be the two binary, commutative, and associative partial operations on $\mathcal{D}^{\mathrm{F}}$ defined recursively, for any $\ell \geqslant 0$, $\mathfrak{f}_1, \dots, \mathfrak{f}_\ell \in \mathcal{D}^{\mathrm{T}}$, $\mathfrak{f}'_1, \dots, \mathfrak{f}'_\ell \in \mathcal{D}^{\mathrm{T}}$, and $\mathfrak{f}, \mathfrak{f}', \mathfrak{f}'' \in \mathcal{D}^{\mathrm{F}}$, by

$$\mathfrak{f}_1 \dots \mathfrak{f}_\ell \wedge \mathfrak{f}'_1 \dots \mathfrak{f}'_\ell := (\mathfrak{f}_1 \wedge \mathfrak{f}'_1) \dots (\mathfrak{f}_\ell \wedge \mathfrak{f}'_\ell),$$

$$\circ(\mathfrak{f}) \wedge \circ(\mathfrak{f}') := \circ(\mathfrak{f} \wedge \mathfrak{f}'), \qquad \bullet(\mathfrak{f}) \wedge \bullet(\mathfrak{f}') := \bullet(\mathfrak{f} \wedge \mathfrak{f}'),$$

$$\circ(\mathfrak{f}) \wedge \bullet(\mathfrak{f}'\mathfrak{f}'') := \circ(\mathfrak{f} \wedge \mathfrak{f}' \wedge \mathfrak{f}''),$$

$$\mathfrak{f}_1 \dots \mathfrak{f}_\ell \vee \mathfrak{f}'_1 \dots \mathfrak{f}'_\ell := (\mathfrak{f}_1 \vee \mathfrak{f}'_1) \dots (\mathfrak{f}_\ell \vee \mathfrak{f}'_\ell),$$

$$\circ(\mathfrak{f}) \vee \circ(\mathfrak{f}') := \circ(\mathfrak{f} \vee \mathfrak{f}'), \qquad \bullet(\mathfrak{f}) \vee \bullet(\mathfrak{f}') := \bullet(\mathfrak{f} \vee \mathfrak{f}'),$$

$$\circ(\mathfrak{f}) \vee \bullet(\mathfrak{f}'\mathfrak{f}'') := \bullet((\mathfrak{f} \vee \mathfrak{f}')(\mathfrak{f} \vee \mathfrak{f}'')).$$

– Proposition [G., 2022] **–**

Given a duplicative forest $\mathfrak{f}$, the poset $(\mathfrak{f}^*, \ll)$ is a **lattice** for the operations $\wedge$ and $\vee$.

This can be proved by structural induction on duplicative forests.

From Mockingbird terms to duplicative trees

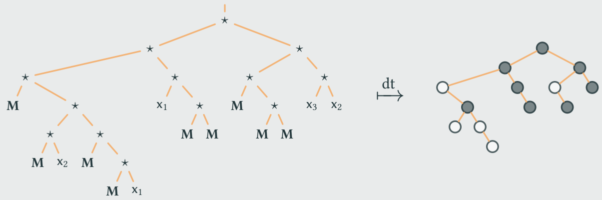
Let $\mathrm{dt} : \mathfrak{T}(\mathfrak{G}) \to \mathcal{D}^{\mathrm{T}}$ be the map defined recursively, for any $\mathrm{x}_i \in \mathbb{X}$ and $\mathfrak{t}, \mathfrak{t}' \in \mathfrak{T}(\mathfrak{G})$, by

$$\mathrm{dt}(\mathrm{x}_i) := \epsilon,$$

$$\mathrm{dt}(\mathbf{M}) := \epsilon,$$

$$\mathrm{dt}(\mathfrak{t} \star \mathfrak{t}') := \begin{cases} \circ(\mathrm{dt}(\mathfrak{t}')) & \text{if } \mathfrak{t} = \mathbf{M} \text{ and } \mathfrak{t}' \neq \mathbf{M}, \\ \bullet(\mathrm{dt}(\mathfrak{t})\ \mathrm{dt}(\mathfrak{t}')) & \text{otherwise.} \end{cases}$$

– Example –



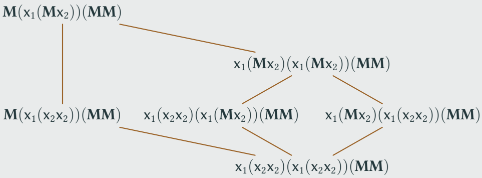
Poset isomorphism

– Proposition [G., 2022] –

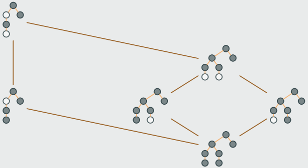
For any $\mathfrak{t} \in \mathfrak{T}(\mathfrak{G})$, the **posets** $(\mathfrak{t}^*, \preccurlyeq)$ and $(\mathrm{dt}(\mathfrak{t})^*, \ll)$ are **isomorphic** and dt is such an isomorphism.

