

Pattern avoidance in trees, operads, and enumeration

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Syntax trees and patterns

Operads and enumeration

Examples

Syntax trees and patterns

Syntax trees

An alphabet is a graded set $\mathfrak{G} := \bigsqcup_{n \geq 1} \mathfrak{G}(n)$.

Syntax trees

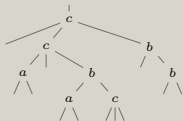
An alphabet is a graded set $\mathfrak{G} := \bigsqcup_{n \geq 1} \mathfrak{G}(n)$.

Let $S(\mathfrak{G})$ be the set of \mathfrak{G} -syntax trees defined recursively as

- $\mid \in S(\mathfrak{G})$, where \mid is the leaf;
- if $a \in \mathfrak{G}$ and $t_1, \dots, t_{|a|} \in S(\mathfrak{G})$, then $a(t_1, \dots, t_{|a|}) \in S(\mathfrak{G})$.

- Example -

Let $\mathfrak{G} := \mathfrak{G}(2) \sqcup \mathfrak{G}(3)$ such that $\mathfrak{G}(2) = \{a, b\}$ and $\mathfrak{G}(3) = \{c\}$.



denotes the \mathfrak{G} -tree

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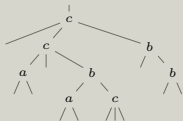
- ▶ $\mid \in S(\mathfrak{G})$, where \mid is the leaf;
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Let $t = a(t_1, \dots, t_{|a|}) \in S(\mathfrak{G})$. Some definitions:

- ▶ the degree $\deg(t)$ of t is its number of internal nodes;
- ▶ the arity $|t|$ of t is its number of leaves;
- ▶ for any $i \in [|a|]$, $t(i)$ is the i -th subtree t_i of t .

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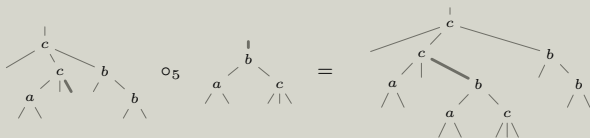
$c(\mid, c(a(\mid, \mid), \mid, b(a(\mid, \mid), c(\mid, \mid, \mid))), b(\mid, b(\mid, \mid)))$

having degree 8 and arity 12.

Compositions of syntax trees

Let $t, s \in \mathbf{S}(\mathcal{G})$. For each $i \in [|t|]$, the partial composition $t \circ_i s$ is the tree obtained by grafting the root of s onto the i -th leaf of t .

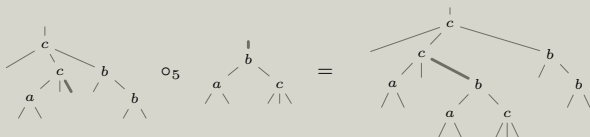
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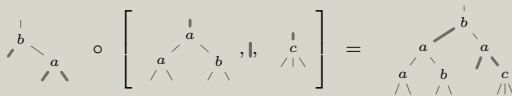
Let $t, s \in \mathbf{S}(\mathfrak{G})$. For each $i \in [|t|]$, the partial composition $t \circ_i s$ is the tree obtained by grafting the root of s onto the i -th leaf of t .

- Example -



Let $t, s_1, \dots, s_{|t|}$ be \mathfrak{G} -trees. The full composition $t \circ [s_1, \dots, s_{|t|}]$ is obtained by grafting simultaneously the roots of each s_i onto the i -th leaf of t .

- Example -



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Factors and prefixes

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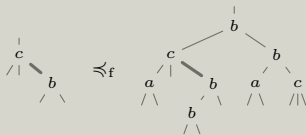
If t decomposes as

$$t = \tau \circ_i (s \circ [\tau_1, \dots, \tau_{|s|}])$$

for some trees $\tau, \tau_1, \dots, \tau_{|s|}$, and $i \in [| \tau |]$, then s is a factor of t .

This property is denoted by $s \preccurlyeq_f t$.

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If in the previous decomposition $\tau = |$, then

$$t = s \circ [\tau_1, \dots, \tau_{|s|}],$$

and s is a prefix of t .

This property is denoted by $s \preccurlyeq_p t$.

- Example -



Pattern avoidance and enumeration

A \mathcal{G} -tree t avoids (resp. prefix-avoids) a \mathcal{G} -tree s if $s \not\prec_f t$
(resp. $s \not\prec_p t$).

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Enumerate $A(\mathcal{P})$ w.r.t. the arities of the trees.

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- Examples -

► $A\left(\begin{array}{cccc} \begin{array}{c} a \\ \swarrow \quad \searrow \\ a \quad b \end{array} & \begin{array}{c} a \\ \swarrow \quad \searrow \\ b \quad a \end{array} & \begin{array}{c} b \\ \swarrow \quad \searrow \\ a \quad b \end{array} & \begin{array}{c} b \\ \swarrow \quad \searrow \\ b \quad a \end{array} \end{array}\right)$ is enumerated by $1, 2, 4, 8, 16, 32, 64, 128, \dots$

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► $A\left(\begin{array}{cccc} & a & & \\ & \swarrow \searrow & & \\ a & & c & \\ \swarrow \searrow & & \swarrow \searrow & \end{array}\right)$ is enumerated by $1, 1, 2, 4, 9, 21, 51, 127, \dots$ (A001006).

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Enumerate $A(\mathcal{P})$ w.r.t. the arities of the trees.

Let $\mathcal{P} \subseteq \mathbf{S}(\mathfrak{G}) \setminus \{\emptyset\}$ and $a \in \mathfrak{G}(k)$.

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Let \mathcal{P}_a be the subset of \mathcal{P} of the trees whose roots are labeled by a .

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A sequence $\mathcal{S} := (\mathcal{S}_1, \dots, \mathcal{S}_k)$, where each \mathcal{S}_i is a subset of $\mathbf{S}(\mathfrak{G})$, is \mathcal{P}_a -consistent if for any $\mathfrak{s} \in \mathcal{P}_a$, there is an $i \in [k]$ such that $\mathfrak{s}(i) \neq \mid$ and $\mathfrak{s}(i) \in \mathcal{S}_i$.

Consistent words

Let $\mathcal{P} \subseteq \mathbf{S}(\mathfrak{G}) \setminus \{\downarrow\}$ and $a \in \mathfrak{G}(k)$.

