

# Combinatorial realizations of algebraic structures

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Journée Opérades LAGA-LIPN 2020  
Université Paris 13

November 5, 2020

# Outline

1. Universal algebra
2. Operads and varieties from monoids
3. Clones and realizations of semigroups

# Outline

## 1. Universal algebra

# Types of algebraic structures

Combinatorics deals with sets (or spaces) of structured objects:

- monoids;
- groups;
- lattices;
- associative alg.;
- Hopf bialg.;
- Lie alg.;
- pre-Lie alg.;
- dendriform alg.;
- duplicial alg.

Such **types of algebras** are specified by

1. a collection of operations;
2. a collection of relations between operations.

## – Example –

The type of monoids can be specified by

1. the operations  $\star$  (binary) and  $\mathbb{1}$  (nullary);
2. the relations  $(x_1 \star x_2) \star x_3 = x_1 \star (x_2 \star x_3)$  and  $x \star \mathbb{1} = x = \mathbb{1} \star x$ .

# Universal algebra

**Universal algebra** is a formalism to work with such structures.

A **signature** is a graded set  $\mathfrak{G} := \bigsqcup_{k \geq 0} \mathfrak{G}(k)$  wherein each  $\mathbf{a} \in \mathfrak{G}(k)$  is an **operation** of arity  $k$ .

A  **$\mathfrak{G}$ -term** is

- either a variable  $x$  from the set  $\mathbb{X} := \{x_1, x_2, \dots\}$ ;
- either a pair  $(\mathbf{a}, (t_1, \dots, t_k))$  where  $\mathbf{a} \in \mathfrak{G}(k)$  and each  $t_i$  is a  $\mathfrak{G}$ -term.

The set of all  $\mathfrak{G}$ -terms is denoted by  $\mathfrak{T}(\mathfrak{G})$ .

## – Example –



This is the tree representation of the  $\mathfrak{G}$ -term

$$(\times, ((+, (x_1, x_2)), (+, ((\times, (x_1, x_1)), x_3))))$$

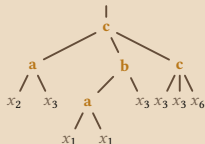
where  $\mathfrak{G} := \mathfrak{G}(2) := \{+, \times\}$ .

## More on terms

Let  $t$  be a  $\mathcal{G}$ -term.

The **frontier** of  $t$  is the sequence of all variables appearing in  $t$ .

### – Example –



The frontier of this term is  $(x_2, x_3, x_1, x_1, x_3, x_3, x_3, x_6)$   
and its ground arity is 6.

The **ground arity** of  $t$  is the greatest integer  $n$  such that  $x_n$  is a variable appearing in  $t$ .

The term  $t$  is

- **planar** if its frontier is  $(x_1, \dots, x_n)$ ;
- **standard** if its frontier is a permutation of  $(x_1, \dots, x_n)$ ;
- **linear** if there are no multiple occurrences of the same variable in the frontier of  $t$ .

# Varieties

A  $\mathfrak{G}$ -**equation** is a pair  $(t, t')$  where  $t$  and  $t'$  are both  $\mathfrak{G}$ -terms.

A **variety** is a pair  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G}$  is a signature and  $\mathfrak{R}$  is a set of  $\mathfrak{G}$ -equations. We denote by  $t \mathfrak{R} t'$  the fact that  $(t, t') \in \mathfrak{R}$ .

## – Example –

The **variety of groups** is the pair  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(0) \sqcup \mathfrak{G}(1) \sqcup \mathfrak{G}(2)$  with  $\mathfrak{G}(0) := \{1\}$ ,  $\mathfrak{G}(1) := \{i\}$ , and  $\mathfrak{G}(2) := \{\star\}$ , and  $\mathfrak{R}$  is the set of  $\mathfrak{G}$ -equations satisfying

$$\begin{array}{c} | \\ \star \\ / \quad \backslash \\ \star \quad x_3 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} | \\ \star \\ / \quad \backslash \\ x_1 \quad \star \\ / \quad \backslash \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \star \\ / \quad \backslash \\ x_1 \quad 1 \end{array} \mathfrak{R} \begin{array}{c} | \\ i \\ / \quad \backslash \\ 1 \quad \star \\ / \quad \backslash \\ 1 \quad x_1 \end{array}, \quad \begin{array}{c} | \\ \star \\ / \quad \backslash \\ i \quad x_1 \\ / \quad \backslash \\ x_1 \quad 1 \end{array} \mathfrak{R} \begin{array}{c} | \\ \star \\ / \quad \backslash \\ x_1 \quad i \\ / \quad \backslash \\ x_1 \quad 1 \end{array}.$$

## – Example –

The **variety of semilattices** is the pair  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\wedge\}$ , and  $\mathfrak{R}$  is the set of  $\mathfrak{G}$ -equations satisfying

$$\begin{array}{c} | \\ \wedge \\ / \quad \backslash \\ \wedge \quad x_3 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} | \\ \wedge \\ / \quad \backslash \\ x_1 \quad \wedge \\ / \quad \backslash \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \wedge \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} | \\ \wedge \\ / \quad \backslash \\ x_2 \quad x_1 \end{array}, \quad \begin{array}{c} | \\ \wedge \\ / \quad \backslash \\ x_1 \quad x_1 \end{array} \mathfrak{R} \begin{array}{c} | \\ i \\ / \quad \backslash \\ x_1 \quad 1 \end{array}.$$

# Algebras of a variety

Let  $\mathcal{A}$  be a nonempty set. An  $\mathcal{A}$ -**substitution** is a map  $\sigma : \mathbb{X} \rightarrow \mathcal{A}$ .

