# Operads of musical phrases and generation

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## Outline

1. Operads and constructions

2. The music box operad

3. Bud generating systems

## **Outline**

1. Operads and constructions

# Nonsymmetric operads

A **nonsymmetric operad** is a triple  $(\mathcal{O}, \circ, \mathbb{1})$  where

- lacksquare  $\mathcal{O}$  is a graded set  $\mathcal{O} = \bigsqcup_{n \geqslant 1} \mathcal{O}(n)$ ;
- lacksquare  $\circ$  :  $\mathcal{O}(n) \times (\mathcal{O}(m_1) \times \cdots \times \mathcal{O}(m_n)) \rightarrow \mathcal{O}(m_1 + \cdots + m_n)$  is a map called **full composition**;
- 1 is an element of  $\mathcal{O}(1)$  called **unit**.

The following relations have to hold:

• for all  $x \in \mathcal{O}(n)$ ,  $y_i \in \mathcal{O}(m_i)$ , and  $z_{i,j} \in \mathcal{O}$ ,  $i \in [n]$ ,  $j \in [m_i]$ ,

$$(x \circ y_1 \dots y_n) \circ z_{1,1} \dots z_{1,m_1} \dots z_{n,1} \dots z_{n,m_n} = x \circ (y_1 \circ z_{1,1} \dots z_{1,m_1}) \dots (y_n \circ z_{n,1} \dots z_{n,m_n});$$

• for all  $x \in \mathcal{O}(n)$ ,

$$1 \circ x = x = x \circ 1^n.$$

### Intuition

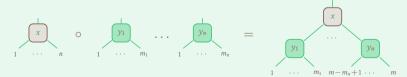
Each element x of  $\mathcal{O}(n)$  can be seen as a **planar operator**, that is an entity having n inputs and a single output:



The **arity** |x| of x is its number n of inputs, numbered from 1 to n from left to right.

The full composition  $x \circ y_1 \dots y_n$  consists in **grafting** the output of each  $y_i$  onto the *i*-th input of x.

This produces a new operator



of arity  $m := m_1 + \cdots + m_n$ .

# Partial and homogeneous composition maps

Let  $\mathcal{O}$  be an operad.

The **partial composition map** of  $\mathcal{O}$  is the map  $\circ_i : \mathcal{O}(n) \times \mathcal{O}(m) \to \mathcal{O}(n+m-1)$  defined, for any  $x \in \mathcal{O}(n)$ ,  $i \in [n]$ , and  $y \in \mathcal{O}(m)$ , by

$$x \circ_i y := x \circ \mathbb{1}^{i-1} y \mathbb{1}^{n-i}.$$

Conversely, we recover  $\circ$  from the  $\circ_i$  by setting, for any  $x \in \mathcal{O}(n)$  and  $y_1 \dots y_n \in \mathcal{O}^n$ ,

$$x \circ y_1 \dots y_n := (\dots (x \circ_n y_n) \circ_{n-1} y_{n-1} \dots) \circ_1 y_1.$$

The **homogeneous composition map** of  $\mathcal{O}$  is the map  $\odot$  :  $\mathcal{O}(n) \times \mathcal{O}(m) \to \mathcal{O}(nm)$  defined, for any  $x \in \mathcal{O}(n)$  and  $y \in \mathcal{O}(m)$ , by

$$x \odot y := x \circ y^n$$
.

The triple  $(\mathcal{O}, \odot, \mathbb{1})$  is a monoid.

# **Example: the duplicial operad**

The **duplicial operad** [Loday, 2008] is the operad (**Dup**,  $\circ_i$ ,  $\mathbb{1}$ ) such that

- for any  $n \ge 1$ , **Dup**(n) is the set of all binary trees with n internal nodes;
- for any  $\mathfrak{t} \in \mathbf{Dup}(n)$ ,  $i \in [n]$ , and  $\mathfrak{s} \in \mathbf{Dup}$ ,  $\mathfrak{t} \circ_i \mathfrak{s}$  is the binary tree obtained by replacing the i-th internal node u of  $\mathfrak{t}$  (for the infix traversal) by  $\mathfrak{s}$  and by grafting onto the leftmost (resp. rightmost) leaf of  $\mathfrak{s}$  the left (resp. right) child of u;
- 1 is the binary tree with exactly one internal node.