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- Example -

Let the set of patterns

$$\mathcal{P} := \left\{ \begin{array}{c} \downarrow \\ a \\ \swarrow \quad \searrow \\ c \quad \downarrow \end{array}, \begin{array}{c} \downarrow \\ c \\ \swarrow \quad \searrow \\ a \quad \downarrow \end{array}, \begin{array}{c} \downarrow \\ c \\ \swarrow \quad \searrow \\ b \quad b \\ \swarrow \quad \searrow \end{array}, \begin{array}{c} \downarrow \\ c \\ \swarrow \quad \searrow \\ b \quad a \\ \swarrow \quad \searrow \end{array}, \begin{array}{c} \downarrow \\ c \\ \swarrow \quad \searrow \\ c \quad a \\ \swarrow \quad \searrow \end{array} \right\}.$$

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The sequence

$$\mathcal{S} := \left(\left\{ \begin{array}{c} \perp \\ \diagup \quad \diagdown \\ a \quad \perp \end{array} \right\}, \left\{ \begin{array}{c} \perp \\ \diagup \quad \diagdown \\ b \quad \perp \end{array}, \begin{array}{c} \perp \\ \diagup \quad \diagdown \\ \perp \quad a \end{array} \right\}, \left\{ \begin{array}{c} \perp \\ \diagup \quad \diagdown \\ a \quad \perp \end{array}, \begin{array}{c} \perp \\ \diagup \quad \diagdown \\ \perp \quad a \end{array} \right\} \right)$$

is \mathcal{P}_c -consistent.

Let $\mathcal{P} \subseteq \mathbf{S}(\mathfrak{G}) \setminus \{\cdot\}$, $a \in \mathfrak{G}(k)$, and $\mathcal{S} := (\mathcal{S}_1, \dots, \mathcal{S}_k)$ be a \mathcal{P}_a -consistent sequence.

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and the \mathcal{P}_c -consistent word

$$\mathcal{S} := \left(\left\{ \begin{array}{c} | \\ a \\ / \backslash \end{array} \right\}, \left\{ \begin{array}{c} | \\ b \\ / \backslash \\ a \end{array}, \begin{array}{c} | \\ c \\ / \backslash \\ a \end{array} \right\}, \left\{ \begin{array}{c} | \\ a \\ / \backslash \\ a \end{array} \right\} \right).$$

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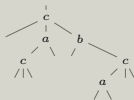
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The tree



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Minimal consistent words

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We have $\mathfrak{M}(\mathcal{P}_a) = \left\{ \left(\begin{array}{c} c \\ \swarrow \quad \searrow \\ a \quad b \end{array}, \emptyset \right) \right\},$

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We have $\mathfrak{M}(\mathcal{P}_a) = \left\{ \left(\left\{ \begin{array}{c} c \\ \swarrow \quad \searrow \\ a \quad b \end{array} \right\}, \emptyset \right) \right\}$, $\mathfrak{M}(\mathcal{P}_b) = \{(\emptyset, \emptyset)\}$,

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Let $\mathfrak{M}(\mathcal{P}_a)$ be the set of all minimal \mathcal{P}_a -consistent words.

- Examples -

Let the set of patterns

$$\mathcal{P} := \left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array}, \begin{array}{c} c \\ \swarrow \searrow \\ a \end{array}, \begin{array}{c} c \\ \swarrow \searrow \\ b \end{array}, \begin{array}{c} c \\ \swarrow \searrow \\ a \end{array}, \begin{array}{c} c \\ \swarrow \searrow \\ a \end{array} \right\}.$$

We have $\mathfrak{M}(\mathcal{P}_a) = \left\{ \left(\left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array} \right\}, \emptyset \right) \right\}$, $\mathfrak{M}(\mathcal{P}_b) = \{(\emptyset, \emptyset)\}$,

$$\mathfrak{M}(\mathcal{P}_c) = \left\{ \left(\left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array} \right\}, \left\{ \begin{array}{c} b \\ \swarrow \searrow \\ c \end{array} \right\}, \left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array} \right\} \right), \left(\left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array}, \begin{array}{c} b \\ \swarrow \searrow \\ c \end{array} \right\}, \emptyset, \left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array} \right\} \right), \right. \\ \left. \left(\left\{ \begin{array}{c} a \\ \swarrow \searrow \\ c \end{array}, \begin{array}{c} b \\ \swarrow \searrow \\ c \end{array}, \begin{array}{c} c \\ \swarrow \searrow \\ c \end{array} \right\}, \emptyset, \emptyset \right) \right\}.$$

- Lemma -

Let $\mathcal{P} \subseteq \mathbf{S}(\mathfrak{G}) \setminus \{|\}$ and t be a \mathfrak{G} -tree with root labeled by a .

The following assertions are equivalent:

1. t prefix-avoids \mathcal{P} ;
2. there exists a minimal \mathcal{P}_a -consistent word S such that t is S -admissible.

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- Lemma -

Let $\mathcal{P}, \mathcal{Q} \subseteq \mathbf{S}(\mathfrak{G}) \setminus \{|\}$ and t be a \mathfrak{G} -tree with root labeled by $a \in \mathfrak{G}(k)$.

The following assertions are equivalent:

1. t avoids \mathcal{P} and prefix-avoids \mathcal{Q} ;
2. for all $i \in [k]$, $t(i)$ avoid \mathcal{P} and there exists a minimal $(\mathcal{P} \cup \mathcal{Q})_a$ -consistent word S such that t is S -admissible.

Formal power series

Let \mathbb{K} be the field $\mathbb{Q}(q_0, q_1, q_2, \dots)$ and \mathcal{X} be a set.

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Endowed with the pointwise addition

$$\langle x, \mathbf{f} + \mathbf{g} \rangle := \langle x, \mathbf{f} \rangle + \langle x, \mathbf{g} \rangle$$

and the pointwise multiplication by a scalar

$$\langle x, \lambda \mathbf{f} \rangle := \lambda \langle x, \mathbf{f} \rangle,$$

the set $\mathbb{K} \langle\langle \mathcal{X} \rangle\rangle$ is a vector space.

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The sum notation of \mathbf{f} is

$$\mathbf{f} = \sum_{x \in \mathcal{X}} \langle x, \mathbf{f} \rangle x.$$

Tree series

A tree series is an element of $\mathbb{K}\langle\langle\mathbf{S}(\mathfrak{G})\rangle\rangle$.

