

NATURAL HOPF ALGEBRAS OF OPERADS AND THEIR POLYNOMIAL REALIZATION

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Séminaire ADA, LMPA

April 11, 2024

Objectives

Present a **polynomial realization** of some Hopf algebras constructed from operads.

Main points:

1. Combinatorial Hopf algebras.
2. Polynomial realizations.
3. Nonsymmetric operads.
4. Natural Hopf algebras of nonsymmetric operads.
5. A polynomial realization of natural Hopf algebras of free operads.

Combinatorial Hopf algebras

All algebraic structures are over a field \mathbb{K} of characteristic zero.

A combinatorial Hopf algebra (CHA) \mathcal{H} is a graded vector space decomposing as

$$\mathcal{H} = \bigoplus_{n \in \mathbb{N}} \mathcal{H}(n)$$

endowed with

- an associative unital graded **product** $\star : \mathcal{H}(n_1) \otimes \mathcal{H}(n_2) \rightarrow \mathcal{H}(n_1 + n_2)$
- a coassociative counital cograded **coproduct** $\Delta : \mathcal{H}(n) \rightarrow \bigoplus_{n=n_1+n_2} \mathcal{H}(n_1) \otimes \mathcal{H}(n_2)$

such that each $\mathcal{H}(n)$ is finite dimensional, $\dim \mathcal{H}(0) = 1$, and

$$\Delta(x_1 \star x_2) = \Delta(x_1) \bar{\star} \Delta(x_2).$$

Let **FQSym** be the space such that **FQSym**(n) is the linear span of $\mathfrak{S}(n)$, the set of **permutations** of size n .

The set $\{F_\sigma : \sigma \in \mathfrak{S}\}$ is the fundamental basis of **FQSym**.

Let \star be the shifted shuffle product on **FQSym**.

Example — Product of FQSym on the F-basis

$$F_{231} \star F_{12} = F_{23145} + F_{23415} + F_{23451} + F_{24315} + F_{24351} + F_{24531} + F_{42315} + F_{42351} + F_{42531} + F_{45231}$$

Let Δ be the standardized deconcatenation coproduct on **FQSym**.

Example — Coproduct of FQSym on the F-basis

$$\Delta(F_{24351}) = F_\epsilon \otimes F_{24351} + F_1 \otimes F_{3241} + F_{12} \otimes F_{231} + F_{132} \otimes F_{21} + F_{1324} \otimes F_1 + F_{24351} \otimes F_\epsilon$$

This is the Malvenuto-Reutenauer CHA [Malvenuto, Reutenauer, 1995].

Polynomial realizations

For any alphabet A , let $\mathbb{K}\langle A \rangle$ be the space of noncommutative polynomials on A having a possibly **infinite support** but a **finite degree**.

Example — Some noncommutative polynomials

Set $A_{\mathbb{N}} := \{\mathbf{a}_0, \mathbf{a}_1, \mathbf{a}_2, \dots\}$.

□ An element in $\mathbb{K}\langle A_{\mathbb{N}} \rangle$:

$$\sum_{0 \leq i_1 < i_2} \mathbf{a}_{i_1} \mathbf{a}_{i_2} = \mathbf{a}_0 \mathbf{a}_1 + \mathbf{a}_0 \mathbf{a}_2 + \dots + \mathbf{a}_1 \mathbf{a}_2 + \mathbf{a}_1 \mathbf{a}_3 + \dots$$

□ An element which is not in $\mathbb{K}\langle A_{\mathbb{N}} \rangle$:

$$\sum_{n \geq 0} \mathbf{a}_0^n = 1 + \mathbf{a}_0 + \mathbf{a}_0^2 + \mathbf{a}_0^3 + \dots$$

The space $\mathbb{K}\langle A \rangle$ is endowed with the product of noncommutative polynomials.

A polynomial realization of a CHA \mathcal{H} is a quadruple $(\mathcal{A}, \sharp, r_A, \mathbb{A})$ such that

1. \mathcal{A} is a class of alphabets (a class of sets possibly endowed with relations).
2. \sharp is an associative operation of disjoint union on \mathcal{A} .
3. For any alphabet A of \mathcal{A} ,

$$r_A : \mathcal{H} \rightarrow \mathbb{K}\langle A \rangle$$

is an associative algebra morphism.

4. For any $x \in \mathcal{H}$ and any **mutually commuting** alphabets A' and A'' of \mathcal{A} in $\mathbb{K}\langle A' \sharp A'' \rangle$,

$$r_{A' \sharp A''}(x) = (r_{A'} \otimes r_{A''}) \circ \Delta(x).$$

5. \mathbb{A} is an alphabet of \mathcal{A} such that $r_{\mathbb{A}}$ is injective.