# - Example -

## The construction T

Let  $(\mathcal{M}, \star, e)$  be a monoid and let  $(T\mathcal{M}, \circ_i, \mathbb{1})$  be the triple such that

- for any  $n \ge 1$ ,  $T\mathcal{M}(n)$  is the set  $\mathcal{M}^n$ ;
- for any  $u \in T\mathcal{M}(n)$ ,  $i \in [n]$ , and  $v \in T\mathcal{M}$ ,

$$u \circ_i v := u(1, i-1) \ (u(i) \star v(1)) \dots (u(i) \star v(\ell(v))) \ u(i+1, \ell(u));$$

■ 1 is the element e seen as a word of length 1.

## - Examples -

Set 
$$\mathcal{M} := (\{\mathbf{a}, \mathbf{b}\}^*, ., \epsilon)$$
. In  $T\mathcal{M}$ ,

$$(aa, \epsilon, bab, \epsilon, b) \in T\mathcal{M}(5)$$

and

$$(\mathbf{b},\mathbf{ab},\epsilon,\mathbf{a})\circ_2(\epsilon,\mathbf{a},\mathbf{aa})=(\mathbf{b},\mathbf{ab}.\epsilon,\mathbf{ab}.\mathbf{a},\mathbf{ab}.\mathbf{aa},\epsilon,\mathbf{a})=(\mathbf{b},\mathbf{ab},\mathbf{aba},\mathbf{abaa},\epsilon,\mathbf{a}).$$

## - Theorem [G., 2015] -

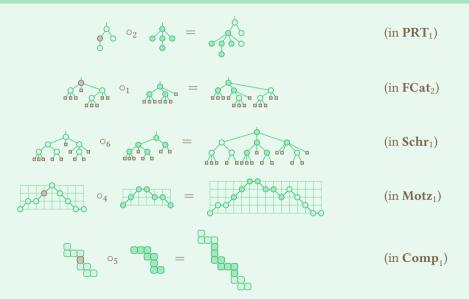
For any monoid  $\mathcal{M},$   $T\mathcal{M}$  is an operad.

# Operads from the construction T

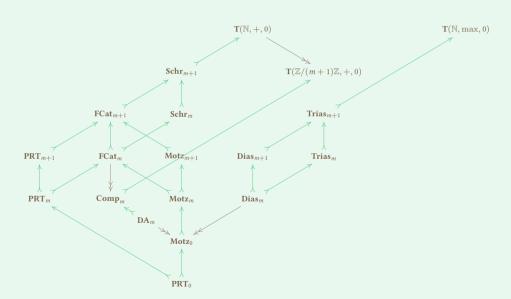
The operads  $T\mathcal{M}$  are large enough to contain a lot of suboperads realizable in combinatorial terms. As main examples:

- For any  $m \ge 0$ , with  $\mathcal{M} := (\mathbb{N}, +, 0)$ ,
  - **PRT**<sub>m</sub>, generated by  $\{01, \ldots, 0m\}$ , on primitive m-Dyck paths;
  - **FCat**<sub>m</sub>, gen. by  $\{00, 01, ..., 0m\}$ , on m-trees;
  - **Schr**<sub>m</sub>, gen. by  $\{01, \ldots, 0m, 00, m0, \ldots, 10\}$ , on some Schröder trees;
  - $Motz_m$ , gen. by  $\{00,000,010,\ldots,0m0\}$ , on colored Motzkin paths.
- For any  $m \ge 0$ , with  $\mathcal{M} := (\mathbb{Z}/(m+1)\mathbb{Z}, +, 0)$ ,
  - **Comp**<sub>m</sub>, gen. by  $\{00, 01, ..., 0m\}$ , on m-words;
  - **DA**<sub>m</sub>, gen. by  $\{00, 01, \dots, 0(m-1)\}$ , on some directed animals.
- For any  $m \ge 0$ ,  $\mathcal{M} := (\mathbb{N}, \max, 0)$ ,
  - $\mathbf{Dias}_m$ , gen. by  $\{01, \ldots, 0m, m0, \ldots, 10\}$ , is the m-pluriassociative operad [Loday, 2001] [G., 2016];
  - **Trias**<sub>m</sub>, gen. by  $\{01, \ldots, 0m, 00, m0, \ldots, 10\}$ , is the m-pluritriassociative operad [**Loday**, **Ronco**, 2004] [**G.**, 2016].