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- Example -

For $x \in \mathfrak{G}$, let f_x be the $S(\mathfrak{G})$ -series wherein $\langle t, f_x \rangle$ is the number of occurrences of x in t . For instance,

$$f_a = \begin{array}{c} | \\ a \\ / \quad \backslash \end{array} + 2 \begin{array}{c} | \\ a \\ / \quad \backslash \\ a \quad \quad \backslash \end{array} + \begin{array}{c} | \\ a \\ / \quad \backslash \\ \quad b \quad \backslash \end{array} + \begin{array}{c} | \\ a \\ / \quad \backslash \\ \quad \quad b \end{array} + 2 \begin{array}{c} | \\ a \\ / \quad \backslash \\ a \quad \quad a \end{array} + 3 \begin{array}{c} | \\ a \\ / \quad \backslash \\ a \quad \quad a \\ / \quad \backslash \quad / \quad \backslash \end{array} + \dots$$

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- Example -

Let f_{\mid} be the $S(\mathfrak{G})$ -series wherein $\langle t, f_{\mid} \rangle := |t|$. Hence,

$$f_{\mid} = \mid + 2 \begin{array}{c} a \\ \diagup \quad \diagdown \end{array} + 2 \begin{array}{c} b \\ \diagup \quad \diagdown \end{array} + 3 \begin{array}{c} c \\ \diagup \quad \diagdown \end{array} + 3 \begin{array}{c} a \\ \diagup \quad \diagdown \\ a \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} + 3 \begin{array}{c} a \\ \diagup \quad \diagdown \\ b \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} + 3 \begin{array}{c} a \\ \diagup \quad \diagdown \\ a \quad \diagup \quad \diagdown \\ b \quad \diagup \quad \diagdown \end{array} + 3 \begin{array}{c} a \\ \diagup \quad \diagdown \\ a \quad \diagup \quad \diagdown \\ a \quad \diagup \quad \diagdown \end{array} + \dots$$

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- Example -

Let $f_{|}$ be the $S(\mathfrak{G})$ -series wherein $\langle t, f_{|} \rangle := |t|$. Hence,

$$f_{|} = | + 2 \begin{array}{c} | \\ \wedge \end{array} a + 2 \begin{array}{c} | \\ \wedge \end{array} b + 3 \begin{array}{c} | \\ \wedge \end{array} c + 3 \begin{array}{c} | \\ \wedge \end{array} \begin{array}{c} | \\ \wedge \end{array} a + 3 \begin{array}{c} | \\ \wedge \end{array} \begin{array}{c} | \\ \wedge \end{array} b + 3 \begin{array}{c} | \\ \wedge \end{array} \begin{array}{c} | \\ \wedge \end{array} b + 3 \begin{array}{c} | \\ \wedge \end{array} \begin{array}{c} | \\ \wedge \end{array} a + \dots$$

- Example -

In the tree series $f_a + f_b + f_c$, the coefficient of a tree is its degree.

In the tree series $f_{|} + f_a + f_b + f_c$, the coefficient of a tree is its number of edges.

Characteristic series

Let \mathcal{X} be a set and $\mathcal{S} \subseteq \mathcal{X}$.

The characteristic series of \mathcal{S} is the series

$$\mathbf{f}_{\mathcal{S}} := \sum_{x \in \mathcal{S}} x$$

of $\mathbb{K}\langle\langle\mathcal{X}\rangle\rangle$.

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- Lemma -

Let \mathcal{X} be a set and $\mathcal{S}_1, \dots, \mathcal{S}_n$, $n \geq 0$, be subsets of \mathcal{X} .

Then, the characteristic series of $\mathcal{S}_1 \cup \dots \cup \mathcal{S}_n$ expresses as

$$\mathbf{f}_{\mathcal{S}_1 \cup \dots \cup \mathcal{S}_n} = \sum_{\substack{\ell \geq 1 \\ \{i_1, \dots, i_\ell\} \subseteq [n]}} (-1)^{\ell+1} \mathbf{f}_{\mathcal{S}_{i_1} \cap \dots \cap \mathcal{S}_{i_\ell}}.$$

Enumeration map and compatible operations

Let \mathcal{X} be a set endowed with a size map $|-|: \mathcal{X} \rightarrow \mathbb{N}$.

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$$\text{en} : \mathbb{K} \langle\langle \mathcal{X} \rangle\rangle \rightarrow \mathbb{K} \langle\langle z \rangle\rangle$$

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When \mathcal{X} is combinatorial (that is, each fiber $|n|^{-1}$ is finite), $\text{en}(f_{\mathcal{X}})$ is the generating series of \mathcal{X} , enumerating its elements w.r.t. their sizes.

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A k -ary product $\star: \mathbb{K} \langle\langle \mathcal{X} \rangle\rangle^{\otimes k} \rightarrow \mathbb{K} \langle\langle \mathcal{X} \rangle\rangle$ is enumeration-compatible if

$$\text{en}(\star(\mathbf{f}_1, \dots, \mathbf{f}_k)) = \prod_{i \in [k]} \text{en}(\mathbf{f}_i)$$

for all \mathcal{X} -series $\mathbf{f}_1, \dots, \mathbf{f}_k$.

Composition of tree series

The composition of the $S(\mathfrak{G})$ -series f and g_1, \dots, g_n is the series

$$f \circ [g_1, \dots, g_n] := \sum_{\substack{t \in S(\mathfrak{G})(n) \\ s_1, \dots, s_n \in S(\mathfrak{G})}} \left(\langle t, f \rangle \prod_{i \in [n]} \langle s_i, g_i \rangle \right) t \circ [s_1, \dots, s_n].$$

Observe that this product is linear in all its arguments.

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$$\left(\begin{array}{c} | \\ a \\ / \backslash \end{array} + \begin{array}{c} | \\ b \\ / \backslash \\ \quad b \\ / \backslash \end{array} + \begin{array}{c} | \\ c \\ / \backslash \end{array} \right) \circ \left[|, \begin{array}{c} | \\ c \\ / \backslash \end{array}, \begin{array}{c} | \\ a \\ / \backslash \end{array} + \begin{array}{c} | \\ b \\ / \backslash \end{array} \right] = \begin{array}{c} | \\ b \\ / \backslash \\ c \quad a \\ / \backslash \end{array} + \begin{array}{c} | \\ b \\ / \backslash \\ c \quad b \\ / \backslash \end{array} + \begin{array}{c} | \\ c \\ / \backslash \\ a \quad \end{array} + \begin{array}{c} | \\ c \\ / \backslash \\ c \quad b \\ / \backslash \end{array}$$

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For any $t \in S(\mathfrak{G})(n)$, let $\bar{o}_t : \mathbb{K} \langle\langle S(\mathfrak{G}) \rangle\rangle^{\otimes n} \rightarrow \mathbb{K} \langle\langle S(\mathfrak{G}) \rangle\rangle$ be the product defined by

$$\bar{o}_t(g_1, \dots, g_n) := f_{\{t\}} \bar{o}[g_1, \dots, g_n].$$

for all tree series g_1, \dots, g_n .

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These products $\bar{\circ}_t$ are enumeration-compatible.