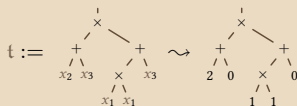
An  $\mathcal{A}$ -**interpretation** of a signature  $\mathfrak{G}$  is a set

$$\mathfrak{G}_{\mathcal{A}} := \left\{ \mathbf{a}_{\mathcal{A}} : \mathcal{A}^k \rightarrow \mathcal{A} : \mathbf{a} \in \mathfrak{G}(k) \text{ for a } k \geq 0 \right\}.$$

The **evaluation** of a  $\mathfrak{G}$ -term  $t$  under an  $\mathcal{A}$ -substitution  $\sigma$  and an  $\mathcal{A}$ -interpretation  $\mathfrak{G}_{\mathcal{A}}$  is defined by induction as

$$\text{ev}_{\mathcal{A}}^{\sigma}(t) := \begin{cases} \sigma(x) & \text{if } t = x \text{ is a variable,} \\ \mathbf{a}_{\mathcal{A}}(\text{ev}_{\mathcal{A}}^{\sigma}(t_1), \dots, \text{ev}_{\mathcal{A}}^{\sigma}(t_k)) & \text{otherwise, where } t = (\mathbf{a}, (t_1, \dots, t_k)). \end{cases}$$

## – Example –



With  $\mathcal{A} := \mathbb{N}$ ,  $\mathfrak{G}_{\mathcal{A}}$  defined naturally, and  $\sigma$  satisfying  $\sigma(x_1) := 1$ ,  $\sigma(x_2) := 2$ , and  $\sigma(x_3) := 0$ , one obtains  $\text{ev}_{\mathcal{A}}^{\sigma}(t) = 2$ .

An **algebra** of a variety  $(\mathfrak{G}, \mathfrak{R})$  is a pair  $(\mathcal{A}, \mathfrak{G}_{\mathcal{A}})$  where for any  $(t, t') \in \mathfrak{R}$  and  $\mathcal{A}$ -substitution  $\sigma$ ,  $\text{ev}_{\mathcal{A}}^{\sigma}(t) = \text{ev}_{\mathcal{A}}^{\sigma}(t')$ .



# Equivalent terms

Two  $\mathfrak{G}$ -terms  $t$  and  $t'$  are  $\mathfrak{R}$ -equivalent if for all algebras  $(\mathcal{A}, \mathfrak{G}_{\mathcal{A}})$  of  $(\mathfrak{G}, \mathfrak{R})$  and for all  $\mathcal{A}$ -substitutions  $\sigma$ , one has  $\text{ev}_{\mathcal{A}}^{\sigma}(t) = \text{ev}_{\mathcal{A}}^{\sigma}(t')$ . This property is denoted by  $t \equiv_{\mathfrak{R}} t'$ .

## – Example –

In the variety of groups,

$$\begin{array}{c} | \\ i \\ | \\ \star \\ / \backslash \\ x_1 \quad x_2 \end{array} \equiv_{\mathfrak{R}} \begin{array}{c} | \\ \star \\ / \backslash \\ i \quad i \\ | \quad | \\ x_2 \quad x_1 \end{array} .$$

## – Questions –

1. Design an algorithm to decide if two  $\mathfrak{G}$ -terms are  $\equiv_{\mathfrak{R}}$ -equivalent. This is known as the **word problem**.
2. Construct a system of representatives  $\mathcal{C}$  of the  $\equiv_{\mathfrak{R}}$ -equivalence classes. The set  $\mathcal{C}$  is a **combinatorial realization** of the variety.
3. Enumerate the  $\equiv_{\mathfrak{R}}$ -equivalence classes of (planar/standard/linear)  $\mathfrak{G}$ -terms w.r.t. their ground arity.

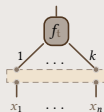
# Abstract operations and compositions

To tackle these issues, we need a formalization and an abstraction of the notion of **composition of terms** in order to consider operations over operations.

An  $\mathcal{G}$ -term  $t$  on the variables  $\{x_1, \dots, x_n\}$  is an **abstract operation**

$$(x_1, \dots, x_n) \mapsto f_t(x_1, \dots, x_n)$$

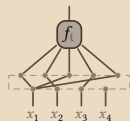
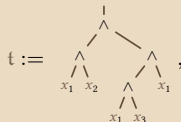
depicted as



where  $k$  is the length of the frontier of  $t$ .

## – Example –

For the signature  $\mathcal{G}$  of the variety of semilattices, here is a  $\mathcal{G}$ -term seen on the set  $\{x_1, \dots, x_4\}$  of variables and the abstract operation it denotes:

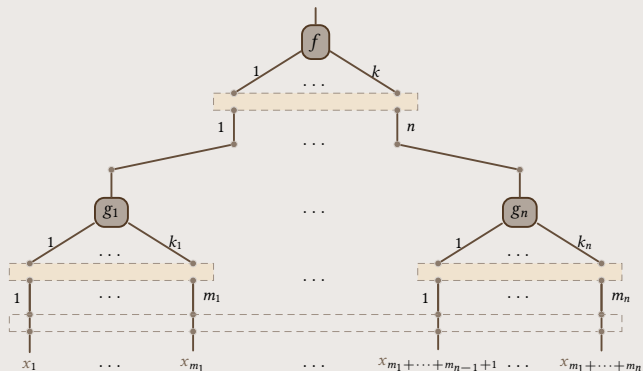


# A first paradigm for composition

If  $f$  is an abstract operation of arity  $n$  and  $g_1, \dots, g_n$  are abstract operations of respective arities  $m_1, \dots, m_n$ , then  $f \circ [g_1, \dots, g_n]$  is the abstract operation satisfying

$$(x_1, \dots, x_{m_1+\dots+m_n}) \mapsto f(g_1(x_1, \dots, x_{m_1}), \dots, g_n(x_{m_1+\dots+m_{n-1}+1}, \dots, x_{m_1+\dots+m_n})).$$

