The Complexity of Transducer Synthesis from Multi-Sequential Specifications

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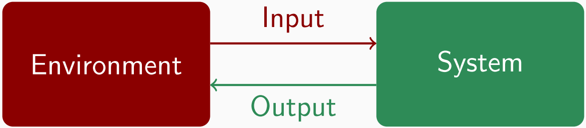
From verification to synthesis

Reactive systems



Interaction $\rightsquigarrow$ $i_1 o_1 i_2 o_2 i_3 o_3 \cdots \in (I\,O)^\omega$ or $(I\,O)^*$

Verification

Check that a system satisfies a specification

System $\|$ Env $\models$ Specification

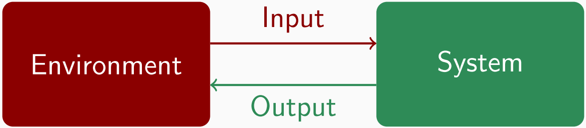
Synthesis

Generate a system from a specification

? $\|$ Env $\models$ Specification

From verification to synthesis

Reactive systems



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System $\parallel$ Env $\models$ Specification

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Generate a system from a specification

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Synthesis

$$? \parallel \text{Environment} \models \text{Specification}$$

➔ Generate a system from a specification

Implementing a specification

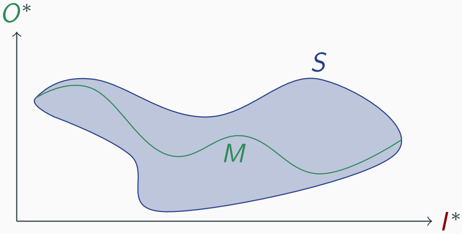
Input words I^*

Output words O^*

Implementation $M : I^* \to O^*$

Specification $S \subseteq I^* \times O^*$

M fulfils S, written $M \models S$, if for all $in \in I^*, (in, M(in)) \in S$



The realisability and synthesis problem

$\mathcal{S}$ = Class of specifications

$S \subseteq I^* \times O^*$

$\mathcal{M}$ = Class of target implementations

$M : I^* \to O^*$

Synthesis problem from $\mathcal{S}$ to $\mathcal{M}$

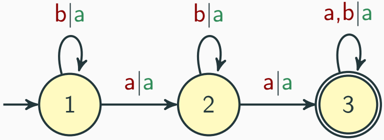
Input: Specification $S \in \mathcal{S}$

Output:
- Implementation $M \in \mathcal{M}$ s.t. $M \models S$ if it exists
- **No** otherwise

Realisability problem from $\mathcal{S}$ to $\mathcal{M}$

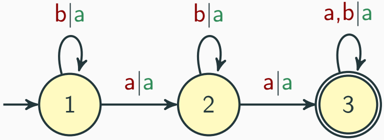
➔ Corresponding decision problem

Finite transducers: automata with outputs



Replace every letter with an a when there are at least two a's

Finite transducers: automata with outputs



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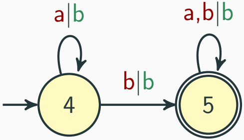

Replace every letter with a b when there is at least one b

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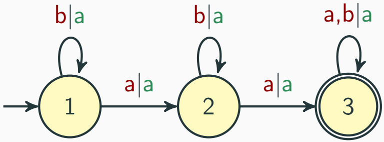

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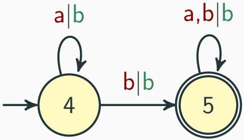
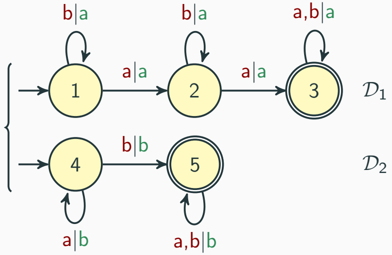


Replace every letter with a b when there is at least one b

Sequential transducer

The transition and output letter are *determined* by the input letter

Multi-sequential transducers



Running example

Multi-sequential transducer

Union of sequential transducers

$$\mathcal{T} = \biguplus_{i=1}^{k} \mathcal{D}_i$$

Multi-sequentiality

A relation is *multi-sequential* if it can be defined by a multi-sequential transducer

- Decidable for functions [?]
- Membership in **PTime** [?]