# Some partial compositions on combinatorial objects



# Full diagram



#### The construction U

Let  $(\mathcal{M}, \star, e)$  be a monoid and let  $(U\mathcal{M}, \circ_i, \mathbb{1})$  be the triple such that

- for any  $n \ge 1$ ,  $U\mathcal{M}(n)$  is the set  $\mathcal{M}^{n+1}$ ;
- for any  $u \in \mathbf{U}\mathcal{M}(n)$ ,  $i \in [n]$ , and  $v \in \mathbf{U}\mathcal{M}$ ,

$$u \circ_i v := u(1, i-1) \ (u(i) \star v(1)) \ v(2, \ell(v) - 1) \ (v(\ell(v)) \star u(i+1)) \ u(i+2, \ell(u)));$$

■ 1 is the element *ee*.

## - Examples -

Set 
$$\mathcal{M} := (\{\mathbf{a}, \mathbf{b}\}^*, ., \epsilon)$$
. In  $\mathbf{U}\mathcal{M}$ ,

$$(aa, \epsilon, bab, \epsilon, b) \in U\mathcal{M}(4)$$

and

$$(\mathbf{ba},\mathbf{aa},\mathbf{b},\epsilon,\mathbf{a})\circ_2(\mathbf{a},\mathbf{bb},\mathbf{b})=(\mathbf{ba},\mathbf{aa}.\mathbf{a},\mathbf{bb},\mathbf{b}.\mathbf{b},\epsilon,\mathbf{a})=(\mathbf{ba},\mathbf{aaa},\mathbf{bb},\mathbf{bb},\epsilon,\mathbf{a}).$$

## - **Theorem** [G., 2021-] -

For any monoid  $\mathcal{M},$   $U\mathcal{M}$  is an operad.

# Outline

2. The music box operad

# Degree patterns

A **degree** is an integer.

A **degree pattern** is a nonempty word **d** of degrees. The **arity**  $|\mathbf{d}|$  of **d** is its length.

Let  $\mathcal{N}$  be the set of all musical notes  $n_k$  where n is the pitch class and k is the octave of the note.

A degree interpretation is a map  $\rho: \mathbb{Z} \to \mathcal{N}$  sending each degree to a note. The  $\rho$ -interpretation  $\rho(\mathbf{d})$  of  $\mathbf{d}$  is the sequence  $\rho(\mathbf{d}(1)) \dots \rho(\mathbf{d}(|\mathbf{d}|))$  of notes.

## - Examples -

Let the degree pattern

$$\mathbf{d} := 10\overline{2}\overline{3}507.$$

A negative value has a bar above its absolute value. The arity of  ${\bf d}$  is 7.

If  $\rho$  sends 0 to 0<sub>4</sub> and the other degrees according with the minor pentatonic scale,

$$\rho(\mathbf{d}) = 3_4 \ 0_4 \ 7_3 \ 5_3 \ 0_5 \ 0_4 \ 5_5.$$

Instead, if  $\rho$  sends 0 to 0<sub>4</sub> and the other degrees according with the major natural scale,

$$\rho(\mathbf{d}) = 2_4 \ 0_4 \ 9_3 \ 7_3 \ 9_4 \ 0_4 \ 0_5.$$

# Rhythm patterns

A **rhythm pattern** is a nonempty word **r** on the alphabet  $\{\Box, \Box\}$  having at least one occurrence of  $\Box$ . The **arity**  $|\mathbf{r}|$  of **r** is its number of occurrences of  $\Box$ .

The **duration sequence** of a rhythm pattern  $\mathbf{r}$  is the unique sequence  $\sigma := \sigma(1) \dots \sigma(|\mathbf{r}| + 1)$  of nonnegative integers such that

$$\mathbf{r} = \square^{\sigma(1)} \square \square^{\sigma(2)} \dots \square \square^{\sigma(|\mathbf{r}|+1)}$$
.

A **rhythm interpretation** is a positive integer value  $\delta$  specifying the duration of the unit of time. The  $\delta$ -interpretation  $\delta(\mathbf{r})$  of  $\mathbf{r}$  is the sequence  $\sigma(1)\delta, (\sigma(2)+1)\delta, \ldots, (\sigma(|\mathbf{r}|+1)+1)\delta$ . It specifies the duration of the initial rest and the durations of the other beats.