This is the abstract operation depicted as

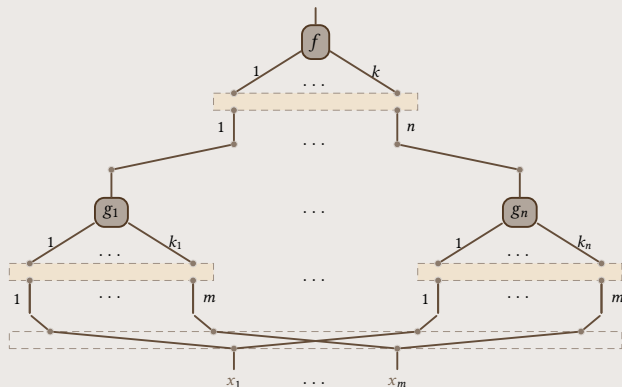


## A second paradigm for composition

If  $f$  is an abstract operation of arity  $n$  and  $g_1, \dots, g_n$  are abstract operations all of arity  $m$ , then  $f \odot [g_1, \dots, g_n]$  is the operation satisfying

$$(x_1, \dots, x_m) \mapsto f(g_1(x_1, \dots, x_m), \dots, g_n(x_1, \dots, x_m)).$$

This is the abstract operation depicted as



# Duplicial algebras

A **duplicial algebra** [Brouder, Frabetti, 2003] is a set  $\mathcal{A}$  endowed with two binary operations

$$\ll, \gg: \mathcal{A}^2 \rightarrow \mathcal{A}$$

satisfying the three relations

$$(x_1 \ll x_2) \ll x_3 = x_1 \ll (x_2 \ll x_3),$$

$$(x_1 \gg x_2) \ll x_3 = x_1 \gg (x_2 \ll x_3),$$

$$(x_1 \gg x_2) \gg x_3 = x_1 \gg (x_2 \gg x_3).$$

## – Example –

On  $\mathbb{N}^*$ , let  $\ll$  and  $\gg$  be the operations defined by

$$u \ll v := u(v \uparrow_{\max(u)}), \quad u \gg v := u(v \uparrow_{|u|}).$$

Then, for instance,

$$0211 \ll 14 = 021136, \quad 0211 \gg 14 = 021158.$$

This structure is a duplicial algebra [Novelli, Thibon, 2013].

# Duplicial operations and equivalence

Let us describe a way to test if two planar duplicial operations are equivalent.

By the duplicial relations, we have

$$\begin{array}{c} | \\ \ll \\ \swarrow \quad \searrow \\ \ll \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \equiv \begin{array}{c} | \\ \ll \\ \swarrow \quad \searrow \\ x_1 \quad \ll \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ \ll \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \equiv \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ x_1 \quad \ll \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ \gg \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \equiv \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ x_1 \quad \gg \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}.$$

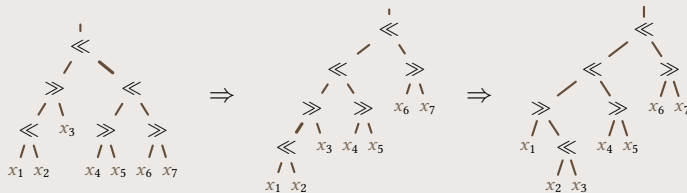
We orient these equations as

$$\begin{array}{c} | \\ \ll \\ \swarrow \quad \searrow \\ \ll \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \leftarrow \begin{array}{c} | \\ \ll \\ \swarrow \quad \searrow \\ x_1 \quad \ll \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ \ll \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \rightarrow \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ x_1 \quad \ll \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ \gg \quad x_3 \\ \swarrow \quad \searrow \\ x_1 \quad x_2 \end{array} \leftarrow \begin{array}{c} | \\ \gg \\ \swarrow \quad \searrow \\ x_1 \quad \gg \\ \swarrow \quad \searrow \\ x_2 \quad x_3 \end{array}.$$

in order to obtain a **rewrite relation**  $\Rightarrow$  on the set of all the duplicial operations by performing local moves.

# Testing equivalence of duplicial operations

We have for instance the sequence



of rewritings.

## – Proposition –

Two planar duplicial operations  $t$  and  $t'$  are equivalent iff there is a duplicial operation  $s$  such that  $t \Rightarrow^* s$  and  $t' \Rightarrow^* s$ .

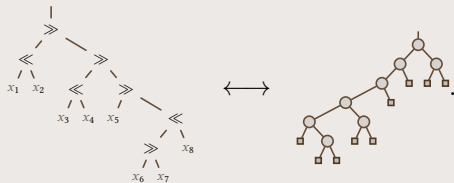
To prove this, we have to establish the fact that  $\Rightarrow$  is a **terminating and confluent** rewrite relation.

# Enumerating duplicial operations

## – Proposition –

The set of normal forms for  $\Rightarrow$  of planar duplicial operations with  $n \geq 0$  inputs is in one-to-one correspondence with the set of all binary trees with  $n$  internal nodes.

A possible bijection puts the following two trees in correspondence:



Therefore, there are

$$\frac{1}{n+1} \binom{2n}{n}$$

pairwise nonequivalent planar duplicial operations with  $n$  inputs.



# Distributive lattices

A **distributive lattice** is a set  $\mathcal{A}$  endowed with two binary operations

$$\wedge, \vee : \mathcal{A}^2 \rightarrow \mathcal{A}$$

satisfying the relations

$$\begin{aligned}(x_1 \wedge x_2) \wedge x_3 &= x_1 \wedge (x_2 \wedge x_3), & (x_1 \vee x_2) \vee x_3 &= x_1 \vee (x_2 \vee x_3), \\ x_1 \wedge x_2 &= x_2 \wedge x_1, & x_1 \vee x_2 &= x_2 \vee x_1, \\ x_1 \wedge (x_1 \vee x_2) &= x_1, & x_1 \vee (x_1 \wedge x_2) &= x_1, \\ x_1 \vee (x_2 \wedge x_3) &= (x_1 \vee x_2) \wedge (x_1 \vee x_3), & x_1 \wedge (x_2 \vee x_3) &= (x_1 \wedge x_2) \vee (x_1 \wedge x_3).\end{aligned}$$

## – Example –

- On  $[n]$ ,  $\vee$  defined as the union and  $\wedge$  as the intersection is a finite distributive lattice.
- The set of all Young diagrams is an infinite lattice for the intersection and the union of diagrams.

