# Combinatorics of braids 

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## A PhD about braids?

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Some questions of interest...

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(1) What are braids?

## A PhD about braids?



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(1) What are braids?

Mathematical objects interacting with each other.

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(1) What are braids? Mathematical objects interacting with each other.
(2) What is a complicated braid?

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Define notions of complexity.

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(1) What are braids? Mathematical objects interacting with each other.
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(3) How do complicated braids typically behave?

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## Some questions of interest. . .

(1) What are braids? Mathematical objects interacting with each other.
(2) What is a complicated braid?

Define notions of complexity.
(3) How do complicated braids typically behave?

Choose a dynamic framework/probability measure.

## What are braids? - Algebra

## Isotopy classes of braid diagrams (Artin 1926)



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Finitely generated monoid

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\mathcal{B}_{n}=\langle
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What are braids? - Geometry

## Isotopy classes of laminations of the punctured plane (Birman 1975)

Trivial lamination


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Braid acting on a lamination


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Braids are isotopy classes of which laminations?

Open lamination


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What are braids? - Checking braid equality
Garside normal form (Garside 1969, Adian 1984)
(1) The monoid of positive braids $\mathcal{B}_{n}^{+}=\left\langle\sigma_{1}, \ldots, \sigma_{n-1}\right\rangle^{+}$is a lattice for the divisibility ordering $\leqslant . \quad\left(\alpha \leqslant \beta \Leftrightarrow \exists \gamma \in \mathcal{B}_{n}^{+}, \alpha \gamma=\beta\right)$

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\left(\alpha \leqslant \beta \Leftrightarrow \exists \gamma \in \mathcal{B}_{n}^{+}, \alpha \gamma=\beta\right)
$$

(2) There exists a Garside element $\Delta_{n}=\bigvee\left\{\sigma_{1}, \ldots, \sigma_{n-1}\right\}$.
(3) The Garside normal form of a positive braid $\alpha \in \mathcal{B}_{n}^{+}$is the smallest word $\operatorname{Gar}(\alpha)=a_{1} \cdot a_{2} \cdot \ldots \cdot a_{k}$ such that:

$$
\alpha=a_{1} a_{2} \ldots a_{k} ; \quad \quad a_{i}=\Delta_{n} \wedge\left(\left(a_{1} \ldots a_{i-1}\right)^{-1} \alpha\right) .
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The Garside normal form:

- can be extended to the group $\mathcal{B}_{n}$;
- is automatic: for all $i \in\{1, \ldots, n-1\}$, the languages

$$
\left\{\left(\mathbf{G} \operatorname{ar}(\alpha), \mathbf{G a r}\left(\alpha \sigma_{i}\right)\right): \alpha \in \mathcal{B}_{n}\right\} ; \quad \bullet\left\{\left(\mathbf{G} \operatorname{ar}(\alpha), \mathbf{\operatorname { G a r }}\left(\sigma_{i} \alpha\right)\right): \alpha \in \mathcal{B}_{n}\right\}
$$ are regular;

- solves the equality problem: $\alpha=\beta$ iff $\operatorname{Gar}(\alpha)=\boldsymbol{\operatorname { G a r }}(\beta)$.

What are braids? - Checking braid equality
Tight laminations/curve diagrams
A lamination/curve diagram is tight if it minimises crossings $\pm$ or $\|^{\boldsymbol{\sim}}$.

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A lamination/curve diagram is tight if it minimises crossings $\underset{+}{\perp}$ or $-\mid$.
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- Two tight laminations/curve diagrams represent the same braid iff they are visibly isotopic to each other.


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What are braids? - Checking complexity
Which braid is the most complicated?


What are braids? - Checking complexity
Which braid is the most complicated?


Several approaches to braid complexity

- "Naive" Artin length: $\quad \mathrm{E}>\mathrm{F}>\mathrm{D}>\mathrm{A} \approx \mathrm{B} \approx \mathrm{C}$;
- $\mathbf{A}=\sigma_{1} \sigma_{2} \sigma_{1}:$
$|\mathbf{A}|=3 ;$
- $\mathbf{D}=\sigma_{1}^{4}$ :
$|D|=4 ;$
- $\mathbf{B}=\sigma_{1} \sigma_{2} \overline{\sigma_{1}}:$
$|B|=3 ;$
- $\mathbf{E}=\left(\sigma_{1} \sigma_{2}\right)^{3}$ :
$|E|=6 ;$
- $\mathbf{C}=\sigma_{1} \overline{\sigma_{2}} \sigma_{1}$ :
$|\mathbf{C}|=3 ;$
- $\mathbf{F}=\sigma_{1}^{2} \sigma_{2} \overline{\sigma_{1} \sigma_{2}}$ :
$|\mathbf{F}|=5$.