## - Examples -

Let the rhythm pattern

$$\mathbf{r} := \square \square \square \square \square \square \square$$
.

The arity of  $\mathbf{r}$  is 3.

The duration sequence of  $\mathbf{r}$  is 2013 so that  $\mathbf{r}$  specifies the rhythm consisting in a rest lasting 2 units of time, a note lasting 1 unit of time, a note lasting 2 units of time, and a note lasting 4 units of time.

## **Patterns**

A **pattern p** is a pair  $(\mathbf{d}, \mathbf{r})$  such that  $|\mathbf{d}| = |\mathbf{r}|$ . The **arity**  $|\mathbf{p}|$  of **p** is the arity of **d** (or of **r**).

Patterns are denoted concisely by replacing each  $\blacksquare$  by the corresponding degree. In this way, patterns are words on the alphabet  $\{\Box\} \cup \mathbb{Z}$ .

— Example — The pattern  $\mathbf{p}:=(1\bar{1}2,\blacksquare \blacksquare \blacksquare \blacksquare \blacksquare)$  is of arity 3.

– Example –

The previous patterns writes as  $\mathbf{p} = 1\overline{1} \square \square \square 2 \square$ .

An **interpretation** is a pair  $(\rho, \delta)$  such that  $\rho$  is a degree interpretation and  $\delta$  is a rhythm interpretation. The  $(\rho, \delta)$ -interpretation of a pattern  ${\bf p}$  is the sequence of notes with their durations obtained from the  $\rho$ -interpretation of  ${\bf p}$  and the  $\delta$ -interpretation of  ${\bf p}$ .

By convention, in the following musical scores, each unit of time lasts  $\frac{1}{8}$  the duration of a whole note.

## – Example –

Let the pattern  $\mathbf{p} := 0$   $\square$   $12\overline{1}$   $\square$   $01\overline{2}$   $\square$   $\overline{100}$   $\square$   $\square$ . By setting  $\rho$  as the degree interpretation specifying the  $9_3$  harmonic minor scale and  $\delta$  as the rhythm interpretation specifying 192 as tempo, we obtain the musical phrase



# **Multi-patterns**

A **multi-pattern** is a nonempty sequence **m** of patterns such that for all  $i, j \in [\ell(\mathbf{m})]$ ,  $|\mathbf{m}(i)| = |\mathbf{m}(j)|$  and  $\ell(\mathbf{m}(i)) = \ell(\mathbf{m}(j))$ .

The **arity**  $|\mathbf{m}|$  of  $\mathbf{m}$  is the arity of any  $\mathbf{m}(i)$ .

The **length**  $\ell(\mathbf{m})$  of  $\mathbf{m}$  is the length of any  $\mathbf{m}(i)$ .

The **multiplicity**  $\mathfrak{m}(\mathbf{m})$  of **m** is  $\ell(\mathbf{m})$ .

Multi-patterns are denoted as matrices by stacking their patterns.

## – Example –

Let the multi-pattern

$$\mathbf{m} := \begin{bmatrix} 0 & 4 & \square & 4 & 0 & 0 & \square & \square \\ \overline{7} & \overline{7} & 0 & \square & \overline{3} & \overline{3} & \square & \square & \square \end{bmatrix}.$$

The arity of  ${\bf m}$  is 5, its length is 9, and its multiplicity is 2.

By interpreting each pattern of a multi-pattern through an interpretation, a multi-pattern specifies a musical phrase consisting in stacked voices.

## - Example -

The previous multi-pattern, interpreted through the degree interpretation specifying the  $9_3$  harmonic minor scale and the rhythm interpretation specifying 128 as tempo gives the musical phrase



# Operads of degree patterns

A **degree monoid** is a monoid  $(D, \star, e)$  such that  $D \subseteq \mathbb{Z}$ .

The *D*-degree pattern operad is the operad  $\mathbf{DP}^D := \mathbf{T}D$ . The elements of  $\mathbf{DP}^D$  are degree patterns on *D*.