# Combinatorial realization

A **normal term** is a term  $t$  expressing as

$$t = s_1 \vee \dots \vee s_m, \quad m \geq 0, \quad \text{where} \quad s_i = x_{f_{i,1}} \wedge \dots \wedge x_{f_{i,k_i}}, \quad k_i \geq 1,$$

for any  $i, i' \in [k]$ ,  $x_{f_{i,r}} = x_{f_{i',r'}}$  implies  $r = r'$ , and  $\{f_{i,1}, \dots, f_{i,k_i}\} \subseteq \{f_{i',1}, \dots, f_{i',k_{i'}}\}$  implies  $i = i'$ .

## – Examples –

$(x_2 \wedge x_3 \wedge x_5) \vee (x_3 \wedge x_7) \vee (x_3 \wedge x_4) \vee x_6$  is a normal term.

$(x_2 \wedge x_3 \wedge x_5) \vee (x_2 \wedge x_5)$  is not.

## – Proposition –

The set of all sets of sets of positive integers  $\{\{f_{1,1}, \dots, f_{1,k_1}\}, \dots, \{f_{m,1}, \dots, f_{m,k_m}\}\}$  satisfying the above properties is a combinatorial realization of the variety of distributive lattices.

Pairwise nonequivalent distributive lattice operations with  $n$  inputs are enumerated by the **Dedekind numbers** whose sequence begins with (only these few terms are known today)

1, 2, 5, 19, 167, 7580, 7828353, 2414682040997, 56130437228687557907787.

## 2. Operads and varieties from monoids

# Operads

Nonsymmetric operads provide a formalization of **planar operations** under the **first paradigm for composition**.

A **nonsymmetric operad** is a triple  $(\mathcal{O}, \circ, \mathbb{1})$  where

- $\mathcal{O}$  is a graded set

$$\mathcal{O} = \bigsqcup_{n \geq 0} \mathcal{O}(n);$$

- $\circ$  is a map

$$\circ : \mathcal{O}(n) \times \mathcal{O}(m_1) \times \cdots \times \mathcal{O}(m_n) \rightarrow \mathcal{O}(m_1 + \cdots + m_n)$$

called **full composition map**;

- $\mathbb{1}$  is an element of  $\mathcal{O}(1)$  called **unit**.

This data has to satisfy some axioms.

# Operad axioms and partial composition maps

The following relations have to be satisfied:

- For all  $x \in \mathcal{O}$ ,

$$\mathbb{1} \circ [x] = x = x \circ [\mathbb{1}, \dots, \mathbb{1}].$$

This says that  $\mathbb{1}$  is the identity operation.

- For all  $x \in \mathcal{O}(n)$ ,  $y_i \in \mathcal{O}(m_i)$ , and  $z_{i,j} \in \mathcal{O}$ ,

$$\begin{aligned} (x \circ [y_1, \dots, y_n]) \circ [z_{1,1}, \dots, z_{1,m_1}, \dots, z_{n,1}, \dots, z_{n,m_n}] \\ = x \circ [y_1 \circ [z_{1,1}, \dots, z_{1,m_1}], \dots, y_n \circ [z_{n,1}, \dots, z_{n,m_n}]]. \end{aligned}$$

This says that the two ways to compose elements to form an operation having three layers (by starting from top or by starting from bottom) give the same operation.

The **partial composition map** of  $\mathcal{O}$  is the map  $\circ_i : \mathcal{O}(n) \times \mathcal{O}(m) \rightarrow \mathcal{O}(n + m - 1)$  where  $i \in [n]$  and defined by

$$x \circ_i y := x \circ [\overbrace{\mathbb{1}, \dots, \mathbb{1}}^{i-1}, y, \overbrace{\mathbb{1}, \dots, \mathbb{1}}^{n-i}].$$

# Free operads

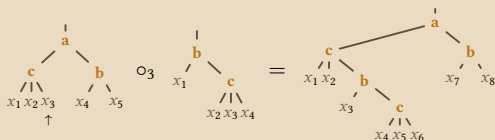
Let  $\mathfrak{G}$  be a signature.

The **free operad** on  $\mathfrak{G}$  is the operad  $(\mathfrak{P}(\mathfrak{G}), \circ, \mathbb{1})$  where

- $\mathfrak{P}(\mathfrak{G})$  is the set of all planar  $\mathfrak{G}$ -terms graded by the ground arity;
- $\circ_i$  is defined as follows. The  $\mathfrak{G}$ -term  $\mathfrak{t} \circ_i \mathfrak{s}$  is obtained by replacing the variable  $x_i$  of  $\mathfrak{t}$  by the root of  $\mathfrak{s}$ , and by setting  $(x_1, x_2, \dots)$  for the frontier of the obtained term;
- $\mathbb{1}$  is the  $\mathfrak{G}$ -term  $\overset{\textstyle |}{x_1}$ .

## – Example –

By setting  $\mathfrak{G} := \mathfrak{G}(2) \sqcup \mathfrak{G}(3)$  where  $\mathfrak{G}(2) := \{\mathbf{a}, \mathbf{b}\}$  and  $\mathfrak{G}(3) := \{\mathbf{c}\}$ , one has



in the free operad  $\mathfrak{P}(\mathfrak{G})$ .

# A variety from a monoid

Let  $(\mathcal{M}, \cdot, \epsilon)$  be a monoid.

Let the signature  $\mathfrak{G}_{\mathcal{M}} := \mathfrak{G}_{\mathcal{M}}(1) \sqcup \mathfrak{G}_{\mathcal{M}}(2)$  where  $\mathfrak{G}_{\mathcal{M}}(1) := \mathcal{M}$  and  $\mathfrak{G}_{\mathcal{M}}(2) := \{\mathbf{a}\}$ , and let  $\mathfrak{R}_{\mathcal{M}}$  be the set of  $\mathfrak{G}_{\mathcal{M}}$ -equations satisfying

$$\begin{array}{c} \mathbf{a} \\ \swarrow \searrow \\ \mathbf{a} \quad x_3 \\ \swarrow \searrow \\ x_1 \quad x_2 \end{array} \mathfrak{R}_{\mathcal{M}} \begin{array}{c} \mathbf{a} \\ \swarrow \searrow \\ x_1 \quad \mathbf{a} \\ \swarrow \searrow \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} \alpha_1 \\ | \\ \alpha_2 \\ | \\ x_1 \end{array} \mathfrak{R}_{\mathcal{M}} \begin{array}{c} \alpha_1 \cdot \alpha_2 \\ | \\ x_1 \end{array}, \quad \begin{array}{c} | \\ \epsilon \\ | \\ x_1 \end{array} \mathfrak{R}_{\mathcal{M}} \begin{array}{c} | \\ x_1 \end{array}, \quad \begin{array}{c} \mathbf{a} \\ \swarrow \searrow \\ \alpha \quad \alpha \\ | \quad | \\ x_1 \quad x_2 \end{array} \mathfrak{R}_{\mathcal{M}} \begin{array}{c} \alpha \\ | \\ \mathbf{a} \\ \swarrow \searrow \\ x_1 \quad x_2 \end{array}$$

for any  $\alpha, \alpha_1, \alpha_2 \in \mathcal{M}$ .