Transducer realisability problem
Known results

$\mathcal{M}$ = sequential transducers

$\mathcal{S}$	**Complexity**
MSO	Nonelementary [?]
LTL	2-ExpTime-c [?]
Finite Transducers	ExpTime-c

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LTL	2-ExpTime-c [?]
Finite Transducers	ExpTime-c

Question: Class of transducers with better complexity?

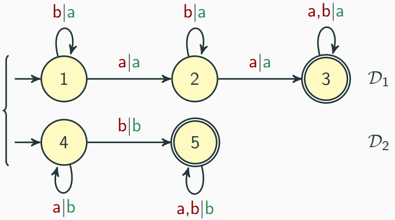
Realisability of multi-sequential specifications

$\mathcal{S}$ = Multi-seq. transducers

Unions of sequential transducers $\mathcal{T} = \uplus_{i=1}^{k} \mathcal{D}_i$

$\mathcal{M}$ = Seq. transducers

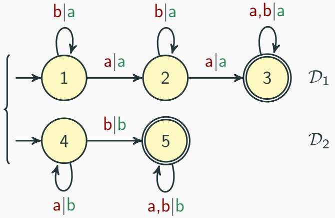
Output letter and transition is *determined* by input letter



Theorem

Sequential transducer synthesis from multi-sequential specifications is **PSpace**-complete.

Realisability of multi-sequential specifications
Proof



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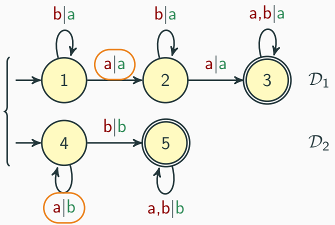

➔ On input a, need to *drop* one transducer

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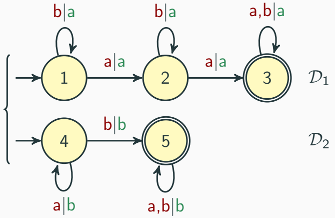

➔ On input a, need to *drop* one transducer

Critical prefix u

At least two runs on u disagree on their output

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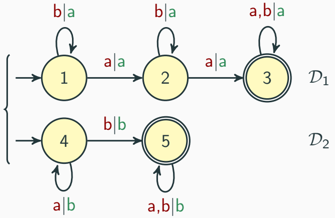

→ On input a, need to *drop* one transducer

Critical prefix u

At least two runs on u disagree on their output

Residual property

For all critical prefix u, there exists $P \subsetneq \{\mathcal{D}_1, \ldots, \mathcal{D}_k\}$ s.t.:

1. All transducers in P produce the same output on u
2. The domain is still covered: $u^{-1}\mathrm{dom}(\mathcal{T}) = \bigcup_{i \in P} u^{-1}\mathrm{dom}(\mathcal{D}_i)$
3. The residual specification $u^{-1} \left[\!\left[\biguplus_{i \in P} \mathcal{D}_i \right]\!\right]$ is realisable

Realisability of multi-sequential specifications
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Sequential transducer realisability from multi-sequential specifications is **PSpace**-complete.

Easiness

The *residual property* can be checked in **PSpace**.

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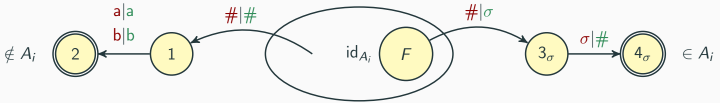
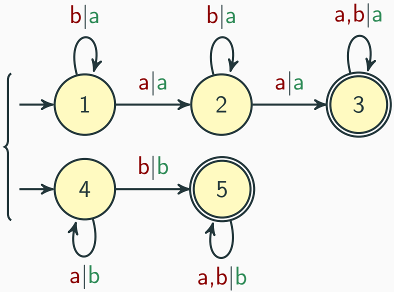
Hardness

$\rightsquigarrow$ Emptiness problem of the intersection of n DFAs

$$
S: \quad \begin{array}{lcll} w\#\sigma & \mapsto & w\sigma\# & \text{if } \exists i, w \in L(A_i) \qquad\qquad (\sigma \in \{a, b\}) \\ w\#\sigma & \mapsto & w\#\sigma & \text{if } \exists i, w \notin L(A_i) \end{array}
$$