## - Examples -

By denoting by  $\mathbb{Z}$  the monoid  $(\mathbb{Z}, +, 0)$ , we have in  $\mathbf{DP}^{\mathbb{Z}}$ ,

$$012\overline{1} \circ_3 024 = 01\ 246\ \overline{1},$$

$$012\overline{1} \odot 024 = 024 \ 135 \ 246 \ \overline{1}13.$$

By denoting, for any  $k \ge 1$ , by  $\mathbb{C}_k$  the cyclic monoid  $(\mathbb{Z}/k\mathbb{Z}, +, 0)$ , we have in  $\mathbf{DP}^{\mathbb{C}_3}$ ,

$$20010 \circ_4 2120 = 20002010.$$

By denoting, for any subset Z of  $\mathbb{Z}$  having a lower bound z, by  $\mathbb{M}_Z$  the monoid  $(Z, \max, z)$ , we have in  $\mathbf{DP}^{\mathbb{M}[0,2]}$ ,

$$20010 \circ_4 2120 = 20021210.$$

# **Operad of rhythm patterns**

Let  $\mathbb{N}$  be the monoid  $(\mathbb{N}, +, 0)$ .

The **rhythm pattern operad** is the operad  $RP := U\mathbb{N}$ . The elements of RP are duration sequences.

## - Example -

In RP,

$$001_21 \circ_3 110 = 00 212 1.$$

Since duration sequences and rhythm patterns are in one-to-one correspondence, **RP** can be seen as an operad on rhythm patterns.

On rhythm patterns, the partial composition of **RP** expresses as follows: if **r** and **r**' are two rhythm patterns, then  $\mathbf{r} \circ_i \mathbf{r}'$  is obtained by replacing the *i*-th beat of **r** by  $\mathbf{r}'$ .

## - Example -

The previous composition, seen on rhythm patterns, translates as



# Operads of patterns

The **Hadamard product** of two operads  $\mathcal{O}$  and  $\mathcal{O}'$  is the operad  $\mathcal{O} \boxtimes \mathcal{O}'$  such that for any  $n \geqslant 1$ ,

$$(\mathcal{O} \boxtimes \mathcal{O}')(n) := \mathcal{O}(n) \times \mathcal{O}'(n)$$

and for any  $(x, x'), (y, y') \in \mathcal{O} \boxtimes \mathcal{O}'$ ,

$$(x,x')\circ_i(y,y')=(x\circ_i y,x'\circ_i y').$$

Let *D* be a degree monoid.

The *D*-pattern operad is the operad  $\mathbf{P}^D := \mathbf{DP}^D \boxtimes \mathbf{RP}$ . The elements of  $\mathbf{P}^D$  are pairs  $(\mathbf{d}, \mathbf{r})$  such that  $|\mathbf{d}| = |\mathbf{r}|$ . Therefore, the elements of  $\mathbf{P}^D$  are patterns.

## - Examples -

In  $\mathbf{P}^{\mathbb{Z}}$ , we have

$$(\overline{2}31, \square\square\square\square\square) \circ_2 (0\overline{1}, \square\square\square) = (\overline{2}\ 32\ 1, \square\square\square\square\square\square\square)$$

which translates through the concise notation for patterns as

$$\square \bar{2}3 \square 1 \circ_2 0 \square \bar{1} = \square \bar{2}3 \square 2 \square 1.$$

# **Operads of multi-patterns**

Let *D* be a degree monoid and  $m \ge 1$ . Let  $\mathbf{P}_m^{D'}$  be the operad defined as

$$\mathbf{P}_{m}^{D'} := \underbrace{\mathbf{P}^{D} \boxtimes \cdots \boxtimes \mathbf{P}^{D}}_{m \text{ terms}}.$$

Let also  $\mathbf{P}_m^D$  be the subset of the underlying graded set of  $\mathbf{P}_m^{D'}$  restrained to the sequences  $\mathbf{m}(1) \dots \mathbf{m}(m)$  such that  $\ell(\mathbf{m}(1)) = \dots = \ell(\mathbf{m}(m))$ .

For any degree monoid D and any positive integer m,  $\mathbf{P}_m^D$  is an operad.

We call  $\mathbf{P}_m^D$  the *D*-music box operad.