Series of trees avoiding patterns

For any $\mathcal{P}, \mathcal{Q} \subseteq \mathbf{S}(\mathfrak{G})$, let

$$\mathbf{F}(\mathcal{P}, \mathcal{Q}) := \sum_{\substack{t \in \mathbf{S}(\mathfrak{G}) \\ t \in A(\mathcal{P}) \\ \forall s \in \mathcal{Q}, s \not\prec_P t}} t.$$

This is the formal sum of all the \mathfrak{G} -trees avoiding as factors all patterns of \mathcal{P} and avoiding as prefixes all patterns of \mathcal{Q} .

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Since

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then the series $\mathbf{F}(\mathcal{P}, \mathcal{Q})$ contains all the enumerative data about the trees avoiding \mathcal{P} .

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then the series $\mathbf{F}(\mathcal{P}, \mathcal{Q})$ contains all the enumerative data about the trees avoiding \mathcal{P} .

We describe a system of equations for $\mathbf{F}(\mathcal{P}, \emptyset)$ using sums, multiplications by scalars, and operations \bar{o}_a , $a \in \mathfrak{G}$.

System of equations

When \mathfrak{G} , \mathcal{P} , and \mathcal{Q} satisfy some conditions, $\mathbf{F}(\mathcal{P}, \mathcal{Q})$ expresses as an inclusion-exclusion formula involving simpler terms $\mathbf{F}(\mathcal{P}, \mathcal{S}_i)$.

- Theorem [G., 2019] -

The series $\mathbf{F}(\mathcal{P}, \mathcal{Q})$ satisfies

$$\mathbf{F}(\mathcal{P}, \mathcal{Q}) = \mathbb{I} + \sum_{\substack{k \geq 1 \\ a \in \mathfrak{G}(k)}} \sum_{\substack{\ell \geq 1 \\ \{\mathcal{R}^{(1)}, \dots, \mathcal{R}^{(\ell)}\} \subseteq \mathfrak{M}((\mathcal{P} \cup \mathcal{Q})_a) \\ (\mathcal{S}_1, \dots, \mathcal{S}_k) = \mathcal{R}^{(1)} \dot{+} \dots \dot{+} \mathcal{R}^{(\ell)}}} (-1)^{1+\ell} \bar{o}_a(\mathbf{F}(\mathcal{P}, \mathcal{S}_1), \dots, \mathbf{F}(\mathcal{P}, \mathcal{S}_k)).$$

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The series $\mathbf{F}(\mathcal{P}, \mathcal{Q})$ satisfies

$$\mathbf{F}(\mathcal{P}, \mathcal{Q}) = 1 + \sum_{\substack{k \geq 1 \\ a \in \mathfrak{G}(k)}} \sum_{\substack{\ell \geq 1 \\ \{\mathcal{R}^{(1)}, \dots, \mathcal{R}^{(\ell)}\} \subseteq \mathfrak{M}((\mathcal{P} \cup \mathcal{Q})_a) \\ (\mathcal{S}_1, \dots, \mathcal{S}_k) = \mathcal{R}^{(1)} \dot{+} \dots \dot{+} \mathcal{R}^{(\ell)}}} (-1)^{1+\ell} \bar{o}_a(\mathbf{F}(\mathcal{P}, \mathcal{S}_1), \dots, \mathbf{F}(\mathcal{P}, \mathcal{S}_k)).$$

This leads to a system of equations for the generating series of $A(\mathcal{P})$.

Indeed, the generating series of $A(\mathcal{P})$ is the series $F(\mathcal{P}, \emptyset)$ where

$$F(\mathcal{P}, \mathcal{Q}) = z + \sum_{\substack{k \geq 1 \\ a \in \mathfrak{G}(k)}} \sum_{\substack{\ell \geq 1 \\ \{\mathcal{R}^{(1)}, \dots, \mathcal{R}^{(\ell)}\} \subseteq \mathfrak{M}((\mathcal{P} \cup \mathcal{Q})_a) \\ (\mathcal{S}_1, \dots, \mathcal{S}_k) = \mathcal{R}^{(1)} \dot{+} \dots \dot{+} \mathcal{R}^{(\ell)}}} (-1)^{1+\ell} \prod_{i \in [k]} F(\mathcal{P}, \mathcal{S}_i).$$

System of equations

- Example -

Let

$$\mathcal{P} := \left\{ \begin{array}{c} \text{a} \\ \swarrow \quad \searrow \\ \text{a} \quad \text{b} \end{array} \right\}.$$

We obtain the system of formal power series of trees

$$\begin{aligned} \mathbf{F}(\mathcal{P}, \emptyset) &= | + \bar{o}_a(\mathbf{F}(\mathcal{P}, \{a\}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_a(\mathbf{F}(\mathcal{P}, \emptyset), \mathbf{F}(\mathcal{P}, \{b\})) \\ &\quad - \bar{o}_a(\mathbf{F}(\mathcal{P}, \{a\}), \mathbf{F}(\mathcal{P}, \{b\})) + \bar{o}_b(\mathbf{F}(\mathcal{P}, \emptyset), \mathbf{F}(\mathcal{P}, \emptyset)), \\ \mathbf{F}(\mathcal{P}, \{a\}) &= | + \bar{o}_b(\mathbf{F}(\mathcal{P}, \emptyset), \mathbf{F}(\mathcal{P}, \emptyset)), \\ \mathbf{F}(\mathcal{P}, \{b\}) &= | + \bar{o}_a(\mathbf{F}(\mathcal{P}, \{a\}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_a(\mathbf{F}(\mathcal{P}, \emptyset), \mathbf{F}(\mathcal{P}, \{b\})) \\ &\quad - \bar{o}_a(\mathbf{F}(\mathcal{P}, \{a\}), \mathbf{F}(\mathcal{P}, \{b\})). \end{aligned}$$

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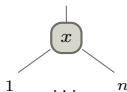
As a consequence, $F(\mathcal{P}, \emptyset)$ satisfies

$$z - F(\mathcal{P}, \emptyset) + (2 + z)F(\mathcal{P}, \emptyset)^2 - F(\mathcal{P}, \emptyset)^3 + F(\mathcal{P}, \emptyset)^4 = 0.$$

Operads and enumeration

Operators

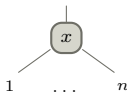
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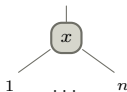
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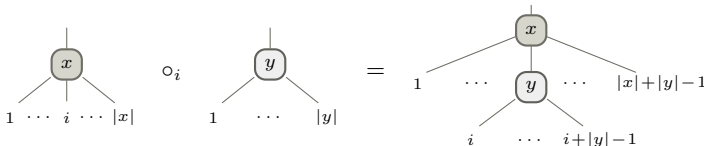


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This produces a new operator



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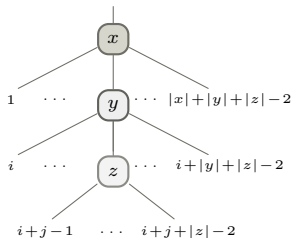
This data has to satisfy some axioms.