Point 4. offers a way to compute the coproduct of \mathcal{H} by expressing the realization of x on the sum of two alphabets. This is the **alphabet doubling trick**.

Let A be an alphabet **endowed with a total order** \preccurlyeq .

The **standardization** of $u \in A^*$ is the word of positive integers $\text{std}(u)$ such that

$$\text{std}(u)_i = \#\{j : j \leq i \text{ and } u_j \leq u_i\} + \#\{j : i < j \text{ and } u_i > u_j\}.$$

Example — Standardization of a word

Let on the alphabet $A_{\mathbb{N}}$ the total order relation \preccurlyeq satisfying $\mathbf{a}_{i_1} \preccurlyeq \mathbf{a}_{i_2}$ iff $i_1 \leq i_2$.

$$\text{std}(\mathbf{a}_7 \mathbf{a}_2 \mathbf{a}_0 \mathbf{a}_0 \mathbf{a}_2 \mathbf{a}_6 \mathbf{a}_2 \mathbf{a}_0 \mathbf{a}_4) = \text{std}(7 \underline{2} \underline{0} \underline{0} \underline{2} \underline{6} \underline{2} \underline{0} \underline{4}) = 9 \ 4 \ 1 \ 2 \ 5 \ 8 \ 6 \ 3 \ 7$$

Observations :

- std is a map from A^* to \mathfrak{S} ;
- std is surjective iff A is infinite;
- $\text{std}(u)$ is the unique permutation having the same inversion set as the one of u .

Let $r_A : \mathbf{FQSym} \rightarrow \mathbb{K}\langle A \rangle$ be the map defined by

$$r_A(F_\sigma) := \sum_{\substack{u \in A^* \\ \text{std}(u) = \sigma^{-1}}} u$$

Example — The polynomial of a basis element

$$r_{A_{\mathbb{N}}}(F_{312}) = \sum_{0 \leq i_1 < i_2 \leq i_3} \mathbf{a}_{i_2} \mathbf{a}_{i_3} \mathbf{a}_{i_1} = \mathbf{a}_1 \mathbf{a}_1 \mathbf{a}_0 + \mathbf{a}_1 \mathbf{a}_2 \mathbf{a}_0 + \cdots + \mathbf{a}_2 \mathbf{a}_2 \mathbf{a}_0 + \mathbf{a}_2 \mathbf{a}_3 \mathbf{a}_0 + \cdots + \mathbf{a}_2 \mathbf{a}_2 \mathbf{a}_1 + \cdots$$

The disjoint sum $A' \uplus A''$ of the totally ordered alphabets A' and A'' is the ordinal sum of A' and A'' .

Theorem [Duchamp, Hivert, Thibon, 2002]

The class of totally ordered alphabets together with the ordinal sum operation \uplus , the map r_A , and the alphabet $A_{\mathbb{N}}$ forms a polynomial realization of \mathbf{FQSym} .

Example — An alphabet doubling in FQSym

$$r_{A' \# A''}(\mathbf{F}_{312}) = \sum_{\substack{u \in (A' \# A'')^* \\ \text{std}(u) = 231}} u = \sum_{\substack{u_1, u_2, u_3 \in A' \# A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1$$

$$= \sum_{\substack{u_1, u_2, u_3 \in A' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_1, u_2 \in A', u_3 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_1, u_3 \in A', u_2 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_2, u_3 \in A', u_1 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1$$

$$+ \sum_{\substack{u_1 \in A', u_2, u_3 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_2 \in A', u_1, u_3 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_3 \in A', u_1, u_2 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1 + \sum_{\substack{u_1, u_2, u_3 \in A'' \\ u_1 \preccurlyeq u_2 \preccurlyeq u_3 \\ u_1 \neq u_2}} u_2 u_3 u_1$$

$$= r_{A'}(\mathbf{F}_{312}) \otimes r_{A''}(\mathbf{F}_\epsilon) + r_{A'}(\mathbf{F}_{21}) \otimes r_{A''}(\mathbf{F}_1) + 0 + 0$$

$$+ r_{A'}(\mathbf{F}_1) \otimes r_{A''}(\mathbf{F}_{12}) + 0 + 0 + r_{A'}(\mathbf{F}_\epsilon) \otimes r_{A''}(\mathbf{F}_{312})$$

$$= (r_{A'} \otimes r_{A''}) \circ \Delta(\mathbf{F}_{312})$$

The following CHAs admit polynomial realizations:

- **NCSF**, the noncommutative symmetric functions CHA [Gelfand, Krob, Lascoux, Leclerc, Retakh, Thibon, 1995];
- **FQSym**, the Malvenuto-Reutenauer CHA [Duchamp, Hivert, Thibon, 2002];
- **PBT**, the Loday-Ronco CHA [Hivert, Novelli, Thibon, 2005];
- **WQSym**, the packed words CHA [Novelli, Thibon, 2006];
- **PQSym**^{*}, the dual parking functions CHA [Novelli, Thibon, 2007];
- **UBP**, the uniform block permutations CHA [Maurice, 2013];
- *CK* and **NCK**, the commutative and noncommutative Connes-Kreimer CHAs [Foissy, Novelli, Thibon, 2014];
- **H_{FG}**, the CHA on Feynman graphs [Foissy, 2020].

There are several advantages to polynomial realizations.

1. They produce from a CHA \mathcal{H} a family of polynomials **generalizing symmetric functions**.
2. They lead to a **unified encoding** of the elements of a CHA \mathcal{H} by polynomials, no matter how complex are the product and coproduct of \mathcal{H} .
3. They lead to links between a CHA \mathcal{H} and other CHAs by **specializing the alphabet** on which \mathcal{H} is realized.
4. Given a space \mathcal{A} endowed with a product and a coproduct, the existence of a polynomial realization of \mathcal{A} **proves** that \mathcal{A} is a **Hopf algebra**.

Nonsymmetric operads

A nonsymmetric operad (operad) is a set

$$\mathcal{O} = \bigsqcup_{n \in \mathbb{N}} \mathcal{O}(n)$$

endowed with

□ a unit $\mathbb{1} \in \mathcal{O}(1)$;

□ a composition map $-[-, \dots, -] : \mathcal{O}(n) \times \mathcal{O}(m_1) \times \dots \times \mathcal{O}(m_n) \rightarrow \mathcal{O}(m_1 + \dots + m_n)$

such that

$$\mathbb{1}[x] = x = x[\mathbb{1}, \dots, \mathbb{1}]$$

and

$$x[y_1, \dots, y_n][z_{1,1}, \dots, z_{1,m_1}, \dots, z_{n,1}, \dots, z_{n,m_n}] = x[y_1[z_{1,1}, \dots, z_{1,m_1}], \dots, y_n[z_{n,1}, \dots, z_{n,m_n}]].$$

The arity $\text{ar}(x)$ of $x \in \mathcal{O}$ is the unique integer n such that $x \in \mathcal{O}(n)$.

Let \mathcal{O} be an operad.

An element $x \in \mathcal{O}(n)$ is **finitely factorizable** if the set of pairs $(y, (z_1, \dots, z_n))$ satisfying

$$x = y[z_1, \dots, z_n]$$

is finite.

When all elements of \mathcal{O} are finitely factorizable, by extension, \mathcal{O} is **finitely factorizable**.

A map $\mathrm{dg} : \mathcal{O} \rightarrow \mathbb{N}$ is a **grading** of \mathcal{O} if

$$\square \mathrm{dg}^{-1}(0) = \{1\};$$

$$\square \text{ for any } y \in \mathcal{O}(n) \text{ and } z_1, \dots, z_n \in \mathcal{O},$$

$$\mathrm{dg}(y[z_1, \dots, z_n]) = \mathrm{dg}(y) + \mathrm{dg}(z_1) + \dots + \mathrm{dg}(z_n).$$

When such a map exists, \mathcal{O} is **graded**.

The nonsymmetric associative operad As is the operad such that

- $As(0) = \emptyset$ and for any $n \geq 1$, $As(n)$ is the set $\{\alpha_n\}$;
- the unit is α_1 ;
- the composition map satisfies

$$\alpha_n[\alpha_{m_1}, \dots, \alpha_{m_n}] = \alpha_{m_1 + \dots + m_n}.$$

Example — A composition in As

$$\alpha_4[\alpha_2, \alpha_2, \alpha_3, \alpha_1] = \alpha_{2+2+3+1} = \alpha_8$$

The map dg defined by $dg(\alpha_n) := n - 1$ is a grading of As .

The operad As is finitely factorizable.

Natural Hopf algebras of nonsymmetric operads

Let \mathcal{O} be an operad.

The **reduced** $\text{rd}(w)$ of $w \in \mathcal{O}^*$ is the word obtained by removing the letters $\mathbb{1}$ in w .