What are braids? - Checking complexity
Which braid is the most complicated?


Several approaches to braid complexity

- "Naive" Artin length:
$\mathrm{E}>\mathrm{F}>\mathrm{D}>\mathrm{A} \approx \mathrm{B} \approx \mathrm{C}$;
- "Real" Artin length:
$\mathrm{E}>\mathrm{D}>\mathrm{A} \approx \mathrm{B} \approx \mathrm{C}=\mathrm{F}$;
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- Symmetric Garside length:
$\mathrm{D}>\mathrm{C}=\mathrm{F}>\mathrm{B} \approx \mathrm{E}>\mathrm{A}$;
- $\mathbf{A}=\sigma_{1} \sigma_{2} \sigma_{1}$ :
$|\mathbf{A}|=1 ;$
- $\mathbf{D}=\sigma_{1} \cdot \sigma_{1} \cdot \sigma_{1} \cdot \sigma_{1}:$
$|D|=4 ;$
- $\mathbf{B}=\overline{\sigma_{1}} \cdot \sigma_{1} \sigma_{2}$ :
$|B|=2 ;$
- $E=\Delta_{3} \cdot \Delta_{3}$ :
$|E|=2 ;$
- $\mathbf{C}=\overline{\sigma_{2} \sigma_{1}} \cdot \sigma_{2} \sigma_{1} \cdot \sigma_{1}$ :
$|\mathbf{C}|=3 ;$
- $\mathrm{F}=\mathrm{C}$ :
$|\mathbf{F}|=3$.

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- Symmetric Garside length:
$\mathrm{D}>\mathrm{C}=\mathrm{F}>\mathrm{B} \approx \mathrm{E}>\mathrm{A} ;$
- Open laminated complexity: $\mathbf{C}=\mathrm{F}>\mathrm{D} \approx \mathrm{E}>\mathrm{B}>\mathrm{A}$;
$|A|=6$
$|\mathbf{B}|=8$
$|\mathbf{C}|=|\mathbf{F}|=14$
$|\mathbf{D}|=10$
$|\mathbf{E}|=10$


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Several approaches to braid complexity

- "Naive" Artin length:
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- Open laminated complexity: $\mathbf{C}=\mathrm{F}>\mathrm{D} \approx \mathrm{E}>\mathrm{B}>\mathrm{A}$;
- Diagrammatic complexity: $\mathrm{C}=\mathrm{F}>\mathrm{D} \approx \mathrm{E}>\mathrm{B}>\mathrm{A}$.




## What are braids? - Checking complexity

## How fast can you compute the complexity of a braid $\alpha \in \mathcal{B}_{n}$ of length $k$ ?

- Artin length:

$$
\begin{array}{lr}
\text { coNP-complete }(n, k) & \text { (Paterson \& Razborov 1991); } \\
\text { polynomial }(n \leqslant 3, k) & \text { (Sabalka 2003); } \\
\text { open }(n \geqslant 5, k) ; &
\end{array}
$$

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## How fast can you compute the complexity of a braid $\alpha \in \mathcal{B}_{n}$ of length $k$ ?

- Artin length: coNP-complete( $n, k$ ) (Paterson \& Razborov 1991); polynomial $(n \leqslant 3, k) \quad$ (Sabalka 2003); open $(n \geqslant 5, k)$;
- Symmetric Garside length: polynomial $(n, k)$
(Thurston 1988);


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(Thurston 1988);
- Open laminated complexity: polynomial( $n, k$ ) (Dynnikov \& Wiest 2004);


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- Symmetric Garside length: polynomial $(n, k)$
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(Dynnikov \& Wiest 2004);
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Algebra
Geometry

Class of words
$\mathcal{B}_{n}=\langle$ generators $|$ relations $\rangle$
Garside normal form (regular)

Class of drawings
$\mathcal{B}_{n}=\{$ tight drawings $\}$

Algebra
Geometry

## Braid

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Relaxation normal form (regular) Part 1/4

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Relaxation normal form (regular) Part 1/4


Garside: $\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|}$ rational Relaxation: $\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|}$ rational
Artin: $\sum_{\alpha \in \mathcal{B}_{3}} z^{|\alpha|}$ rational
Artin: $\sum_{\alpha \in \mathcal{B}_{n \geqslant 4}} z^{|\alpha|}$ ?