## - Example -

In 
$$\mathbf{P}_2^{\mathbb{Z}}$$
,

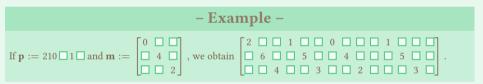
$$\begin{bmatrix} \square & \bar{2} & \bar{1} & \square & 0 \\ 0 & 1 & \square & \square & 1 \end{bmatrix} \circ_2 \begin{bmatrix} 1 & \square \\ \bar{3} & \square \end{bmatrix} = \begin{bmatrix} \square & \bar{2} & 0 & \square & 0 \\ 0 & \bar{2} & \square & \square & 1 \end{bmatrix}.$$

# Operations on musical phrases

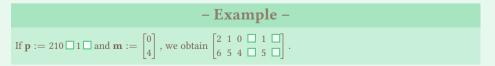
Every element of  $\mathbf{P}_m^D(n)$  can be seen as an operator of arity n acting on multi-patterns.

Here are some examples.

■ If **m** is an arpeggio shape and **p** is a pattern,  $\mathbf{p}' \odot \mathbf{m}$  is an arpeggiation of **p**, where  $\mathbf{p}'$  is the multi-pattern obtained by stacking  $\mathfrak{m}(\mathbf{m})$  equal voices from **p**.



■ If **m** is a chord shape and **p** is a pattern,  $\mathbf{p}' \odot \mathbf{m}$  is an harmonization of **p**, where  $\mathbf{p}'$  is the multi-pattern obtained by stacking  $\mathfrak{m}(\mathbf{m})$  equal voices from **p**.



## **Outline**

3. Bud generating systems

# **Colored operads**

Let  $\mathfrak C$  be a nonempty set, called **set of colors**.

A  $\mathfrak{C}$ -colored operad is a triple  $(\mathcal{C}, \circ, \mathbb{1})$  where

 $\mathbf{C}$  is a set

$$C = \bigsqcup_{\substack{a \in \mathfrak{C} \\ u \in \mathfrak{C}^+}} C(a, u);$$

- $\circ: \mathcal{C}(a, u) \times (\mathcal{C}(u(1), v_1) \times \cdots \times \mathcal{C}(u(\ell(u)), v_{\ell(u)})) \to \mathcal{C}(a, v_1 \dots v_{\ell(u)})$  is a map called **colored full composition**;
- $1 : \mathfrak{C} \to \mathcal{C}(a, a)$  is a map called **colored unit**.

This data satisfies similar relations than the ones of operads.

Intuitively, in a colored operad, each element has an output color and a color for each input. The composition is defined only when colors match.

# **Bud operads**

Let  $(\mathcal{O}, \circ, \mathbb{1})$  be an operad and  $\mathfrak{C}$  be a set of colors.

Let  $(B_{\mathfrak{C}}\mathcal{O}, \circ', \mathbb{1}')$  be the triple such that

• for any  $a \in \mathfrak{C}$  and  $u \in \mathfrak{C}^+$ ,

$$(\mathsf{B}_{\mathfrak{C}}\mathcal{O})(a,u) := \{(a,x,u) : x \in \mathcal{O}(\ell(u))\};$$

• for any  $(a, x, u) \in B_{\mathfrak{C}}\mathcal{O}$  and  $(u(i), y_i, v_i) \in B_{\mathfrak{C}}\mathcal{O}$ ,  $i \in [\ell(u)]$ ,

$$(a, x, u) \circ' (u(1), y_1, v_1) \dots (u(\ell(u)), y_{\ell(u)}, v_{\ell(u)}) = (a, x \circ y_1 \dots y_{\ell(u)}, v_1 \dots v_{\ell(u)});$$

• for any  $c \in \mathfrak{C}$ , 1/(c) := (c, 1, c).

#### - **Proposition** [G., 2019] -

For any operad  $\mathcal{O}$  and any set of colors  $\mathfrak{C}$ ,  $B_{\mathfrak{C}}\mathcal{O}$  is a  $\mathfrak{C}$ -colored operad.

# **Bud generating systems**

A bud generating system is a tuple  $\mathcal{B}:=(\mathcal{O},\mathfrak{C},\mathcal{R},b)$  where

- $(\mathcal{O}, \circ_i, \mathbb{1})$  is an operad, called **ground operad**;
- C is a finite set of colors;
- $\mathcal{R}$  is a finite subset of  $B_{\mathfrak{C}}(\mathcal{O})$ , called **set of rules**;
- **b** is a color of C, called **initial color**.