The associativity relation

$$(x \circ_i y) \circ_{i+j-1} z = x \circ_i (y \circ_j z)$$

$$1 \leq i \leq |x|, 1 \leq j \leq |y|$$

says that the pictured operation
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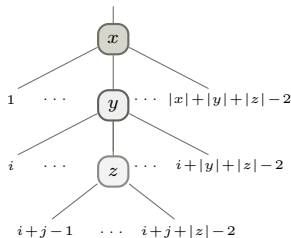
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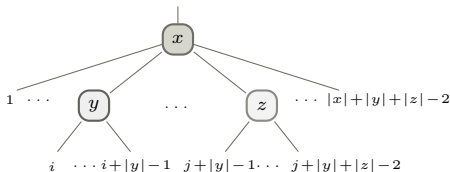


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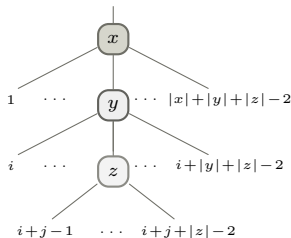
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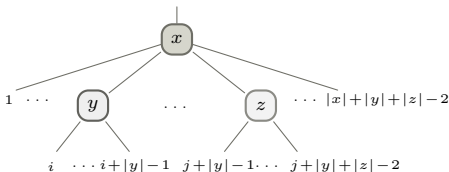


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The unitality relation

$$1 \circ_1 x = x = x \circ_i 1$$

$$1 \leq i \leq |x|$$

says that 1 is the identity map.



Let \mathcal{G} be an alphabet.

Free operads

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- elements of arity n are the \mathcal{G} -trees of arity n ;

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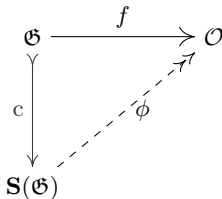
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Free operads satisfy the following universality property.

For any alphabet \mathfrak{G} , any operad \mathcal{O} , and any map $f: \mathfrak{G} \rightarrow \mathcal{O}$ preserving the arities, there exists a unique operad morphism $\phi: S(\mathfrak{G}) \rightarrow \mathcal{O}$ such that $f = \phi \circ c$.



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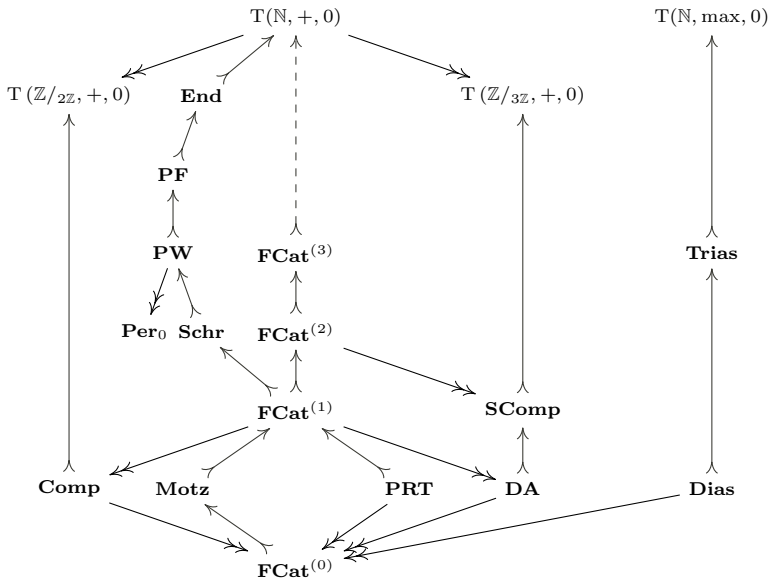
- Theorem [G., 2015] -

For any monoid \mathcal{M} , TM is an operad.

Some combinatorial suboperads

Monoid	Operad	Generators	First dimensions	Comb. objects
$(\mathbb{N}, +, 0)$	End	--	1, 4, 27, 256, 3125	Endofunctions
	PF	--	1, 3, 16, 125, 1296	Parking functions
	PW	--	1, 3, 13, 75, 541	Packed words
	Per₀	--	1, 2, 6, 24, 120	Permutations
	PRT	01	1, 1, 2, 5, 14, 42	Planar rooted trees
	FCat^(m)	00, 01, ..., 0m	Fuss-Catalan numbers	m-trees
	Schr	00, 01, 10	1, 3, 11, 45, 197	Schröder trees
	Motz	00, 010	1, 1, 2, 4, 9, 21, 51	Motzkin words
$(\mathbb{Z}/2\mathbb{Z}, +, 0)$	Comp	00, 01	1, 2, 4, 8, 16, 32	Compositions
$(\mathbb{Z}/3\mathbb{Z}, +, 0)$	DA	00, 01	1, 2, 5, 13, 35, 96	Directed animals
	SComp	00, 01, 02	1, 3, 27, 81, 243	Seg. compositions
$(\mathbb{N}, \max, 0)$	Dias	01, 10	1, 2, 3, 4, 5	Some bin. words
	Trias	00, 01, 10	1, 3, 7, 15, 31	Some bin. words

Diagram of operads



Let \mathcal{O} be an operad.

A congruence of \mathcal{O} is an equivalence relation \equiv on \mathcal{O} preserving the arities and such that $x \equiv x'$ and $y \equiv y'$ imply $x \circ_i y \equiv x' \circ_i y'$ for all $i \in [|x|]$.

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A basis of \mathcal{O} is a subset \mathcal{B} of $\mathbf{S}(\mathfrak{G})$ such that for any $[t]_{\equiv} \in \mathbf{S}(\mathfrak{G})/\equiv$, there exists a unique $s \in \mathcal{B}$ such that $s \in [t]_{\equiv}$.

Operads, presentations, and patterns

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- Link between bases and pattern avoidance -

In most cases, \mathcal{B} can be described as the set of \mathfrak{G} -trees avoiding a subset $\mathcal{P}_{\mathcal{B}}$ of $\mathbf{S}(\mathfrak{G})$.

Let $\mathcal{X} = \bigsqcup_{n \geq 1} \mathcal{X}(n)$ be a family of combinatorial objects, and consider that we want to describe the generating series

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By the previous results,

$$\mathbf{F}(\mathcal{P}_{\mathcal{B}}, \emptyset) = \mathbf{f}_{\mathcal{X}},$$

so that $\text{en}(\mathbf{F}(\mathcal{P}_{\mathcal{B}}, \emptyset))$ is the generating series of \mathcal{X} .