Example — Reduced word of As^*

$$\text{rd}(\alpha_2 \alpha_2 \alpha_1 \alpha_4 \alpha_1 \alpha_1) = \alpha_2 \alpha_2 \alpha_4$$

The **natural space** $\mathbf{N}(\mathcal{O})$ of \mathcal{O} is the linear span of the set of reduced elements of \mathcal{O}^* .

The set $\{E_w : w \in \text{rd}(\mathcal{O}^*)\}$ is the **elementary basis** of $\mathbf{N}(\mathcal{O})$.

If \mathcal{O} admits a grading dg , then $\mathbf{N}(\mathcal{O})$ becomes a **graded space** by setting

$$\text{dg}(E_{w_1 \dots w_\ell}) := \text{dg}(w_1) + \dots + \text{dg}(w_\ell).$$

Note that $\text{dg}(E_\epsilon) = 0$.

Let \star be the **product** on $\mathbf{N}(\mathcal{O})$ defined by

$$E_{w_1} \star E_{w_2} := E_{w_1 w_2}.$$

Let Δ be the **coproduct** on $\mathbf{N}(\mathcal{O})$ defined by

$$\Delta(E_x) = \sum_{n \geq 0} \sum_{\substack{(y,w) \in \mathcal{O}(n) \times \mathcal{O}^n \\ x = y[w_1, \dots, w_n]}} E_{\text{rd}(y)} \otimes E_{\text{rd}(w)}.$$

Theorem [van der Laan, 2004] [Méndez, Liendo, 2014]

For any finitely factorizable operad \mathcal{O} , $\mathbf{N}(\mathcal{O})$ is a bialgebra.

If \mathcal{O} is graded, then $\mathbf{N}(\mathcal{O})$ is a Hopf algebra.

$\mathbf{N}(\mathcal{O})$ is the **natural Hopf algebra** of \mathcal{O} .

Let us apply this construction on As endowed with the grading dg satisfying $dg(\alpha_n) = n - 1$.

For any $n \geq 1$, $\dim \mathbf{N}(As)(n) = 2^{n-1}$.

Example — A product in $\mathbf{N}(As)$

$$E_{\alpha_3 \alpha_2 \alpha_2 \alpha_5} \star E_{\alpha_4 \alpha_2} = E_{\alpha_3 \alpha_2 \alpha_2 \alpha_5 \alpha_4 \alpha_2}$$

Example — A coproduct in $\mathbf{N}(As)$

$$\Delta(E_{\alpha_4}) = E_\epsilon \otimes E_{\alpha_4} + 2E_{\alpha_2} \otimes E_{\alpha_3} + E_{\alpha_2} \otimes E_{\alpha_2 \alpha_2} + 3E_{\alpha_3} \otimes E_{\alpha_2} + E_{\alpha_4} \otimes E_\epsilon.$$

Contributions to the coefficient 2 of $E_{\alpha_2} \otimes E_{\alpha_3}$:

$$\alpha_4 = \alpha_2[\alpha_1, \alpha_3], \quad \alpha_4 = \alpha_2[\alpha_3, \alpha_1].$$

Contributions to the coefficient 3 of $E_{\alpha_3} \otimes E_{\alpha_2}$:

$$\alpha_4 = \alpha_3[\alpha_1, \alpha_1, \alpha_2], \quad \alpha_4 = \alpha_3[\alpha_1, \alpha_2, \alpha_1], \quad \alpha_4 = \alpha_3[\alpha_2, \alpha_1, \alpha_1].$$

$\mathbf{N}(As)$ is the **noncommutative Faà di Bruno Hopf algebra FdB** [Figuerola, Gracia-Bondía, 2005] [Foissy, 2008].

Terms and forests

A **signature** is a set \mathcal{S} decomposing as $\mathcal{S} = \bigsqcup_{n \geq 0} \mathcal{S}(n)$.

An **\mathcal{S} -term** is an **ordered rooted tree** decorated on \mathcal{S} such that an internal node decorated by $g \in \mathcal{S}(n)$ has exactly n children.

Let $\mathfrak{T}(\mathcal{S})$ be the set of \mathcal{S} -terms.

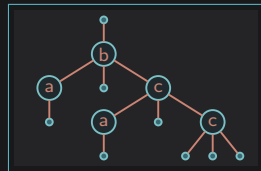
For any $t \in \mathfrak{T}(\mathcal{S})$,

- the **degree** $\text{dg}(t)$ of t is the number of internal nodes of t ;
- the **arity** $\text{ar}(t)$ of t is the number of leaves of t .