Let  $\stackrel{\circ}{\to}$  be the binary relation on  $B_{\mathfrak{C}}\mathcal{O}$  such that  $(a, x, u) \stackrel{\circ}{\to} (a, y, v)$  if there are rules  $r_1, \ldots, r_{|x|} \in \mathcal{R}$  such that

$$(a, y, v) = (a, x, u) \circ r_1 \dots r_{|x|}.$$

An element x of  $\mathcal{O}$  is **fully generated** by  $\mathcal{B}$  if there is an element  $(\mathbf{b}, x, u)$  such that  $(\mathbf{b}, 1, \mathbf{b})$  is in relation with  $(\mathbf{b}, x, u)$  w.r.t. the reflexive and transitive closure of  $\stackrel{\circ}{\rightarrow}$ .

# **Bud** generating systems

## - Example -

Let the bud generating system  $\mathcal{B}:=\left(P_2^\mathbb{Z},\{b_1,b_2,b_3\},\{c_1,c_2,c_3,c_4,c_5\},b_1\right)$  where

$$\begin{split} \mathbf{c}_1 := \left(\mathbf{b}_1, \begin{bmatrix} 0 & 2 & \square & 1 & \square & 0 & 4 \\ \overline{5} & \square & 0 & 0 & 0 & 0 \end{bmatrix}, \mathbf{b}_3 \mathbf{b}_2 \mathbf{b}_1 \mathbf{b}_1 \mathbf{b}_3 \right), \quad \mathbf{c}_2 := \left(\mathbf{b}_1, \begin{bmatrix} 1 & \square & 0 \\ 0 & \square & 1 \end{bmatrix}, \mathbf{b}_1 \mathbf{b}_1 \right), \\ \mathbf{c}_3 := \left(\mathbf{b}_2, \begin{bmatrix} \overline{1} \\ \overline{1} \end{bmatrix}, \mathbf{b}_1 \right), \quad \mathbf{c}_4 := \left(\mathbf{b}_2, \begin{bmatrix} 0 & 0 \\ 0 & 0 \end{bmatrix}, \mathbf{b}_1 \mathbf{b}_1 \right), \quad \mathbf{c}_5 := \left(\mathbf{b}_3, \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \mathbf{b}_3 \right). \end{split}$$

Since

$$\begin{pmatrix} \mathbf{b}_1, \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \mathbf{b}_1 \end{pmatrix} \xrightarrow{\circ} \begin{pmatrix} \mathbf{b}_1, \begin{bmatrix} 0 & 2 & 1 & 1 & 0 & 4 \\ \bar{5} & 1 & 0 & 0 & 0 & 0 \end{bmatrix}, \mathbf{b}_3 \mathbf{b}_2 \mathbf{b}_1 \mathbf{b}_1 \mathbf{b}_3 \end{pmatrix}$$

$$\xrightarrow{\circ} \begin{pmatrix} \mathbf{b}_1, \begin{bmatrix} 0 & 1 & 1 & 2 & 1 & 1 & 0 & 2 & 1 & 0 & 0 & 4 & 4 \\ \bar{5} & 1 & 1 & 0 & 1 & \bar{5} & 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}, \mathbf{b}_3 \mathbf{b}_1 \mathbf{b}_1 \mathbf{b}_3 \mathbf{b}_2 \mathbf{b}_1 \mathbf{b}_1 \mathbf{b}_3 \mathbf{b}_3 \end{pmatrix},$$

the multi-pattern

$$\begin{bmatrix} 0 & 1 & \square & 2 & \square & 1 & \square & 0 & 2 & \square & 1 & \square & 0 & 4 & 4 \\ \overline{5} & \square & \overline{1} & 0 & \square & 1 & \overline{5} & \square & \square & 0 & 0 & 0 & 0 \end{bmatrix}$$

is fully generated by  $\mathcal{B}$ .

# A random generation algorithm

For any color  $a \in \mathfrak{C}$ , we shall denote by  $\mathcal{R}_a$  the set of all rules of  $\mathcal{R}$  having a as output color.

If S is a nonempty finite set, random(S) returns an element of S picked uniformly at random.

Let us consider the following random generation algorithm.

**Data:** A bud generating system  $\mathcal{B} := (\mathcal{O}, \mathfrak{C}, \mathcal{R}, \mathbf{b})$  and an integer  $k \in \mathbb{N}$ .