Examples

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Moreover, the set **Schr**(n) is in one-to-one correspondence with the set of Schröder trees with $n + 1$ leaves.

The first dimensions of **Schr** are hence

$$1, 3, 11, 45, 197, 903, 4279, 20793 \quad (\mathbf{A001003}).$$

Schröder trees

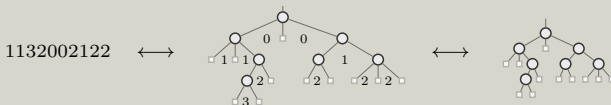
Let **Schr** be the suboperad of $T(\mathbb{N}, +, 0)$ generated by $\mathfrak{S} := \{00, 01, 10\}$.

- Proposition -

The set **Schr** contains exactly all the words u having at least a 0 and, for any letter $u_i \geq 1$, there is a letter $u_j = u_i - 1$ and the letters of the factor of u between u_i and u_j is made of letters $u_k \geq u_i$.

Moreover, the set **Schr**(n) is in one-to-one correspondence with the set of Schröder trees with $n + 1$ leaves.

- Example -



The first dimensions of **Schr** are hence

$$1, 3, 11, 45, 197, 903, 4279, 20793 \quad (\mathbf{A001003}).$$

Schröder trees

- Proposition -

The operad **Schr** admits the presentation (\mathfrak{S}, \equiv) where \equiv is the smallest operad congruence satisfying

$$c(00) \circ_1 c(00) \equiv c(00) \circ_2 c(00),$$

$$c(01) \circ_1 c(10) \equiv c(10) \circ_2 c(01),$$

$$c(00) \circ_1 c(01) \equiv c(00) \circ_2 c(10),$$

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Schröder trees

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$$c(01) \circ_1 c(01) \equiv c(01) \circ_2 c(00),$$

$$c(10) \circ_1 c(00) \equiv c(10) \circ_2 c(10).$$

- Proposition -

The set of the \mathfrak{S} -trees avoiding the set

$$\mathcal{P} := \left\{ \begin{array}{c} \text{00} \\ \text{00} \nearrow \text{00} \searrow \\ \text{00} \nearrow \text{00} \searrow \end{array}, \begin{array}{c} \text{01} \\ \text{10} \nearrow \text{01} \searrow \\ \text{10} \nearrow \text{01} \searrow \end{array}, \begin{array}{c} \text{00} \\ \text{01} \nearrow \text{00} \searrow \\ \text{01} \nearrow \text{00} \searrow \end{array}, \begin{array}{c} \text{01} \\ \text{00} \nearrow \text{01} \searrow \\ \text{00} \nearrow \text{01} \searrow \end{array}, \begin{array}{c} \text{00} \\ \text{10} \nearrow \text{00} \searrow \\ \text{10} \nearrow \text{00} \searrow \end{array}, \begin{array}{c} \text{01} \\ \text{01} \nearrow \text{01} \searrow \\ \text{01} \nearrow \text{01} \searrow \end{array}, \begin{array}{c} \text{10} \\ \text{10} \nearrow \text{10} \searrow \\ \text{10} \nearrow \text{10} \searrow \end{array} \right\}$$

is a basis of **Schr**.

Schröder trees

- Proposition -

The operad **Schr** admits the presentation (\mathfrak{S}, \equiv) where \equiv is the smallest operad congruence satisfying

$$c(00) \circ_1 c(00) \equiv c(00) \circ_2 c(00),$$

$$c(01) \circ_1 c(10) \equiv c(10) \circ_2 c(01),$$

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The set of the \mathfrak{S} -trees avoiding the set

$$\mathcal{P} := \left\{ \begin{array}{c} \text{00} \\ \diagup \quad \diagdown \\ \text{00} \end{array}, \begin{array}{c} \text{01} \\ \diagup \quad \diagdown \\ \text{10} \end{array}, \begin{array}{c} \text{00} \\ \diagup \quad \diagdown \\ \text{01} \end{array}, \begin{array}{c} \text{01} \\ \diagup \quad \diagdown \\ \text{00} \end{array}, \begin{array}{c} \text{00} \\ \diagup \quad \diagdown \\ \text{10} \end{array}, \begin{array}{c} \text{01} \\ \diagup \quad \diagdown \\ \text{01} \end{array}, \begin{array}{c} \text{10} \\ \diagup \quad \diagdown \\ \text{10} \end{array} \right\}$$

is a basis of **Schr**.

The characteristic series of this basis satisfies

$$\begin{aligned} \mathbf{F}(\mathcal{P}, \emptyset) &= | + \bar{o}_{00}(\mathbf{F}(\mathcal{P}, \mathfrak{S}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \mathfrak{S}), \mathbf{F}(\mathcal{P}, \emptyset)) \\ &\quad + \bar{o}_{10}(\mathbf{F}(\mathcal{P}, \emptyset), \mathbf{F}(\mathcal{P}, \{10\})), \end{aligned}$$

$$\mathbf{F}(\mathcal{P}, \mathfrak{S}) = |,$$

$$\mathbf{F}(\mathcal{P}, \{10\}) = | + \bar{o}_{00}(\mathbf{F}(\mathcal{P}, \mathfrak{S}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \mathfrak{S}), \mathbf{F}(\mathcal{P}, \emptyset)).$$

The generating series of these trees enumerated w.r.t. their arities (parameter z) and the numbers of internal nodes by types (parameters q_a , $a \in \mathfrak{G}$) satisfies

$$z + (z(q_{00} + q_{01} + q_{10}) - 1) F(\mathcal{P}, \emptyset) + (z(q_{00}q_{10} + q_{01}q_{10})) F(\mathcal{P}, \emptyset)^2 = 0.$$

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One has

$$\begin{aligned} F(\mathcal{P}, \emptyset) = & z + (q_{00} + q_{01} + q_{10}) z^2 \\ & + \left(q_{00}^2 + 2q_{00}q_{01} + 3q_{00}q_{10} + q_{01}^2 + 3q_{01}q_{10} + q_{10}^2 \right) z^3 \\ & + \left(q_{00}^3 + 3q_{00}^2q_{01} + 6q_{00}^2q_{10} + 3q_{00}q_{01}^2 + 12q_{00}q_{01}q_{10} + 6q_{00}q_{10}^2 \right. \\ & \quad \left. + q_{01}^3 + 6q_{01}^2q_{10} + 6q_{01}q_{10}^2 + q_{10}^3 \right) z^4 + \dots \end{aligned}$$

The generating series of these trees enumerated w.r.t. their arities (parameter z) and the numbers of internal nodes by types (parameters q_a , $a \in \mathfrak{G}$) satisfies

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The coefficients of the bivariate series obtained by specializing q_{10} and q_{01} (resp. q_{00} and q_{01}) to 1 are the ones of Triangle A126216 (resp. Triangle A114656).