Example — An \mathcal{S} -term

Let the signature $\mathcal{S} := \mathcal{S}(1) \sqcup \mathcal{S}(3)$ with $\mathcal{S}(1) := \{a\}$ and $\mathcal{S}(3) := \{b, c\}$.

This \mathcal{S} -term has degree 5 and arity 7.

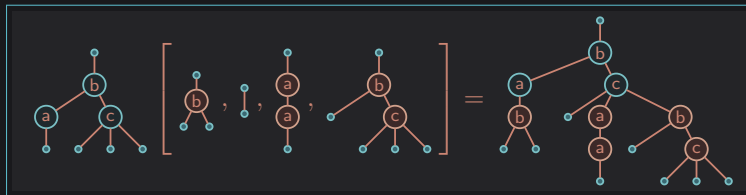


Let \mathcal{S} be a signature.

The free operad on \mathcal{S} is the set $\mathfrak{T}(\mathcal{S})$ such that

- $\mathfrak{T}(\mathcal{S})(n)$ is the set of \mathcal{S} -terms of arity n ;
- the unit is the \mathcal{S} -term containing exactly one leaf \bullet ;
- the composition map is such that $\mathfrak{t}[t_1, \dots, t_n]$ is the \mathcal{S} -term obtained by grafting simultaneously each t_i on the i -th leaf of \mathfrak{t} .

Example — A composition in a free operad



The map dg is a grading of $\mathfrak{T}(\mathcal{S})$ and this operad is finitely factorizable.

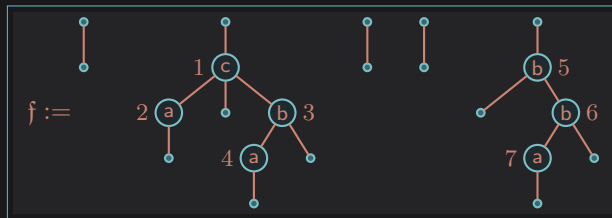
Let \mathcal{S} be a signature.

An \mathcal{S} -forest is a word on $\mathfrak{T}(\mathcal{S})$. Let $\mathfrak{F}(\mathcal{S})$ be the set of \mathcal{S} -forests.

The internal nodes of an \mathcal{S} -forest f are identified with their positions for the **preorder traversal**.

Let \xrightarrow{f}_j be the binary relation on the set of internal nodes of f such that $i_1 \xrightarrow{f}_j i_2$ if i_1 is the j -th child of i_2 in f .

Example — An \mathcal{S} -forest



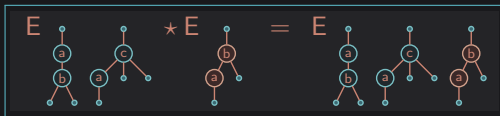
For instance, $1 \xrightarrow{f}_1 2$, $1 \xrightarrow{f}_3 3$, and $5 \xrightarrow{f}_2 6$.

Natural Hopf algebras of free operads

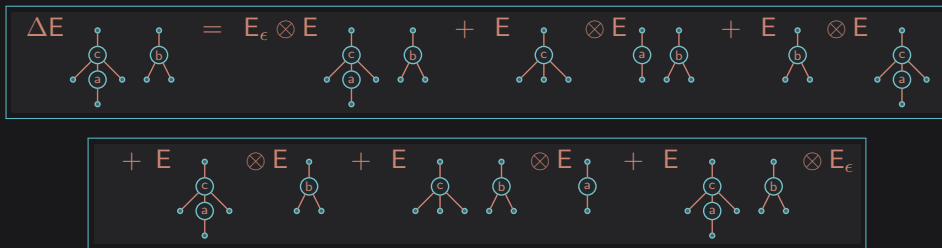
Let \mathcal{S} be a signature.

The bases of $\mathbf{N}(\mathfrak{T}(\mathcal{S}))$ are indexed by the set of **reduced \mathcal{S} -forests**.

Example — A product in $\mathbf{N}(\mathfrak{T}(\mathcal{S}))$



Example — A coproduct in $\mathbf{N}(\mathfrak{T}(\mathcal{S}))$



Polynomial realization

Let \mathcal{S} be a signature.

The class of \mathcal{S} -forest-like alphabets is the class of alphabets A endowed with relations R , D_g , and \prec_j such that

1. R is a unary relation called root relation;
2. for any $g \in \mathcal{S}$, D_g is a unary relation called g -decoration relation;
3. for any $j \geq 1$, \prec_j is a binary relation called j -edge relation.

Let \mathcal{S} be a signature, and A' and A'' be \mathcal{S} -forest-like alphabets.