**Result:** A randomly generated element of  $\mathcal{O}$ .

```
1 begin
2 (a, x, u) \leftarrow (\mathbf{b}, 1, \mathbf{b})
3 \mathbf{for} j \in [k] \mathbf{do}
4 R \leftarrow \mathcal{R}_{u(1)} \times \cdots \times \mathcal{R}_{u(\ell(x))}
5 \mathbf{if} R \neq \emptyset \mathbf{then}
6 (a, x, u) \leftarrow (a, x, u) \circ \mathrm{random}(R)
7 \mathbf{return} x
```

# A random generation algorithm

### - Example -

Let the bud generating system  $\mathcal{B} := (\mathbf{P}_2^{\mathbb{Z}}, \{\mathbf{b}_1, \mathbf{b}_2, \mathbf{b}_3\}, \{\mathbf{c}_1, \mathbf{c}_2, \mathbf{c}_3, \mathbf{c}_4, \mathbf{c}_5\}, \mathbf{b}_1)$  where

$$\begin{aligned} \mathbf{c}_1 &:= \left(\mathbf{b}_1, \begin{bmatrix} 0 & 2 & \square & 1 & \square & 0 & 4 \\ \bar{\mathbf{5}} & \square & 0 & 0 & 0 & 0 \end{bmatrix}, \mathbf{b}_3 \mathbf{b}_2 \mathbf{b}_1 \mathbf{b}_1 \mathbf{b}_3 \right), \quad \mathbf{c}_2 := \left(\mathbf{b}_1, \begin{bmatrix} 1 & \square & 0 \\ 0 & \square & 1 \end{bmatrix}, \mathbf{b}_1 \mathbf{b}_1 \right), \\ \mathbf{c}_3 &:= \left(\mathbf{b}_2, \begin{bmatrix} \bar{1} \\ \bar{1} \end{bmatrix}, \mathbf{b}_1 \right), \quad \mathbf{c}_4 := \left(\mathbf{b}_2, \begin{bmatrix} 0 & 0 \\ 0 & 0 \end{bmatrix}, \mathbf{b}_1 \mathbf{b}_1 \right), \quad \mathbf{c}_5 := \left(\mathbf{b}_3, \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \mathbf{b}_3 \right). \end{aligned}$$

The previous random generation algorithm run with this bud generating system and k:=2 generates the multi-pattern

$$\begin{bmatrix} 2 & 4 & \square & 3 & \square & 2 & 6 & \square & 2 & \square & 1 & \square & 0 & 1 & \square & 2 & \square & 1 & \square & 0 & 4 \\ \overline{5} & \square & \square & 0 & 0 & 0 & \square & 1 & \square & 2 & \square & \overline{4} & \square & 0 & 0 & 1 & \square & 2 & 1 & \square & 2 & 1 \end{bmatrix}.$$

This pattern is obtained from the underlying random syntax tree



# The Bud Music Box program

An implementation of these concepts can be found at

https://github.com/SamueleGiraudo/Bud-Music-Box

#### Here is a Bud Music Box program:

```
scale 2 1 4 1 4
root 57
tempo 192
sounds 3 1
multi_pattern mpat_1 0 2 4 * : -5 * 0 -1
multi_pattern mpat_2 0 * ; * 0
multi_pattern mpat_3 0 ; 0
colorize cpat_1 mpat_1 c1 c3 c1 c3
colorize cpat_2 mpat_1 c1 c3 c2 c3
colorize cpat_3 mpat_2 c2 c1
colorize cpat_4 mpat_3 c3 c3
generate mpat_3 full 8 c1 cpat_1 cpat_2 cpat_3 cpat_4
show
play mpat_3
```

Here is a randomly generated pattern:

```
0 2 4 6 8 10 12 * 14 * 12 * 10 * 8 * * 6 * 4 * ;
-5 * -5 * * -5 * -5 * -5 * -5 * * 0 -1 -1 -1 -1 -1 -1
```

■ Interpreted in the Hirajoshi scale.

Here is another one:

```
0 2 4 6 8 10 12 * * 10 * * 8 * 6 * * 4 * ;
-5 * * -5 * -5 * * -5 * * -5 * 0 -1 -1 -1 -1 -1
```

Interpreted in the Hirajoshi scale.