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Garside: $\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|}$ rational
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Artin: $\sum_{\alpha \in \mathcal{B}_{n \geqslant 4}} z^{|\alpha|}$ ?
Geometric: $\sum_{\alpha \in \mathcal{B}_{3}} z^{|\alpha|} \quad \neg$ rational
Part 2/4 $\quad$-holonomic
Geometric: $\sum_{\alpha \in \mathcal{B}_{n \geqslant 4}} z^{|\alpha|}$ ?

## Depth-first exploration



Random walk

Which normal forms converge?
(Vershik, 2000)

Uniform measure on positive braids of given (Artin) size

What do Garside normal forms of large random braids look like?
(Gebhardt \& Tawn, 2013)

Markov-Ivanovsky normal form
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## Depth-first exploration <br> Width-first exploration



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## Depth-first exploration <br> Width-first exploration

## Braids of large size

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Part 3/4
(with J. Mairesse)

Uniform measure on infinite positive braids
Part 4/4 (with S. Abbes, S. Gouëzel \& J. Mairesse)

## Contents

(1) Geometric aspects of braids

- Right relaxation normal form
- Counting braids with a given geometric complexity
(2) Algebraic aspects of braids
- Garside normal form and random walks
- Drawing infinite braids uniformly at random
(3) Conclusion

What is the right relaxation normal form?

Move your rightmost tensed puncture and relax!

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While your lamination is not trivial:

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While your lamination is not trivial:
(1) Select the rightmost (mobile) puncture that lies inside a bigon;

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While your lamination is not trivial:
(1) Select the rightmost (mobile) puncture that lies inside a bigon;
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Moves performed:
[2ص3]

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$$

Relaxation normal form (RNF):
$[1 \triangleleft 2] \cdot[2 \frown 3] \cdot[2 \frown 3]$
$[k \backsim \ell]=\sigma_{k} \ldots \sigma_{\ell-1}$
$[k \frown \ell]=\overline{\sigma_{k}} \ldots \overline{\sigma_{\ell-1}}$

While your lamination is not trivial:
(1) Select the rightmost (mobile) puncture that lies inside a bigon;
(2) Slide it along its right (or left) neighbour arc (and remember it);
(3) Relax your diagram!

Tight closed lamination, cell map and lamination/arc trees
Tight closed lamination


Tight closed lamination, cell map and lamination/arc trees

Tight closed lamination


Cell map


Tight closed lamination, cell map and lamination/arc trees

Tight closed lamination


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Tight closed lamination, cell map and lamination/arc trees

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Tight closed lamination


Lamination trees (LT)


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Tight closed lamination


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Lamination trees (LT)


Cell map


Arc trees


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Cell map


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Tight closed lamination, cell map and lamination/arc trees

Tight closed lamination


Lamination trees (LT)


Cell map


Arc trees


Tight closed lamination, cell map and lamination/arc trees

Tight closed lamination


Lamination trees (LT)


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## After a careful case analysis. . .

Given two braids $\alpha \in \mathcal{B}_{n}$ and [ $k \frown \ell$ ], remembering small-size subtrees $\mathbf{l t}(\alpha)$ of $\mathbf{L T}(\alpha)$ is enough to:

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## Theorem (J. 2015)

The right relaxation normal form is regular.

## Going further

## Some additional results

- Memory requirements: nearly optimal (up to a ratio $\leqslant 20$ );
- Dehornoy ordering: $\sigma$-positivity $\Leftrightarrow$ RNF in a regular language.


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- Memory requirements: nearly optimal (up to a ratio $\leqslant 20$ );
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## and open questions

- Is the right relaxation normal form (bi-)automatic?
(Yes if $n \leqslant 3$ )
- Regularity of other transmission-relaxation normal forms? (wide open)


## Contents

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Knowing the open laminated complexity of $\alpha$, can we compute its:

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Tight open lamination


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\end{array}
$$

Tight open laminations

$\left|\sigma_{1} \overline{\sigma_{2} \sigma_{3}}\right|_{o}=11$

$\left|\sigma_{1}^{2}{\overline{\sigma_{3}}}^{2}\right|_{o}=11$

Tight curve diagrams


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Tight open lamination


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Inverse tight curve diagram


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$$

Let us compute geometric generating functions!

$$
\begin{gathered}
\mathcal{O}_{n}(z)=\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|_{0}} \quad \mathcal{C}_{n}(z)=\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|_{c}} \quad \mathcal{D}_{n}(z)=\sum_{\alpha \in \mathcal{B}_{n}} z^{|\alpha|_{d}} \\
\mathcal{C}_{n}(z)=z^{n+3} \mathcal{O}_{n}(z)=z^{n+3} \mathcal{D}_{n}(z)
\end{gathered}
$$

## Generalised tight curve diagrams and coordinates

Generalising tight curve diagrams

Tight curve diagram


Tight generalised curve diagram


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## Generalised tight curve diagrams and coordinates

## Generalising tight curve diagrams

Tight curve diagram


Tight generalised curve diagram

and encoding them!