Any algebra of this variety is a semigroup  $(\mathcal{A}, \mathbf{a})$  endowed with semigroup endomorphisms  $\phi_{\alpha} : \mathcal{A} \rightarrow \mathcal{A}$  with  $\alpha \in \mathcal{M}$  and satisfying

$$\phi_{\epsilon}(x) = x,$$

$$\phi_{\alpha_1} \circ \phi_{\alpha_2} = \phi_{\alpha_1 \cdot \alpha_2}$$

for any  $\alpha_1, \alpha_2 \in \mathcal{M}$  and  $x \in \mathcal{M}$ .

# Orientation of the equations

Let the orientation  $\rightarrow$  of  $\mathfrak{R}_{\mathcal{M}}$  satisfying

$$\begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \rightarrow \begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ x_1 \quad \begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ x_2 \quad x_3 \end{array} \end{array}, \quad \begin{array}{c} | \\ \alpha_1 \\ | \\ \alpha_2 \\ | \\ x_1 \end{array} \rightarrow \begin{array}{c} | \\ \alpha_1 \cdot \alpha_2 \\ | \\ x_1 \end{array}, \quad \begin{array}{c} | \\ \epsilon \\ | \\ x_1 \end{array} \rightarrow \begin{array}{c} | \\ x_1 \end{array}, \quad \begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ \alpha \quad \alpha \\ | \quad | \\ x_1 \quad x_2 \end{array} \leftarrow \begin{array}{c} | \\ \alpha \\ | \\ \text{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array},$$

## – Proposition –

Two planar  $\mathfrak{G}_{\mathcal{M}}$ -terms  $t$  and  $t'$  are equivalent iff there is a  $\mathfrak{G}_{\mathcal{M}}$ -term  $s$  such that  $t \xRightarrow{*} s$  and  $t' \xRightarrow{*} s$ .

This is a consequence of the fact that  $\Rightarrow$  is a convergent rewrite relation.

The set of normal forms for  $\Rightarrow$  of planar  $\mathfrak{G}_{\mathcal{M}}$ -terms is the set of the terms of the form

$$\begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ s_1 \quad \begin{array}{c} | \\ \text{a} \\ / \quad \backslash \\ s_2 \quad \dots \quad \text{a} \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad s_{n-1} \quad s_n \end{array} \end{array} \quad \text{where} \quad s_i \in \left\{ \begin{array}{c} | \\ x_i \end{array}, \begin{array}{c} | \\ \alpha_i \\ | \\ x_i \end{array} \right\}, \quad \alpha_i \in \mathcal{M} \setminus \{\epsilon\}.$$



# Combinatorial realization

Let  $T\mathcal{M}$  be the set of all words on  $\mathcal{M}$ , graded by their length, let  $\circ_i$  be the partial composition map defined by

$$u \circ_i v := u(1) \dots u(i-1) (u(i) \bar{\cdot} v) u(i+1) \dots u(n),$$

where for any  $\alpha \in \mathcal{M}$  and  $w \in \mathcal{M}^*$ ,

$$\alpha \bar{\cdot} w := (\alpha \cdot w(1)) \dots (\alpha \cdot w(|w|)),$$

and let  $\mathbb{1}$  be  $\epsilon$  seen as a word of length 1.

## – Example –

In  $T(\mathbb{N}, +, 0)$ ,

$$2100213 \circ_5 3001 = 2100\ 5223\ 13.$$

## – Theorem [G., 2015] –

For any monoid  $\mathcal{M}$ , the triple  $(T\mathcal{M}, \circ, \mathbb{1})$  is an operad.

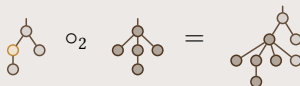
Moreover, this operad is a combinatorial realization of the variety  $(\mathfrak{G}_{\mathcal{M}}, \mathfrak{R}_{\mathcal{M}})$ .

## Some suboperads

The operads  $\mathcal{T}\mathcal{M}$  are large enough to contain a lot of suboperads realizable in combinatorial terms. As main examples:

- For any  $m \geq 0$ , with  $\mathcal{M} := (\mathbb{N}, +, 0)$ ,
  - **PRT** $_m$ , generated by  $\{01, \dots, 0m\}$ , on primitive  $m$ -Dyck paths;
  - **FCat** $_m$ , generated by  $\{00, 01, \dots, 0m\}$ , on  $m$ -trees;
  - **Schr** $_m$ , generated by  $\{01, \dots, 0m\} \cup \{00\} \cup \{10, \dots, m0\}$ , on some Schröder trees;
  - **Motz** $_m$ , generated by  $\{00, 000, 010, \dots, 0m0\}$ , on colored Motzkin paths.
- For any  $m \geq 0$ , with  $\mathcal{M} := (\mathbb{Z}/(m+1)\mathbb{Z}, +, 0)$ ,
  - **Comp** $_m$ , generated by  $\{00, 01, \dots, 0m\}$ , on  $m$ -words;
  - **DA** $_m$ , generated by  $\{00, 01, \dots, 0(m-1)\}$ , on some directed animals.
- For any  $m \geq 0$ ,  $\mathcal{M} := (\mathbb{N}, \max, 0)$ ,
  - **Dias** $_m$ , generated by  $\{01, \dots, 0m\} \cup \{10, \dots, m0\}$ , is the  $m$ -pluriassociative operad [Loday, 2001] [G., 2016];
  - **Trias** $_m$ , generated by  $\{01, \dots, 0m\} \cup \{00\} \cup \{10, \dots, m0\}$ , is the  $m$ -pluritriassociative operad [Loday, Ronco, 2004] [G., 2016].