– Example –

Let $\mathfrak{t} := \mathbf{M}(\mathrm{x}_1(\mathbf{M}\mathrm{x}_2))(\mathbf{MM})$.



The Hasse diagram of the poset $(\mathfrak{t}^*, \preccurlyeq)$.



The Hasse diagram of the poset $(\mathrm{dt}(\mathfrak{t})^*, \ll)$.

– Theorem [G., 2022] –

For any $\mathfrak{t} \in \mathfrak{T}(\mathfrak{G})$, the poset $(\mathfrak{t}^*, \preccurlyeq)$ is a **finite lattice**.

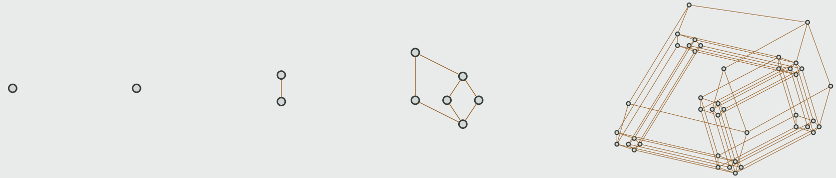
Mockingbird lattices

For any $h \geqslant 0$, the *h*-right comb tree is the $\mathfrak{G}$-term $\mathfrak{r}_h$ satisfying

$$\mathfrak{r}_h = \begin{cases} \mathbf{M} & \text{if } h = 0, \\ \mathbf{M}\,\mathfrak{r}_h & \text{otherwise.} \end{cases}$$

The Mockingbird lattice of order h is the lattice $\mathcal{M}(h) := (\mathfrak{r}_h^*, \preccurlyeq)$.

– Examples –



$\mathcal{M}(0)$ $\mathcal{M}(1)$ $\mathcal{M}(2)$ $\mathcal{M}(3)$ $\mathcal{M}(4)$

For any $h \geqslant 0$, the *h*-ladder is the duplicative tree $\mathfrak{l}_h$ satisfying

$$\mathfrak{l}_h = \begin{cases} \epsilon & \text{if } h = 0, \\ \circ(\mathfrak{l}_{h-1}) & \text{otherwise.} \end{cases}$$

When $h \geqslant 1$, the lattice $\mathcal{M}(h)$ is isomorphic to $\left(\mathfrak{l}_{h-1}^*, \ll\right)$.

3. Enumerative properties

Formal series of duplicative forests

Let $\mathbb{Z}\langle\langle\mathcal{D}^{\mathrm{F}}\rangle\rangle$ be the set of the possibly infinite **formal sums** of duplicative forests with coefficients in $\mathbb{Z}$. These are formal series on $\mathcal{D}^{\mathrm{F}}$ with integer coefficients.

For any $\mathbf{F} \in \mathbb{Z}\langle\langle\mathcal{D}^{\mathrm{F}}\rangle\rangle$ and $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$, the coefficient of $\mathfrak{f}$ in $\mathbf{F}$ is denoted by $\langle \mathfrak{f}, \mathbf{F}\rangle$.

Given a **statistics** $\omega : \mathcal{D}^{\mathrm{F}} \to \mathbb{N}$ and $\mathbf{F} \in \mathbb{Z}\langle\langle\mathcal{D}^{\mathrm{F}}\rangle\rangle$, the ω-specialization of $\mathbf{F}$ is the **generating series**

$$\omega(\mathbf{F}) := \sum_{\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}} \langle \mathfrak{f}, \mathbf{F}\rangle \, z^{\omega(\mathfrak{f})}$$

of $\mathbb{Z}\langle\langle z\rangle\rangle$.

The characteristic series of a set F of duplicative forests is the formal series

$$\mathbf{c}(F) := \sum_{\mathfrak{f} \in F} \mathfrak{f}.$$

In particular, $\mathrm{ht}(\mathbf{c}(F))$ is the generating series of the forests of F enumerated w.r.t. their heights.