The disjoint sum $A' \uplus A''$ of A' and A'' is the \mathcal{S} -forest-like alphabet

$$A := A' \sqcup A''$$

endowed with the relations R , D_g , and \prec_j such that

1. $R := R' \sqcup R''$;

2. $D_g := D'_g \sqcup D''_g$;

3. $a_1 \prec_j a_2$ holds if one of the three following conditions hold:

$$\square a_1 \in A', a_2 \in A', a_1 \prec_{j'} a_2; \quad \square a_1 \in A'', a_2 \in A'', a_1 \prec_{j''} a_2; \quad \square a_1 \in A', a_2 \in A'', a_2 \in R''.$$

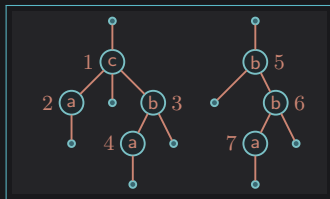
Let \mathcal{S} be a signature, A be an \mathcal{S} -forest-like alphabet, and \mathfrak{f} be a reduced \mathcal{S} -forest.

A word $w \in A^*$ is \mathfrak{f} -compatible, denoted by $w \Vdash^A \mathfrak{f}$, if

1. $\ell(w) = \text{dg}(\mathfrak{f})$;
2. if i is a root of \mathfrak{f} then $w_i \in R$;
3. if i is decorated by $g \in \mathcal{S}$ in \mathfrak{f} then $w_i \in D_g$;
4. if $i_1 \xrightarrow{\mathfrak{f}}_j i_2$ then $w_{i_1} \prec_j w_{i_2}$.

Example — An \mathfrak{f} -compatible word

Considering this reduced forest \mathfrak{f} , any \mathfrak{f} -compatible word $w \in A^*$ satisfies



- $\ell(w) = 7$;
- $w_1, w_5 \in R$;
- $w_2, w_4, w_7 \in D_a$, $w_3, w_5, w_6 \in D_b$, $w_1 \in D_c$;
- $w_1 \prec_1 w_2$, $w_1 \prec_3 w_3$, $w_3 \prec_1 w_4$, $w_5 \prec_2 w_6$, $w_6 \prec_1 w_7$.

Let \mathcal{S} be a signature and A be an \mathcal{S} -forest-like alphabet.

Let $r_A : \mathbf{N}(\mathfrak{T}(\mathcal{S})) \rightarrow \mathbb{K}\langle A \rangle$ be the linear map defined for any $f \in \text{rd}(\mathfrak{F}(\mathcal{S}))$ by

$$r_A(E_f) := \sum_{\substack{w \in A^* \\ w \Vdash^A f}} w.$$

This polynomial is the A -realization of f .

Lemma

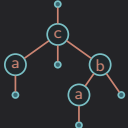
For any signature \mathcal{S} and any \mathcal{S} -forest-like alphabet A , r_A is an associative algebra morphism.

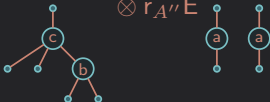
Lemma

For any signature \mathcal{S} , any \mathcal{S} -forest-like alphabets A' and A'' , and any \mathcal{S} -forest \mathfrak{f} ,

$$r_{A' \# A''}(E_{\mathfrak{f}}) = (r_{A'} \otimes r_{A''}) \circ \Delta(E_{\mathfrak{f}}).$$

Example — An alphabet doubling in $N(\mathfrak{T}(\mathcal{S}))$

$$r_{A' \# A''} E = \sum_{\substack{w_1, w_2, w_3, w_4 \in A' \# A'' \\ w_1 \in R \\ w_2, w_4 \in D_a, \quad w_1 \in D_c, \quad w_3 \in D_b \\ w_1 \prec_1 w_2, \quad w_1 \prec_3 w_3, \quad w_3 \prec_1 w_4}} w_1 w_2 w_3 w_4 = \cdots + \sum_{\substack{w_1, w_3 \in A', \quad w_2, w_4 \in A'' \\ w_1, w_2, w_4 \in R \\ w_2, w_4 \in D_a, \quad w_1 \in D_c, \quad w_3 \in D_b \\ w_1 \prec_3 w_3}} w_1 w_2 w_3 w_4 + \cdots$$


$$= \cdots + r_{A'} E \otimes r_{A''} E + \cdots$$


Let $\mathbb{P} := (\mathbb{N} \setminus \{0\})^*$.