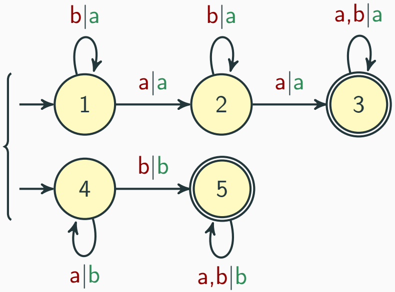


Asynchronicity



Our running example

Asynchronicity



Our running example

Waiting two steps allows to determine whether:

- There is at least one b
- There are at least two a's

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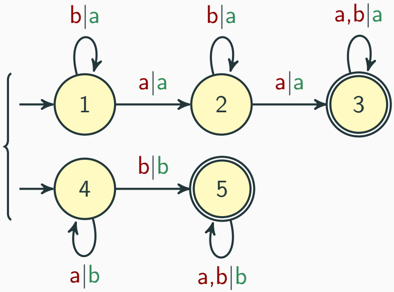

Our running example

Asynchronous transducer

At every transition, reads a letter, outputs a (possibly empty) word.

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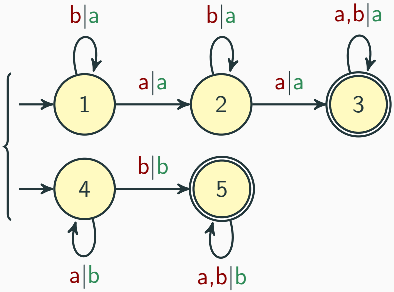

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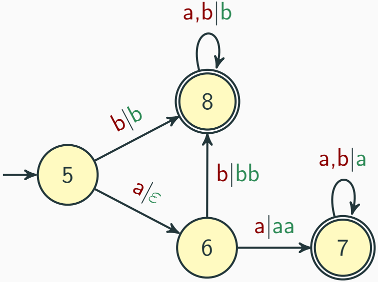

An asynchronous implementation

Asynchronous transducer realisability problem
Known results

$\mathcal{M}$ = Unambiguous functional transducers

Feasible for any asynchronous specification [?]

$\mathcal{M}$ = Sequential transducers

$\mathcal{S}$ (async. transducers)	**Complexity**
Nondeterministic	Undecidable [?]
Finite-valued	3-ExpTime [?]
Multi-sequential	PSpace-c

Asynchronous transducer realisability problem
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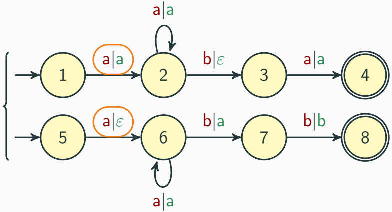
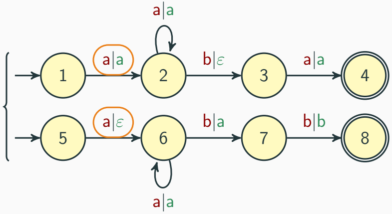


Delay

$$\mathrm{del}(u_1, u_2) = (\ell^{-1}u_1, \ell^{-1}u_2)$$

$$\ell = u_1 \wedge u_2$$

Asynchronous transducer realisability problem
Proof



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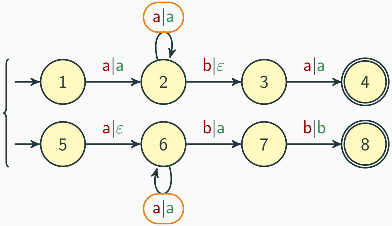

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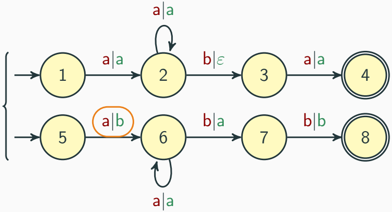
$$\begin{array}{ccc} \mathrm{del}(a, \varepsilon) & = & (a, \varepsilon) \\ \parallel & & \\ \mathrm{del}(aa, a) & = & (a, \varepsilon) \end{array}$$