For any $m \geq 0$, let $\mathbf{FCat}^{(m)}$ be the suboperad of $T(\mathbb{N}, +, 0)$ generated by $\mathfrak{G} := \{00, 01, \dots, 0m\}$.

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- Proposition -

The set $\mathbf{FCat}^{(m)}$ contains exactly all the words u satisfying $u_1 = 0$ and $u_i \in [0, u_{i-1} + m]$ for all valid positions i and $i - 1$.

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Moreover, the set $\mathbf{FCat}^{(m)}(n)$ is in one-to-one correspondence with the set of planar rooted trees with wherein all their n internal nodes have $m + 1$ children.

The dimensions of $\mathbf{FCat}^{(m)}$ are hence Fuss-Catalan numbers.

m -trees

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Moreover, the set $\mathbf{FCat}^{(m)}(n)$ is in one-to-one correspondence with the set of planar rooted trees with wherein all their n internal nodes have $m + 1$ children.

- Example -

For $m := 2$,



The dimensions of $\mathbf{FCat}^{(m)}$ are hence Fuss-Catalan numbers.

- Proposition -

The operad $\mathbf{FCat}^{(m)}$ admits the presentation (\mathfrak{G}, \equiv) where \equiv is the smallest operad congruence satisfying

$$c(0k_3) \circ_1 c(0k_1) \equiv c(0k_1) \circ_2 c(0k_2), \quad k_3 = k_1 + k_2.$$

- Proposition -

The operad $\mathbf{FCat}^{(m)}$ admits the presentation (\mathfrak{G}, \equiv) where \equiv is the smallest operad congruence satisfying

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- Proposition -

The set of the \mathfrak{G} -trees avoiding the set

$$\mathcal{P} := \left\{ \begin{array}{c} | \\ 0k_3 \\ / \quad \backslash \\ 0k_1 \quad \backslash \\ / \quad \backslash \end{array} : 0 \leq k_1 \leq k_3 \leq m \right\}$$

is a basis of $\mathbf{FCat}^{(m)}$.

- Proposition -

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The set of the \mathfrak{G} -trees avoiding the set

$$\mathcal{P} := \left\{ \begin{array}{c} | \\ 0k_3 \\ / \quad \backslash \\ 0k_1 \quad \backslash \\ / \quad \backslash \end{array} : 0 \leq k_1 \leq k_3 \leq m \right\}$$

is a basis of $\mathbf{FCat}^{(m)}$.

The characteristic series of this basis satisfies

$$\mathbf{F}(\mathcal{P}, \emptyset) = 1 + \sum_{0 \leq k \leq m} \bar{o}_{0k} (\mathbf{F}(\mathcal{P}, \mathcal{Q}_k), \mathbf{F}(\mathcal{P}, \emptyset)),$$

$$\mathbf{F}(\mathcal{P}, \mathcal{Q}_k) = 1 + \sum_{k < k' \leq m} \bar{o}_{0k'} (\mathbf{F}(\mathcal{P}, \mathcal{Q}_{k'}), \mathbf{F}(\mathcal{P}, \emptyset)),$$

where

$$\mathcal{Q}_k := \{00, 01, \dots, 0k\}.$$

The generating series of these trees enumerated w.r.t. their arities (parameter z) and the numbers of internal nodes by types (parameters q_a , $a \in \mathfrak{G}$) satisfies

$$-F(\mathcal{P}, \emptyset) + z \prod_{0 \leq k \leq m} (q_{0k} F(\mathcal{P}, \emptyset) + 1) = 0.$$

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$$-F(\mathcal{P}, \emptyset) + z \prod_{0 \leq k \leq m} (q_{0k} F(\mathcal{P}, \emptyset) + 1) = 0.$$

One has, for $m := 1$,

$$\begin{aligned} F(\mathcal{P}, \emptyset) &= z + (q_{00} + q_{01}) z^2 \\ &+ \left(q_{00}^2 + 2q_{00}q_{01} + 2q_{01}^2 \right) z^3 + \left(q_{00}^3 + 3q_{00}^2q_{01} + 5q_{00}q_{01}^2 + 5q_{01}^3 \right) z^4 \\ &+ \left(q_{00}^4 + 4q_{00}^3q_{01} + 9q_{00}^2q_{01}^2 + 14q_{00}q_{01}^3 + 14q_{01}^4 \right) z^5 \\ &+ \left(q_{00}^5 + 5q_{00}^4q_{01} + 14q_{00}^3q_{01}^2 + 28q_{00}^2q_{01}^3 + 42q_{00}q_{01}^4 + 42q_{01}^5 \right) z^6 + \cdots, \end{aligned}$$

The generating series of these trees enumerated w.r.t. their arities (parameter z) and the numbers of internal nodes by types (parameters q_a , $a \in \mathfrak{G}$) satisfies

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The coefficients of the bivariate series obtained by specializing q_{10} (resp. q_{00}) to 1 are the ones of Triangle A033184 (resp. Triangle A009766).

Let **DA** be the suboperad of $(\mathbb{Z}/3\mathbb{Z}, +, 0)$ generated by $\mathfrak{G} := \{00, 01\}$.

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- Proposition -

The set **DA**(n) is in one-to-one correspondence with the set of prefixes of Motkzin paths of $n - 1$ steps.

Directed animals

Let \mathbf{DA} be the suboperad of $(\mathbb{Z}/3\mathbb{Z}, +, 0)$ generated by $\mathfrak{G} := \{00, 01\}$.

- Proposition -

The set $\mathbf{DA}(n)$ is in one-to-one correspondence with the set of prefixes of Motkzin paths of $n - 1$ steps.

- Example -



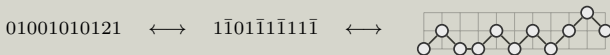
Directed animals

Let \mathbf{DA} be the suboperad of $(\mathbb{Z}/3\mathbb{Z}, +, 0)$ generated by $\mathfrak{G} := \{00, 01\}$.

- Proposition -

The set $\mathbf{DA}(n)$ is in one-to-one correspondence with the set of prefixes of Motzkin paths of $n - 1$ steps.

- Example -



Since prefixes of Motzkin paths are in one-to-one correspondence with directed animals on the square lattice [Gouyou-Beauchamps, Viennot, 1988], \mathbf{DA} is an operad on such objects.