# Some partial compositions on combinatorial objects



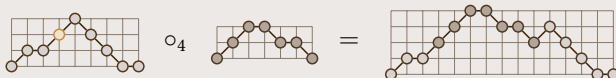
(in  $\mathbf{PRT}_1$ )



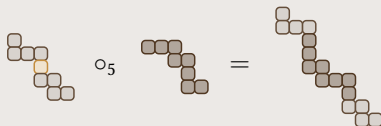
(in  $\mathbf{FCat}_2$ )



(in  $\mathbf{Schr}_1$ )

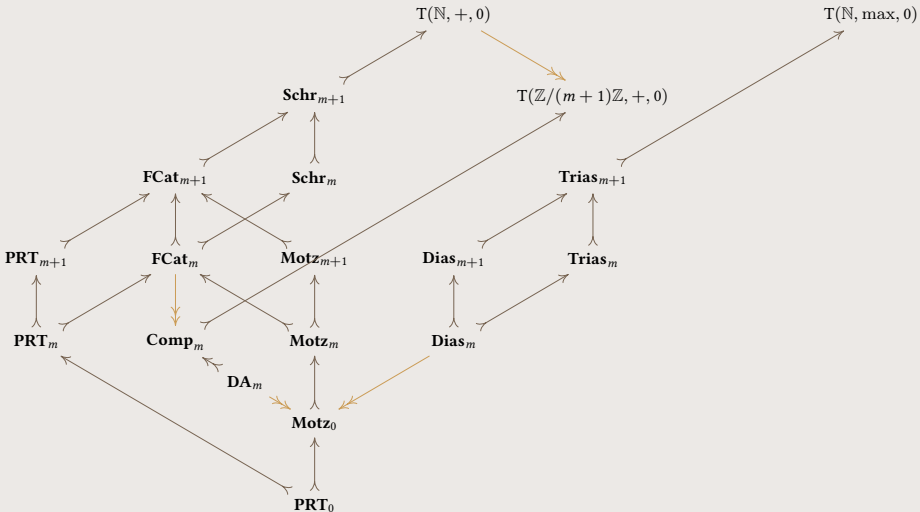


(in  $\mathbf{Motz}_1$ )



(in  $\mathbf{Comp}_1$ )

## Full diagram



## 3. Clones and realizations of semigroups

# Clones

Abstract clones provide a formalization of **general operations** under the **second paradigm for composition**.

An **abstract clone** is a triple  $(\mathcal{C}, \odot, \mathbb{1}_{i,n})$  where

- $\mathcal{C}$  is a graded set

$$\mathcal{C} = \bigsqcup_{n \geq 0} \mathcal{C}(n);$$

- $\odot$  is a map

$$\odot : \mathcal{C}(n) \times \mathcal{C}(m)^n \rightarrow \mathcal{C}(m)$$

called **superposition map**;

- for each  $n \geq 0$  and  $i \in [n]$ ,  $\mathbb{1}_{i,n}$  is an element of  $\mathcal{C}(n)$  called **projection**.

This data has to satisfy some axioms.

# Clone axioms

The following relations have to be satisfied:

- For all  $x_i \in \mathcal{C}(m)$ ,

$$\mathbb{1}_{i,n} \odot [x_1, \dots, x_n] = x_i.$$

This says that  $\mathbb{1}_{i,n}$  is the operation returning its  $i$ -th input.

- For all  $x \in \mathcal{C}(n)$ ,

$$x \odot [\mathbb{1}_{1,n}, \dots, \mathbb{1}_{n,n}] = x,$$

This says that each  $\mathbb{1}_{j,n}$ , put as  $j$ -th input, is an identity operation.

- For all  $x \in \mathcal{C}(n)$ ,  $y_i \in \mathcal{C}(m)$ , and  $z_j \in \mathcal{C}(k)$ ,

$$(x \odot [y_1, \dots, y_n]) \odot [z_1, \dots, z_m] = x \odot [y_1 \odot [z_1, \dots, z_m], \dots, y_n \odot [z_1, \dots, z_m]].$$

This says that the two ways to compose elements to form an operation having three layers (by starting from top or by starting from bottom) give the same operation.

# Free Clones

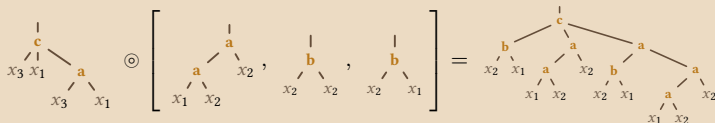
Let  $\mathcal{G}$  be a signature.

The **free clone** on  $\mathcal{G}$  is the clone  $(\mathcal{T}(\mathcal{G}), \odot, \mathbb{1}_{i,n})$  where

- $\mathcal{T}(\mathcal{G})$  is the set of all  $\mathcal{G}$ -terms. Each  $\mathcal{G}$ -term  $t$  is endowed with an integer equal as or greater than its ground arity and called **arity**;
- $\odot$  is defined as follows. The  $\mathcal{G}$ -term  $t \odot [s_1, \dots, s_n]$  is obtained by replacing each occurrence of a variable  $x_i$  of  $t$  by the root of  $s_i$  (without any relabeling);
- $\mathbb{1}_{i,n}$  is the term  $\overset{i}{x}_i$  of arity  $n$ .

## – Example –

By setting  $\mathcal{G} := \mathcal{G}(2) \sqcup \mathcal{G}(3)$  where  $\mathcal{G}(2) := \{a, b\}$  and  $\mathcal{G}(3) := \{c\}$ , one has



in the free clone  $\mathcal{T}(\mathcal{G})$ .