The \mathcal{S} -forest-like alphabet of positions is the \mathcal{S} -forest-like alphabet

$$\mathbb{A}_p(\mathcal{S}) := \{\mathbf{a}_{g,u} : g \in \mathcal{S} \text{ and } u \in \mathbb{P}\}$$

such that

1. the root relation is defined by $R := \{\mathbf{a}_{g,u} \in \mathbb{A}_p(\mathcal{S}) : u = \epsilon\}$;
2. the g -decoration relation D_g is defined by $D_g := \{\mathbf{a}_{g',u} \in \mathbb{A}_p(\mathcal{S}) : g' = g\}$;
3. the j -edge relation \prec_j is defined by $\mathbf{a}_{g,u} \prec_j \mathbf{a}_{g',u.j}$.

Example — An alphabet $\mathbb{A}_p(\mathcal{S})$

Let the signature $\mathcal{S} := \mathcal{S}(1) \sqcup \mathcal{S}(3)$ such that $\mathcal{S}(1) = \{a, b\}$ and $\mathcal{S}(3) = \{c\}$. For instance,

- $\mathbf{a}_{a,\epsilon} \in R$, $\mathbf{a}_{b,121} \notin R$;
- $\mathbf{a}_{c,1761} \in D_c$, $\mathbf{a}_{c,\epsilon} \in D_c$, $\mathbf{a}_{b,211} \notin D_c$;
- $\mathbf{a}_{a,1212} \prec_1 \mathbf{a}_{b,12121}$, $\mathbf{a}_{c,2} \prec_2 \mathbf{a}_{a,22}$.

Example — $\mathbb{A}_p(\mathcal{S})$ -polynomials associated with some forests

$$r_{\mathbb{A}_p(\mathcal{S})} E \text{ (forest with root } b \text{ and one child)} = a_{b,\epsilon}$$

$$r_{\mathbb{A}_p(\mathcal{S})} E \text{ (forest with roots } b \text{ and } a \text{, each with one child)} = a_{b,\epsilon} a_{a,\epsilon}$$

$$r_{\mathbb{A}_p(\mathcal{S})} E \text{ (forest with root } b \text{ having child } a \text{, which has a child)} = a_{b,\epsilon} a_{a,1}$$

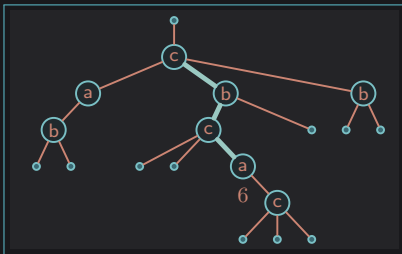
$$r_{\mathbb{A}_p(\mathcal{S})} E \text{ (forest with root } c \text{ having children } a \text{ and } b \text{, and } b \text{ having child } a \text{; and root } b \text{ having children } a \text{ and } b \text{, and } b \text{ having child } a) = a_{c,\epsilon} a_{a,1} a_{b,3} a_{a,31} a_{b,\epsilon} a_{b,2} a_{b,21}$$

Let \mathfrak{f} be an \mathcal{S} -forest.

The decoration $c_i(\mathfrak{f})$ of a node i of \mathfrak{f} is the element of \mathcal{S} decorating it.

The address $p_i(\mathfrak{f})$ of a node i of \mathfrak{f} is the word specifying the positions of the edges to reach i from the root.

Example — The decoration and the address of a node of an \mathcal{S} -term



For this \mathcal{S} -forest \mathfrak{f} , we have $c_6(\mathfrak{f}) = a$ and $p_6(\mathfrak{f}) = 213$.

Lemma

For any signature \mathcal{S} and any \mathcal{S} -forest \mathfrak{f} ,

$$r_{\mathbb{A}_P(\mathcal{S})}(E_{\mathfrak{f}}) = a_{c_1(\mathfrak{f}), p_1(\mathfrak{f})} \cdots a_{c_d(\mathfrak{f}), p_d(\mathfrak{f})}.$$

Lemma

For any signature \mathcal{S} , the map $r_{\mathbb{A}_p(\mathcal{S})} : \mathbf{N}(\mathfrak{T}(\mathcal{S})) \rightarrow \mathbb{K}\langle \mathbb{A}_p(\mathcal{S}) \rangle$ is injective.

Theorem [G., 2024+]

For any signature \mathcal{S} , the class of \mathcal{S} -forest-like alphabets, together with the alphabet disjoint sum operation \sharp , the map r_A , and the alphabet $\mathbb{A}_p(\mathcal{S})$, forms a polynomial realization of the Hopf algebra $\mathbf{N}(\mathfrak{T}(\mathcal{S}))$.