Coordinates: $\left\langle\left(x_{0}, x_{1}, x_{2}, x_{3}, x_{4}\right),\left(y_{1}, y_{2}, y_{3}, y_{4}\right)\right\rangle$

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$$
\sum_{i=1}^{n-1} x_{i}=\frac{|\alpha|_{d}+1-n}{2}: \quad \text { Let us compute } \mathcal{G}_{n}(z)=\sum_{k \geqslant 0} g_{n, k} z^{k}=\sqrt{z}^{1-n} \mathcal{D}_{n}(\sqrt{z})!
$$

## And finally...

## Theorem (J. 2015)

In the 3 -strand braid group $\mathcal{B}_{3}$, we have:

- $\mathcal{G}_{3}(z)=2 \frac{1+2 z-z^{2}}{z^{2}\left(1-z^{2}\right)}\left(\sum_{k \geqslant 3} \varphi(k) z^{k}\right)+\frac{1-3 z^{2}}{1-z^{2}}$ and
- $g_{3, k}=\mathbf{1}_{k=0}+2\left(\varphi(k+2)-\mathbf{1}_{k \in 2 \mathbb{Z}}+2 \sum_{i=1}^{\lfloor k / 2\rfloor} \varphi(k+3-2 i)\right) \mathbf{1}_{k \geqslant 1}$,
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$\mathcal{G}_{3}(z)$ is more complicated than $\sum_{\alpha \in \mathcal{B}_{3}} z^{|\alpha|_{\text {Artin }}}=\frac{(1+z)\left(1-z+z^{2}-2 z^{3}\right)}{(2-z)(1-2 z)\left(1-z-z^{2}\right)}$ !


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## (3) Conclusion

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## Simple positive braids

Simple positive braids are:

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(for all $\alpha \in \mathcal{S}_{n}$ and $\sigma_{i}, \alpha \geqslant \sigma_{i} \Leftrightarrow \alpha \sigma_{i} \notin \mathcal{S}_{n}$ ).

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onsequence on Garside normal forms:
(9) local neighbouring criterion: $w_{1} \cdot w_{2} \cdot \ldots \cdot w_{k} \in \operatorname{Gar}\left(\mathcal{B}_{n}^{+}\right)$iff

- $w_{1}, \ldots, w_{k} \in \mathcal{S}_{n} \backslash\{\mathbf{1}\} ;$
- $\mathbf{R}\left(w_{i}\right) \supseteq \mathbf{L}\left(w_{i+1}\right)$ for all $i<k$.

$$
\left(\mathbf{L}(\alpha)=\left\{\sigma_{i}: \sigma_{i} \leqslant \alpha\right\} \text { and } \mathbf{R}(\alpha)=\left\{\sigma_{i}: \alpha \geqslant \sigma_{i}\right\}\right)
$$

## Generalising braid monoids



- Braid monoid: $\left\langle\sigma_{i} \mid \sigma_{i} \sigma_{i+1} \sigma_{i}=\sigma_{i+1} \sigma_{i} \sigma_{i+1}, i \neq j \pm 1 \Rightarrow \sigma_{i} \sigma_{j}=\sigma_{j} \sigma_{i}\right\rangle^{+}$;


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$$
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$$
\ell(i, j)=+\infty \Rightarrow \text { no relation! }
$$

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- Trace monoid: Artin-Tits with $\ell(i, j) \in\{2,+\infty\}$;


## Generalising braid monoids



- Braid monoid: $\left\langle\sigma_{i} \mid \sigma_{i} \sigma_{i+1} \sigma_{i}=\sigma_{i+1} \sigma_{i} \sigma_{i+1}, i \neq j \pm 1 \Rightarrow \sigma_{i} \sigma_{j}=\sigma_{j} \sigma_{i}\right\rangle^{+}$;
- Garside monoid: finite generating set $\Sigma$, cancellative, $\leqslant-$ and $\geqslant$-lattice, with a Garside element $\Delta$ such that $\{x: x \leqslant \Delta\}=\{x: \Delta \geqslant x\} \supseteq \Sigma$;
- Artin-Tits monoid: $\left\langle\sigma_{i} \mid\left[\sigma_{i} \sigma_{j}\right]^{\ell(i, j)}=\left[\sigma_{j} \sigma_{i}\right]^{\ell(i, j)}\right\rangle^{+}$;
- A-T monoid with spherical type: A-T and Garside;
- Trace monoid: Artin-Tits with $\ell(i, j) \in\{2,+\infty\}$;
- A-T monoid with FC type: A-T with finite 2-way Garside family.