Asynchronous transducer realisability problem
Proof



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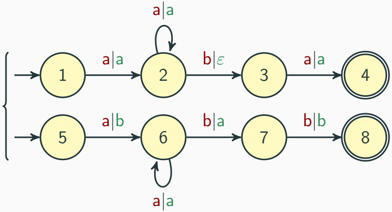
$$\mathrm{del}(a, b) = (a, b)$$

$$\nparallel$$

$$\mathrm{del}(aa, ba) = (aa, ba)$$

Asynchronous transducer realisability problem
Proof



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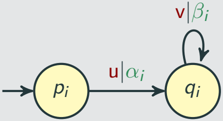
Critical loop

Triple (u, v, X) s.t.:

1. For all $\mathcal{D}_i \in X$,



2. For all $\mathcal{D}_i \notin X$, no run on u
3. For two transducers $\mathcal{D}_i, \mathcal{D}_j \in X$, delays accumulate: $\mathrm{del}(\alpha_i, \alpha_j) \neq \mathrm{del}(\alpha_i\beta_i, \alpha_j\beta_j)$

Asynchronous transducer realisability problem
Proof

Recursive characterisation

$\mathcal{T} = \uplus_{i=1}^{k} \mathcal{D}_i$ is realisable iff for all critical loops (u, v, X), there exists $Y \subsetneq X$ s.t.:

1. Delays do not accumulate:
   $\forall \mathcal{D}_i, \mathcal{D}_j \in Y, \mathrm{del}(\alpha_i, \alpha_j) = \mathrm{del}(\alpha_i\beta_i, \alpha_j\beta_j)$
2. The domain is still covered: $u^{-1}\mathrm{dom}(\mathcal{T}) = \bigcup_{i \in P} u^{-1}\mathrm{dom}(\mathcal{D}_i)$
3. The residual specification $(u, \ell)^{-1} \left[\!\left[\biguplus_{i \in Y} \mathcal{D}_i \right]\!\right]$ is realisable

ℓ longest common prefix of the α_i's

➔ Can be easily checked in **ExpTime**

Asynchronous transducer realisability problem
Proof

Theorem

Asynchronous sequential transducer synthesis from multi-sequential specifications is **PSpace**-complete.

PSpace-easiness: a non-recursive characterisation

Witness of non-satisfaction

- Unfolding of the recursive characterisation
- Reformulation of delay difference

➔ Can be found in **PSpace**

PSpace-hardness

➔ Similar to the synchronous case

Conclusion

Multi-sequential specifications

- Membership decidable in **PTime**
- Sequential realisability is **PSpace**-c both in synchronous and asynchronous cases

➔ Improvement of the general case:

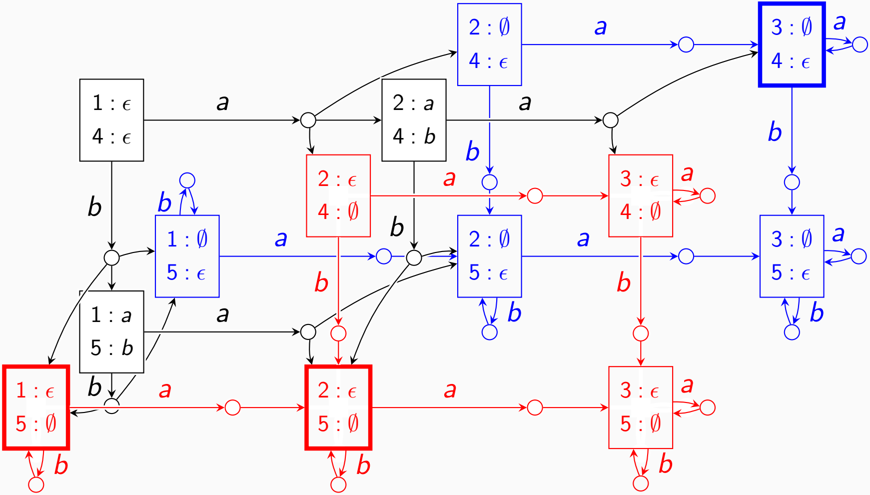
  - synchronous = **ExpTime-c**
  - asynchronous = **undecidable**

Synthesis game

➔ Practical synthesis algorithm

- Suitable for any type of specification defined by transducers (might not terminate)

The synthesis game



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