# Colored words

Let  $(\mathcal{M}, \cdot, \epsilon)$  be a monoid.

Let  $\mathcal{WM}$  be the graded set of all  $\mathcal{M}$ -colored words defined, for any  $n \geq 0$ , by

$$\mathcal{WM}(n) := \bigsqcup_{n \geq 0} \left\{ \binom{u}{c} : u \in [n]^\ell, c \in \mathcal{M}^\ell, \ell \geq 0 \right\}.$$

## – Example –

$$\binom{1 \ 2 \ 1 \ 6}{0 \ 0 \ 1 \ 0} \quad \text{is a } \mathbb{Z}/2\mathbb{Z}\text{-colored word.}$$

Let  $\odot$  be the superposition map defined by

$$\binom{u}{c} \odot \left[ \binom{v_1}{d_1}, \dots, \binom{v_n}{d_n} \right] := \binom{v_{u(1)} \cdots v_{u(\ell)}}{(c(1) \cdot d_{u(1)}) \cdots (c(\ell) \cdot d_{u(\ell)})}.$$

Let finally set  $\mathbb{1}_{i,n} := \binom{i}{\epsilon}$ .

# Clone of colored words

## – Example –

In  $W(\mathbb{N}, +, 0)$ ,

$$\begin{pmatrix} 2 & 2 & 3 \\ 0 & 1 & 0 \end{pmatrix} \odot \left[ \begin{pmatrix} 2 & 1 \\ 1 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 2 \\ 3 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 2 & 2 \\ 1 & 0 \end{pmatrix} \right] = \begin{pmatrix} 1 & 1 & 2 & 1 & 1 & 2 & 2 & 2 \\ 3 & 0 & 0 & 4 & 1 & 1 & 1 & 0 \end{pmatrix}.$$

## – Theorem [G., 2020–] –

For any monoid  $\mathcal{M}$ ,  $(W\mathcal{M}, \odot, \mathbb{1}_{i,n})$  is a clone.

The clone  $W\mathcal{M}$  is in fact the **clone counterpart** of the operad  $T\mathcal{M}$ .

This is due to the fact that they have both the **same presentation**.

# Clone of words and congruences

Let us focus on the case where  $\mathcal{M}$  is the trivial monoid  $\{\epsilon\}$ .

Let **Word** :=  $W\{\epsilon\}$ . We can forget about the colors of the elements of **Word** without any loss of information.

- Let  $\equiv_s$  be the equivalence relation on **Word** wherein  $u \equiv_s v$  if  $u$  and  $v$  have both the same **sorted** version.
- Let  $\equiv_l$  (resp.  $\equiv_r$ ) be the equivalence relation on **Word** wherein  $u \equiv_l v$  (resp.  $u \equiv_r v$ ) if the versions of  $u$  and  $v$  obtained by keeping only the **leftmost** (resp. **rightmost**) among the multiple occurrences of a same letter are equal.

## – Examples –

We have  $311322 \equiv_s 131232$ ,  $223111352 \equiv_l 2333315$ ,  $5142144 \equiv_r 552214$ .

## – Proposition –

The equivalence relations  $\equiv_s$ ,  $\equiv_l$ , and  $\equiv_r$  are clone congruences of **Word**.

# Multisets

Let  $\mathbf{MSet} := \mathbf{Word} / \equiv_s$ .

The elements of  $\mathbf{MSet}$  can be seen as multisets of positive integers. By encoding any such multiset  $u = [1^{a(1)}, \dots, n^{a(n)}]$  by the tuple  $a = (a(1), \dots, a(n))$ , the superposition map of  $\mathbf{MSet}$  expresses as a matrix multiplication

$$a \odot [b_1, \dots, b_n] = \begin{pmatrix} a(1) & \dots & a(n) \end{pmatrix} \begin{pmatrix} b_1(1) & \dots & b_1(m) \\ \vdots & \dots & \vdots \\ b_n(1) & \dots & b_n(m) \end{pmatrix}.$$

## – Proposition –

The clone  $\mathbf{MSet}$  admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\mathbf{a}\}$  and  $\mathfrak{R}$  satisfies

$$\begin{array}{c} \text{a} \\ / \quad \backslash \\ \text{a} \quad x_3 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \text{a} \\ / \quad \backslash \\ x_1 \quad \text{a} \\ / \quad \backslash \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} \text{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \text{a} \\ / \quad \backslash \\ x_2 \quad x_1 \end{array}.$$

Therefore,  $\mathbf{MSet}$  is a combinatorial realization of the variety of **commutative semigroups**.

# Arrangements

Let  $\mathbf{Arr}_1 := \mathbf{Word}/\equiv_1$ .

The elements of  $\mathbf{Arr}_1(n)$  can be seen as arrangements (words without repetitions) on  $[n]$ . For any  $n \geq 0$ ,

$$\#\mathbf{Arr}_1(n) = \sum_{0 \leq k \leq n} \frac{n!}{k!}$$

and this sequence starts by 1, 2, 5, 16, 65, 326, 1957, 13700, 109601.

## – Proposition –

The clone  $\mathbf{Arr}_1$  admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\mathbf{a}\}$  and  $\mathfrak{R}$  satisfies

$$\begin{array}{c} \begin{array}{c} | \\ \mathbf{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} | \\ \mathbf{a} \\ / \quad \backslash \\ x_1 \quad \mathbf{a} \\ \backslash \quad / \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} | \\ \mathbf{a} \\ / \quad \backslash \\ x_1 \quad x_1 \end{array} \mathfrak{R} |_{x_1}, \quad \begin{array}{c} | \\ \mathbf{a} \\ / \quad \backslash \\ \mathbf{a} \quad x_1 \\ \backslash \quad / \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} | \\ \mathbf{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array}.$$

The algebra of this variety are **left-regular bands**, that are idempotent semigroups wherein the operation  $\mathbf{a}$  satisfies the relation  $x_1 \mathbf{a} x_2 \mathbf{a} x_1 = x_1 \mathbf{a} x_2$ .

Analogous properties hold for the quotient  $\mathbf{Arr}_r := \mathbf{Word}/\equiv_r$ , leading to **right-regular bands**.