## Contents

(1) Geometric aspects of braids

- Right relaxation normal form
- Counting braids with a given geometric complexity
(2) Algebraic aspects of braids
- Garside normal form and random walks
- Drawing infinite braids uniformly at random


## (3) Conclusion

## Random walk in a braid monoid

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(1) Select i.i.d. generators $\left(Y_{k}\right)_{k \geqslant 0}$ uniformly chosen in $\left\{\sigma_{1}, \ldots, \sigma_{n-1}\right\}$.

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Several Garside normal forms:
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Chronological $\mathcal{L}$ form:
Left Garside $工$
normal form: 二
$1^{\text {st }}$ step
Right Garside $\mathcal{X}$
normal form: $=$
Left $^{\Delta}$ Garside $\mathcal{X}$ normal form:

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Chronological $\sqrt{25}$
form:
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normal form:
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normal form:


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Blocking patterns: Going to infinity...
Blocking pattern
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Theorem (J. \& Mairesse 2016+)
Prefixes of the words $\operatorname{Gar}_{r}\left(X_{k}\right)_{k \geqslant 0}$ almost surely converge.
... and beyond!

What is our limit object? How fast do we reach it?
... and beyond!

## What is our limit object? How fast do we reach it?

(1) Limit of an infinite-state Markov chain with $L^{1}$ factors;

... and beyond!

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(1) Limit of an infinite-state Markov chain with $L^{1}$ factors;
(2) Ergodic process;
... and beyond!

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... and beyond!

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Computing $\operatorname{Gar}_{r}\left(X_{k+1}\right)$ when knowing $\operatorname{Gar}_{r}\left(X_{k}\right)$ and $Y_{k}$ in expected time $\mathcal{O}(k)$.
... and beyond!

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## Going even further

## Generalised framework

- Braid monoid


## Going even further

## Generalised framework

- Braid monoid $\Rightarrow$ irreducible $A-T$ monoid with spherical type;


## Going even further

## Generalised framework

- Braid monoid $\Rightarrow$ irreducible A-T monoid with spherical type $\Rightarrow$ irreducible $A-T$ group with spherical type;


## Going even further

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(wide open) (wide open)


## Contents

(1) Geometric aspects of braids

- Right relaxation normal form
- Counting braids with a given geometric complexity
(2) Algebraic aspects of braids
- Garside normal form and random walks
- Drawing infinite braids uniformly at random
(3) Conclusion

(with S. Abbes,
S. Gouëzel \&
J. Mairesse)
uniformly at random in $\mathcal{B}_{n}^{k}=\left\{\beta \in \mathcal{B}_{n}^{+}:|\beta|_{\text {Artin }}=k\right\}$

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Smoothening uniform measures on spheres
(1) Choose $\mu_{p}: \alpha \mapsto H_{n}(p) p^{|\alpha|}$

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## Theorem (Abbes, Gouëzel, J. \& Mairesse 2016+)

Uniform probability measures on $\mathcal{B}_{n}^{k}$ converge weakly towards $\mu_{\infty}$ when $k \rightarrow+\infty$.
$\mu_{\infty}$ is a uniform probability measure on infinite braids!

## Going further

## Stable region conjectures (Gebhardt \& Tawn 2014)

(1) The words $\left\{\operatorname{Gar}_{\ell}(\beta) \mid \beta \in \partial \mathcal{B}_{n}^{+}\right\}$contain a geometric number of $\Delta_{n}$;


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Monoid $\mathbb{N} * \mathbb{N}$


Monoid

$$
\langle a, b \mid a b a b=b a b a\rangle^{+}
$$



Monoid $\langle a, b, c \mid a c=c a\rangle^{+}$

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## Uniform measures

(1) Uniform measures on positive spheres converge towards a simple critical Markov process (for all irreducible A-T monoids of FC type).


Do you have questions?

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