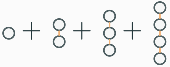
Strategy for the enumeration

– Strategy –

To **enumerate** a family F of duplicative forests, provide a **functional equation** describing $\mathbf{c}(F)$ and then specialize it to obtain a description of $\mathrm{ht}(\mathbf{c}(F))$.

An important ingredient for this is the series of the ladders

$$\mathbf{ld} := \sum_{h \geqslant 0} \mathfrak{l}_h = \epsilon +$$



$$+ \cdots .$$

Observe that $\mathbf{ld}$ satisfies the functional equation

$$\mathbf{ld} = \epsilon + \circ(\mathbf{ld}).$$

We deduce from this that $\mathrm{ht}(\mathbf{ld}) = 1 + z\,\mathrm{ht}(\mathbf{ld})$, implying

$$\mathrm{ht}(\mathbf{ld}) = \frac{1}{1-z}$$

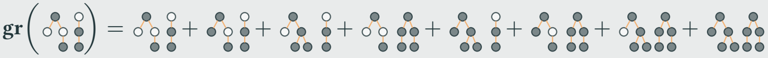
as expected.

Number of elements — 1/2

For any $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$, let

$$\mathbf{gr}(\mathfrak{f}) = \sum_{\mathfrak{f}' \in \mathfrak{f}^*} \mathfrak{f}'.$$

– Example –



By extending **gr** by **linearity**, $\mathbf{gr}(\mathbf{ld})$ is a well-defined formal series on $\mathcal{D}^{\mathrm{F}}$:

$$\mathbf{gr}(\mathbf{ld}) = \sum_{h \geqslant 0} \mathbf{gr}(\mathfrak{l}_h) = \sum_{h \geqslant 0} \sum_{\mathfrak{f} \in \mathfrak{l}_h^*} \mathfrak{f}.$$

The **ht-specialization** of the previous series satisfies

$$\mathrm{ht}(\mathbf{gr}(\mathbf{ld})) = \sum_{h \geqslant 0} (\# \mathfrak{l}_h^*) \, z^h.$$

Number of elements — 2/2

Let $\Delta : \mathbb{Z}\langle\langle \mathcal{D}^{\mathrm{F}} \rangle\rangle \to \mathbb{Z}\langle\langle \mathcal{D}^{\mathrm{F}} \rangle\rangle$ be the linear map satisfying $\Delta(\mathfrak{f}) = \mathfrak{f}\,\mathfrak{f}$ for any $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$.

– Theorem [G., 2022] –

The series $\mathbf{gr}(\mathbf{ld})$ satisfies

$$\mathbf{gr}(\mathbf{ld}) = \epsilon + \circ(\mathbf{gr}(\mathbf{ld})) + \bullet(\mathbf{gr}(\Delta(\mathbf{ld}))).$$

From this, we deduce that the ht-specialization A of $\mathbf{gr}(\mathbf{ld})$ satisfies

$$A = 1 + z\,A + z\,(A \boxtimes A),$$

where $\boxtimes$ is the **Hadamard product** of generating series.

The coefficients $a(h) := \langle z^h, A \rangle$ satisfy $a(0) = 1$ and for any $h \geqslant 1$,

$$a(h) = a(h-1) + a(h-1)^2.$$

The sequence $(a(h))_{h \geqslant 0}$ starts by

$$1, 2, 6, 42, 1806, 3263442, 10650056950806$$

and forms Sequence **A007018**.

Number of intervals — 1/3

For any $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$, let $\mathbf{ns}(\mathfrak{f}) = \mathbf{gr}(\mathbf{gr}(\mathfrak{f}))$. We obtain

$$\mathbf{ns}(\mathfrak{f}) = \sum_{\mathfrak{f}' \in \mathfrak{f}^*} \#[\mathfrak{f}, \mathfrak{f}'] \, \mathfrak{f}'.$$

– Example –

By extending **ns** by **linearity**, $\mathbf{ns}(\mathbf{ld})$ is a well-defined formal series on $\mathcal{D}^{\mathrm{F}}$:

$$\mathbf{ns}(\mathbf{ld}) = \sum_{h \geqslant 0} \mathbf{ns}(\mathfrak{l}_h) = \sum_{h \geqslant 0} \sum_{\mathfrak{f} \in \mathfrak{l}_h^*} \#[\mathfrak{l}_h, \mathfrak{f}] \, \mathfrak{f}.$$

The **ht-specialization** of the previous series satisfies

$$\mathrm{ht}(\mathbf{ns}(\mathbf{ld})) = \sum_{h \geqslant 0} I_h \, z^h$$

where I_h is the **number of intervals** of the poset $(\mathfrak{l}_h^*, \ll)$.