The first dimensions of \mathbf{DA} are

$$1, 2, 5, 13, 35, 96, 267, 750, 2123 \quad (\mathbf{A005773}).$$

- Proposition -

The operad **DA** admits the presentation (\mathfrak{G}, \equiv) where \equiv is the smallest operad congruence satisfying

$$c(00) \circ_1 c(00) \equiv c(00) \circ_2 c(00),$$

$$c(01) \circ_1 c(00) \equiv c(00) \circ_2 c(01),$$

$$c(01) \circ_1 c(01) \equiv c(01) \circ_2 c(00),$$

$$(c(00) \circ_1 c(01)) \circ_2 c(01) \equiv (c(01) \circ_2 c(01)) \circ_3 c(01).$$

- Proposition -

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$$(c(00) \circ_1 c(01)) \circ_2 c(01) \equiv (c(01) \circ_2 c(01)) \circ_3 c(01).$$

- Proposition -

The set of the \mathfrak{G} -trees avoiding the set

$$\mathcal{P} := \left\{ \begin{array}{c} \text{00} \\ / \quad \backslash \\ \text{00} \end{array}, \begin{array}{c} \text{01} \\ / \quad \backslash \\ \text{00} \end{array}, \begin{array}{c} \text{01} \\ / \quad \backslash \\ \text{01} \end{array}, \begin{array}{c} \text{01} \\ / \quad \backslash \\ \text{01} \end{array} \right\}$$

is a basis of **DA**.

The characteristic series of the previous basis of **DA** satisfies

$$\mathbf{F}(\mathcal{P}, \emptyset) = \mathbf{1} + \bar{o}_{00}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, \mathbf{t}\})),$$

$$\mathbf{F}(\mathcal{P}, \{00\}) = \mathbf{1} + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, \mathbf{t}\})),$$

$$\mathbf{F}(\mathcal{P}, \{00, \mathbf{t}\}) = \mathbf{1} + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, 01, \mathbf{t}\})),$$

$$\mathbf{F}(\mathcal{P}, \{00, 01, \mathbf{t}\}) = \mathbf{1}.$$

where $\mathbf{t} := c(01) \circ_2 c(01)$.

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$$\mathbf{F}(\mathcal{P}, \{00\}) = \mathbf{1} + \bar{o}_{01} (\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, \mathfrak{t}\})),$$

$$\mathbf{F}(\mathcal{P}, \{00, \mathfrak{t}\}) = \mathbf{1} + \bar{o}_{01} (\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, 01, \mathfrak{t}\})),$$

$$\mathbf{F}(\mathcal{P}, \{00, 01, \mathfrak{t}\}) = \mathbf{1}.$$

where $\mathfrak{t} := c(01) \circ_2 c(01)$.

The generating series of these trees enumerated w.r.t. their arities (parameter z) and the numbers of internal nodes by types (parameters q_a , $a \in \mathfrak{G}$) satisfies

$$F(\mathcal{P}, \emptyset) = \frac{1 - \sqrt{1 - 2zq_{01} - 3z^2q_{01}^2} - z(2q_{00} + q_{01})}{2z(q_{00}^2 + q_{00}q_{01} + q_{01}^2) - 2q_{00}}$$

The characteristic series of the previous basis of **DA** satisfies

$$\begin{aligned} \mathbf{F}(\mathcal{P}, \emptyset) &= \mathbf{1} + \bar{o}_{00}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, \mathfrak{t}\})), \\ \mathbf{F}(\mathcal{P}, \{00\}) &= \mathbf{1} + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, \mathfrak{t}\})), \\ \mathbf{F}(\mathcal{P}, \{00, \mathfrak{t}\}) &= \mathbf{1} + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, 01, \mathfrak{t}\})), \\ \mathbf{F}(\mathcal{P}, \{00, 01, \mathfrak{t}\}) &= \mathbf{1}. \end{aligned}$$

where $\mathfrak{t} := c(01) \circ_2 c(01)$.

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One has

$$\begin{aligned} F(\mathcal{P}, \emptyset) &= z + (q_{00} + q_{01})z^2 + (q_{00}^2 + 2q_{00}q_{01} + 2q_{01}^2)z^3 \\ &\quad + (q_{00}^3 + 3q_{00}^2q_{01} + 5q_{00}q_{01}^2 + 4q_{01}^3)z^4 + (q_{00}^4 + 4q_{00}^3q_{01} + 9q_{00}^2q_{01}^2 + 12q_{00}q_{01}^3 + 9q_{01}^4)z^5 \\ &\quad + (q_{00}^5 + 5q_{00}^4q_{01} + 14q_{00}^3q_{01}^2 + 25q_{00}^2q_{01}^3 + 30q_{00}q_{01}^4 + 21q_{01}^5)z^6 + \dots \end{aligned}$$

The characteristic series of the previous basis of **DA** satisfies

$$\begin{aligned} \mathbf{F}(\mathcal{P}, \emptyset) &= 1 + \bar{o}_{00}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \emptyset)) + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, t\})), \\ \mathbf{F}(\mathcal{P}, \{00\}) &= 1 + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, t\})), \\ \mathbf{F}(\mathcal{P}, \{00, t\}) &= 1 + \bar{o}_{01}(\mathbf{F}(\mathcal{P}, \{00\}), \mathbf{F}(\mathcal{P}, \{00, 01, t\})), \\ \mathbf{F}(\mathcal{P}, \{00, 01, t\}) &= 1. \end{aligned}$$

where $t := c(01) \circ_2 c(01)$.

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$$F(\mathcal{P}, \emptyset) = \frac{1 - \sqrt{1 - 2zq_{01} - 3z^2q_{01}^2} - z(2q_{00} + q_{01})}{2z(q_{00}^2 + q_{00}q_{01} + q_{01}^2) - 2q_{00}}$$

One has

$$\begin{aligned} F(\mathcal{P}, \emptyset) &= z + (q_{00} + q_{01})z^2 + (q_{00}^2 + 2q_{00}q_{01} + 2q_{01}^2)z^3 \\ &\quad + (q_{00}^3 + 3q_{00}^2q_{01} + 5q_{00}q_{01}^2 + 4q_{01}^3)z^4 + (q_{00}^4 + 4q_{00}^3q_{01} + 9q_{00}^2q_{01}^2 + 12q_{00}q_{01}^3 + 9q_{01}^4)z^5 \\ &\quad + (q_{00}^5 + 5q_{00}^4q_{01} + 14q_{00}^3q_{01}^2 + 25q_{00}^2q_{01}^3 + 30q_{00}q_{01}^4 + 21q_{01}^5)z^6 + \dots \end{aligned}$$

The coefficients of the bivariate series obtained by specializing q_{10} (resp. q_{00}) to 1 are the ones of Triangle A064189 (resp. Triangle A026300).