The \mathcal{S} -forest-like alphabet of extended positions is the \mathcal{S} -forest-like alphabet

$$\mathbb{A}_{\text{ep}}(\mathcal{S}) := \{\mathbf{a}_{g,u} : g \in \mathcal{S} \text{ and } u \in \mathbb{N}^*\}$$

such that

1. the root relation is defined by $R := \{\mathbf{a}_{g,u} \in \mathbb{A}_{\text{ep}}(\mathcal{S}) : u = 0^\ell, \ell \geq 0\}$;
2. the g -decoration relation D_g is defined by $D_g := \{\mathbf{a}_{g',u} \in \mathbb{A}_{\mathcal{S}} : g' = g\}$;
3. the j -edge relation \prec_j is defined by $\mathbf{a}_{g,u} \prec_j \mathbf{a}_{g',u} \text{ } j \text{ } 0^\ell$.

Example — An alphabet $\mathbb{A}_{\text{ep}}(\mathcal{S})$

Let the signature $\mathcal{S} := \mathcal{S}(1) \sqcup \mathcal{S}(3)$ such that $\mathcal{S}(1) = \{a, b\}$ and $\mathcal{S}(3) = \{c\}$.

For instance,

- $\mathbf{a}_{a,\epsilon} \in R, \mathbf{a}_{a,000} \in R, \mathbf{a}_{b,0100} \notin R, \mathbf{a}_{c,210000} \notin R$;
- $\mathbf{a}_{c,010761} \in D_c, \mathbf{a}_{c,000} \in D_c, \mathbf{a}_{b,20101} \notin D_c$;
- $\mathbf{a}_{a,10212} \prec_1 \mathbf{a}_{b,102121000}, \mathbf{a}_{c,00200} \prec_2 \mathbf{a}_{a,002002}$.

Example — $\mathbb{A}_{\text{ep}}(\mathcal{S})$ -polynomials associated with some forests

$$r_{\mathbb{A}_{\text{ep}}(\mathcal{S})} E \text{ (forest with root } b \text{ and two children)} = \sum_{\ell_1 \in \mathbb{N}} a_{b,0^{\ell_1}}$$

$$r_{\mathbb{A}_{\text{ep}}(\mathcal{S})} E \text{ (forest with roots } b \text{ and } a \text{, each with two children)} = \sum_{\ell_1, \ell_2 \in \mathbb{N}} a_{b,0^{\ell_1}} a_{a,0^{\ell_2}}$$

$$r_{\mathbb{A}_{\text{ep}}(\mathcal{S})} E \text{ (forest with root } b \text{ having child } a \text{, which has two children)} = \sum_{\ell_1, \ell_2 \in \mathbb{N}} a_{b,0^{\ell_1}} a_{a,0^{\ell_1}10^{\ell_2}}$$

$$r_{\mathbb{A}_{\text{ep}}(\mathcal{S})} E \text{ (forest with two components: one with root } c \text{ and children } a, b \text{; the other with root } b \text{ and children } b, a) = \sum_{\ell_1, \dots, \ell_7 \in \mathbb{N}} a_{c,0^{\ell_1}} a_{a,0^{\ell_1}10^{\ell_2}} a_{b,0^{\ell_1}30^{\ell_3}} a_{a,0^{\ell_1}30^{\ell_3}10^{\ell_4}} a_{b,0^{\ell_5}} a_{b,0^{\ell_5}20^{\ell_6}} a_{b,0^{\ell_5}20^{\ell_6}10^{\ell_7}}$$

Conclusion

Presented here:

- a **polynomial realization** r_A of $\mathbf{N}(\mathcal{O})$ when \mathcal{O} is a **free operad**.

Not presented here:

- when \mathcal{O} is **not free** and satisfies some properties, we can deduce from r_A a polynomial realization of $\mathbf{N}(\mathcal{O})$;
- by considering **alphabet specializations**, we can construct quotients or Hopf subalgebras of $\mathbf{N}(\mathfrak{T}(\mathcal{S}))$ which are isomorphic to
 - **NCK** and the noncommutative D -decorated Connes-Kreimer CHAs **NCK_D** [Foissy, 2002];
 - **FdB** and deformed noncommutative Faà di Bruno CHAs **FdB_r** [Foissy, 2008];
- fundamental $\{F_f\}$ and homogeneous $\{H_f\}$ bases, constructed through a new partial order on forests.

Topics to explore and open questions:

- study other alphabet specializations linking $\mathbf{N}(\mathcal{O})$ with other CHAs;
- conditions on the operad \mathcal{O} for the fact that $\mathbf{N}(\mathcal{O})$ is self-dual, free, or cofree;
- consider products of alphabets to define internal coproducts.