Number of intervals — 2/3

To provide a functional equation for **ns**(**ld**), we need the following two **tools**.

The *k*-meet decomposition, $k \geqslant 1$, of $\mathfrak{f} \in \mathcal{D}^{\mathrm{F}}$ is

$$\mathbf{md}_k(\mathfrak{f}) = \sum_{\substack{\mathfrak{g}_1, \dots, \mathfrak{g}_k \in \mathfrak{f}^* \\ \mathfrak{g}_1 \wedge \dots \wedge \mathfrak{g}_k = \mathfrak{f}}} \mathfrak{g}_1 \otimes \cdots \otimes \mathfrak{g}_k.$$

For any $u \in \{\circ, \bullet\}^*$, the merging product of $\mathfrak{f}_1, \dots, \mathfrak{f}_\ell \in \mathcal{D}^{\mathrm{F}}$ satisfies

$$\mathrm{mg}_{\circ u}(\mathfrak{f}_1 \otimes \mathfrak{f}_2 \otimes \cdots \otimes \mathfrak{f}_\ell) = \circ(\mathfrak{f}_1) \otimes \mathrm{mg}_u(\mathfrak{f}_2 \otimes \cdots \otimes \mathfrak{f}_\ell),$$

$$\mathrm{mg}_{\bullet u}(\mathfrak{f}_1 \otimes \mathfrak{f}_2 \otimes \mathfrak{f}_3 \otimes \cdots \otimes \mathfrak{f}_\ell) = \bullet(\mathfrak{f}_1 \mathfrak{f}_2) \otimes \mathrm{mg}_u(\mathfrak{f}_3 \otimes \cdots \otimes \mathfrak{f}_\ell).$$

– Example –

Number of intervals — 3/3

– Theorem [G., 2022] –

For any $k \geqslant 1$, the series $\mathbf{md}_k(\mathbf{ns}(\mathbf{ld}))$ satisfies

$$\mathbf{md}_k(\mathbf{ns}(\mathbf{ld})) = \epsilon^{k\otimes} + \sum_{u \in \{\circ, \bullet\}^k} \mathrm{mg}_u\left(\mathbf{md}_{k+|u|_\bullet}(\mathbf{ns}(\mathbf{ld}))\right) + \bullet(\mathbf{md}_k(\mathbf{ns}(\Delta(\mathbf{ld})))).$$

Since $\mathbf{md}_1(\mathbf{ns}(\mathbf{ld})) = \mathbf{ns}(\mathbf{ld})$, this provides a functional equation for $\mathbf{ns}(\mathbf{ld})$.

From this, we deduce that the ht-specialization A_k of $\mathbf{md}_k(\mathbf{ns}(\mathbf{ld}))$ satisfies

$$A_k = 1 + z(A_k \boxtimes A_k) + z \sum_{0 \leqslant i \leqslant k} \binom{k}{i} A_{k+i}.$$

The coefficients $a_k(h) := \left\langle z^h, A_k \right\rangle$ satisfy $a_k(0) = 1$ and for any $h \geqslant 1$,

$$a_k(h) = a_k(h-1)^2 + \sum_{0 \leqslant i \leqslant k} \binom{k}{i} a_{k+i}(h-1).$$

The sequence $(a_1(h))_{h \geqslant 0}$ starts by

$$1, 3, 17, 371, 144513, 20932611523, 438176621806663544657.$$

Conclusion and perspectives

We have studied a very simple CLS, the Mockingbird system, having nevertheless some **rich combinatorics**:

- its rewrite graphs are Hasse diagrams of posets;
- all intervals of these posets are lattices;
- these lattices are not graded, not self-dual, and not semidistributive;
- enumerative data is accessible but nontrivial.

Some **questions** and **projects**:

1. study, under an order theoretic point of view, some other CLS built from some basic combinators of the Enchanted forest of combinator birds;
2. provide necessary and/or sufficient conditions for a CLS to have the poset or the lattice property;
3. realize some well-known posets (like Tamari lattices, Stanley lattices, or Kreweras lattices) as intervals of posets built from specific CLSs.