# Sets

## – Lemma –

$$\equiv_s \circ \equiv_l = \equiv_l \circ \equiv_s$$

Therefore, this composition is a clone congruence of **Word**.

Let us set it as  $\equiv_i$  and let **Set** := **Word**/ $\equiv_i$ .

The elements of **Set** can be seen as sets of positive integers. On such objects, the superposition map of **Set** expresses as

$$U \odot [V_1, \dots, V_n] = \bigcup_{j \in U} V_j.$$

Moreover, for any  $n \geq 0$ ,  $\#\mathbf{Set}(n) = 2^n$ .

## – Proposition –

The clone **Set** admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\mathbf{a}\}$  and  $\mathfrak{R}$  satisfies

$$\begin{array}{c} \mathbf{a} \\ / \quad \backslash \\ \mathbf{a} \quad x_3 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \mathbf{a} \\ / \quad \backslash \\ x_1 \quad \mathbf{a} \\ / \quad \backslash \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} \mathbf{a} \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \mathbf{a} \\ / \quad \backslash \\ x_2 \quad x_1 \end{array}, \quad \begin{array}{c} \mathbf{a} \\ / \quad \backslash \\ x_1 \quad x_1 \end{array} \mathfrak{R} \mid_{x_1}.$$

Therefore, **Set** is a combinatorial realization of the variety of **semilattices**.

# Arrangements of blocks

Let us consider some intersections involving the congruences  $\equiv_s$ ,  $\equiv_l$ , and  $\equiv_r$ .

Let  $\equiv_{sl} := \equiv_s \cap \equiv_l$  and  $\mathbf{ArrB}_l := \mathbf{Word} / \equiv_{sl}$ .

The elements of  $\mathbf{ArrB}_l(n)$  can be seen as arrangements of possibly empty blocks of repeated letters of  $[n]$ .

## – Examples –

The word 3311115526 is such an element of  $\mathbf{ArrB}_l(9)$ . The word 22222333112 is not an element of  $\mathbf{ArrB}_l$ .

## – Proposition –

The clone  $\mathbf{ArrB}_l$  admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\mathbf{a}\}$  and  $\mathfrak{R}$  satisfies

$$\begin{array}{c} \mathbf{a} \\ | \\ \mathbf{a} \quad x_3 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \mathbf{a} \\ | \\ x_1 \quad \mathbf{a} \\ / \quad \backslash \\ x_2 \quad x_3 \end{array}, \quad \begin{array}{c} \mathbf{a} \\ | \\ \mathbf{a} \quad x_1 \\ / \quad \backslash \\ x_1 \quad x_2 \end{array} \mathfrak{R} \begin{array}{c} \mathbf{a} \\ | \\ \mathbf{a} \quad x_2 \\ / \quad \backslash \\ x_1 \quad x_1 \end{array}.$$

Analogous properties hold for the quotient  $\mathbf{ArrB}_r := \mathbf{Word} / \equiv_{sr}$ , where  $\equiv_{sr} := \equiv_s \cap \equiv_r$ .

# Pairs of compatible arrangements

Let  $\equiv_{lr} := \equiv_l \cap \equiv_r$  and  $\mathbf{PArr} = \mathbf{Word} / \equiv_{lr}$ .

The elements of  $\mathbf{PArr}(n)$  can be seen as pairs  $(u, v)$  such that  $u$  and  $v$  are arrangements on  $[n]$ , such that  $j$  appears in  $u$  iff  $j$  appears in  $v$ .

For any  $n \geq 0$ ,

$$\#\mathbf{PArr}(n) = \sum_{0 \leq k \leq n} \frac{n!k!}{(n-k)!}$$

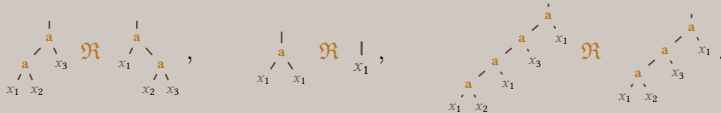
and this sequence starts by 1, 2, 7, 52, 749, 17686, 614227, 29354312, 1844279257.

## – Example –

$(3261, 1263)$  is such an element of  $\mathbf{PArr}(6)$ .

## – Proposition –

The clone  $\mathbf{PArr}$  admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{\mathbf{a}\}$  and  $\mathfrak{R}$  satisfies



Therefore,  $\mathbf{PArr}$  is a combinatorial realization of the variety of **regular bands**.



# Pairs of compatible arrangements of blocks

Let  $\equiv_{\text{slr}} := \equiv_s \cap \equiv_l \cap \equiv_r$  and  $\mathbf{PArrB} := \mathbf{Word} / \equiv_{\text{slr}}$ .

The elements of  $\mathbf{PArrB}(n)$  can be seen as pairs  $(u, v)$  such that  $u$  and  $v$  are arrangements of possibly empty blocks of repeated letters on  $[n]$ , with  $u$  and  $v$  having the same number of occurrences of any letter.

## – Example –

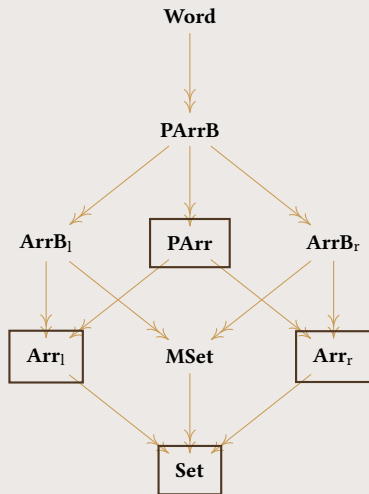
$(3222611, 22211263)$  is such an element of  $\mathbf{PArrB}(6)$ .

## – Proposition –

The clone  $\mathbf{PArrB}$  admits the presentation  $(\mathfrak{G}, \mathfrak{R})$  where  $\mathfrak{G} := \mathfrak{G}(2) := \{a\}$  and  $\mathfrak{R}$  satisfies



# Full diagram



Squared clones are combinatorial.