

Operads, enumeration, and constructions

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Discrete Structures Days

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Combinatorial sets

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Example

The set \mathfrak{S} of the permutations, where the size of a permutation is its length as a word, satisfies

$$\mathfrak{S}(0) = \{\epsilon\}, \quad \mathfrak{S}(1) = \{1\}, \quad \mathfrak{S}(2) = \{12, 21\}, \quad \mathfrak{S}(3) = \{123, 132, 213, 231, 312, 321\},$$

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and its generating series satisfies

$$\mathcal{G}_{\mathfrak{S}}(t) = \sum_{n \in \mathbb{N}} n!t^n = 1 + t + 2t^2 + 6t^3 + 24t^4 + 120t^5 + \dots$$

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On \mathfrak{S} , let $s(\sigma)$ be the number of descents of σ (i.e., the number of i such that $\sigma(i) > \sigma(i+1)$). Then,

$$\mathcal{G}_{\mathfrak{S}}^s(t, q) = 1 + t + (1 + q)t^2 + (1 + 4q + q^2)t^3 + (1 + 11q + 11q^2 + q^3)t^5 + \dots$$

The coefficients of the t^n are the Eulerian polynomials.

Algebraic approach

Aims

Let C be a combinatorial set. We would like to

1. describe the generating series of C ;
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Strategy

Endow C with the structure of an operad \mathcal{O} .

As main benefits,

- ▶ the operations of \mathcal{O} lead to operations on series;
- ▶ presentations of \mathcal{O} by generators and relations highlight elementary building blocks and branching rules for the objects of C ;
- ▶ operad morphisms involving \mathcal{O} show connections between the objects of C and other combinatorial sets.

Formal power series

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Endowed with the pointwise addition

$$\langle x, f + g \rangle := \langle x, f \rangle + \langle x, g \rangle$$

and multiplication by a scalar

$$\langle x, \lambda f \rangle := \lambda \langle x, f \rangle,$$

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The **sum notation** of f is

$$f = \sum_{x \in X} \langle x, f \rangle x.$$

Formal power series and algebraic structures

When X is endowed with an algebraic structure, its operations transfer on X -series. For instance, if $*$: $X \times X \rightarrow X$ is a binary product on X ,

$$\mathbf{f} * \mathbf{g} := \sum_{x,y \in X} \langle x, \mathbf{f} \rangle \langle y, \mathbf{g} \rangle x * y.$$

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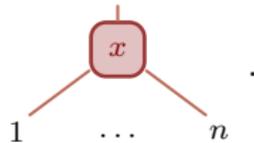
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The most common examples are

Structure on X	Series	Reference
$(\mathbb{N}, +, 0)$	Usual series of $\mathbb{K} \langle\langle t \rangle\rangle$	
Free comm. monoid	Multivariate series	
(A^*, \cdot, ϵ)	Noncomm. series	[Eilenberg, 1974]
Monoid	Series on monoids	[Salomaa, Soittola, 1978]
Operad	Series on operads	[Chapoton, 2002, 2008]

Operators

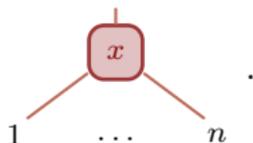
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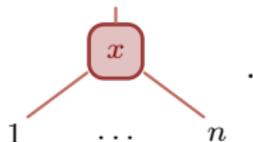
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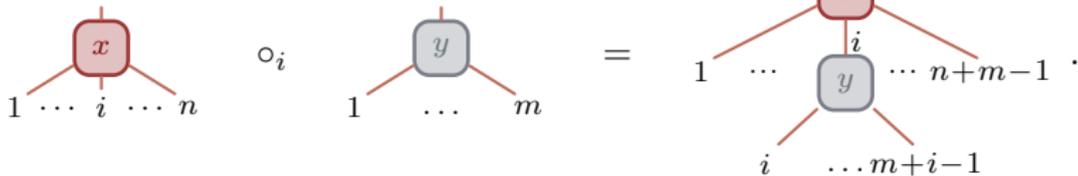


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Composing two operators x et y consists in

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This produces a new operator $x \circ_i y$ of arity $n + m - 1$:



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An (nonsymmetric set-)operad is a triple $(\mathcal{O}, \circ_i, \mathbb{1})$ such that

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This data satisfies some axioms.

Axioms of operads

Associativity:

$$(x \circ_i y) \circ_{i+j-1} z = x \circ_i (y \circ_j z)$$

$$x \in \mathcal{O}(n), y \in \mathcal{O}(m), z \in \mathcal{O}$$

$$i \in [n], j \in [m]$$

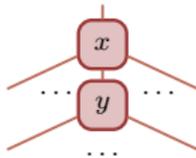
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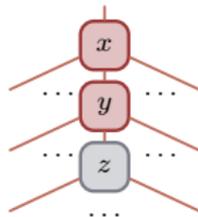
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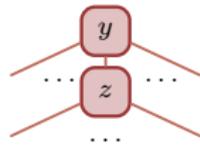
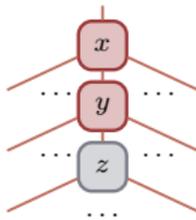
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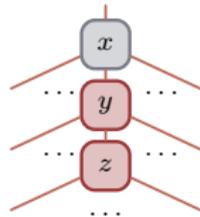
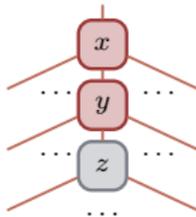
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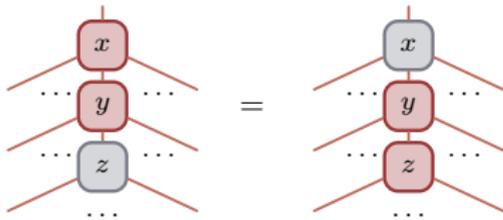
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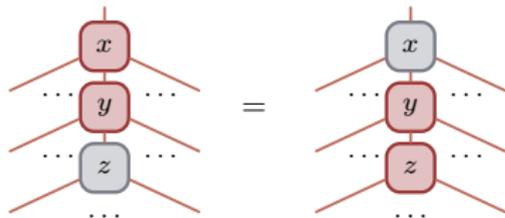
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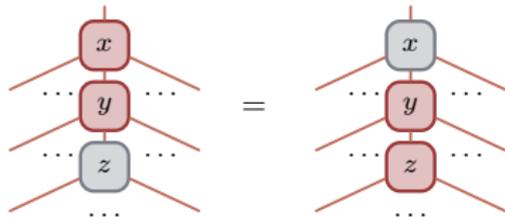
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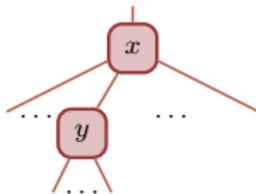


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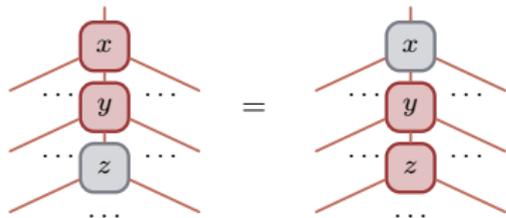
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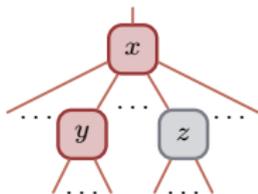


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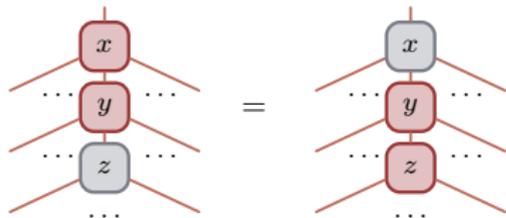
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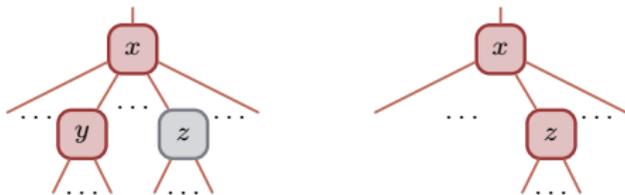


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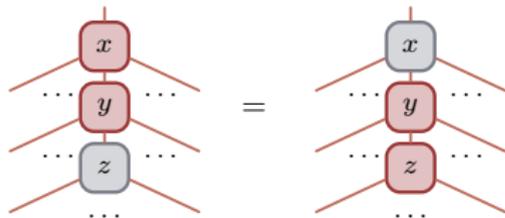
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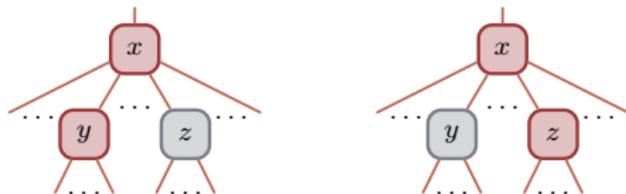


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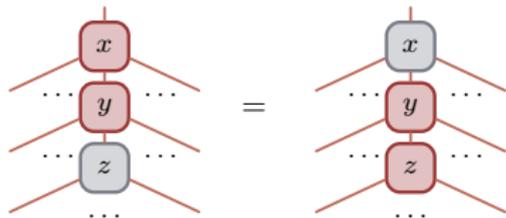
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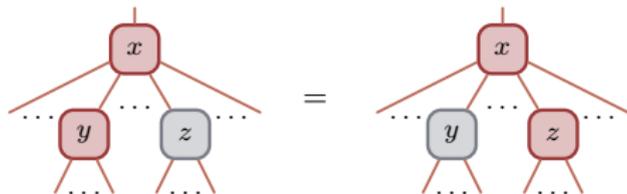


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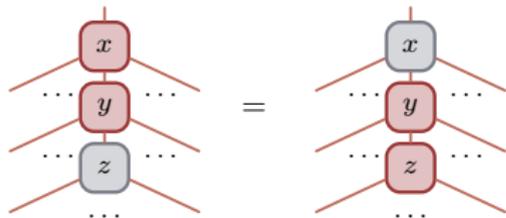
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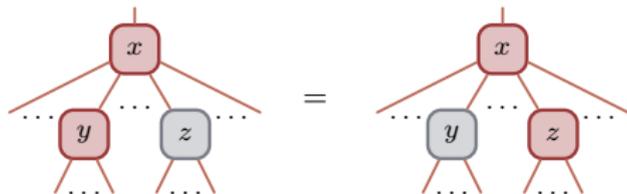


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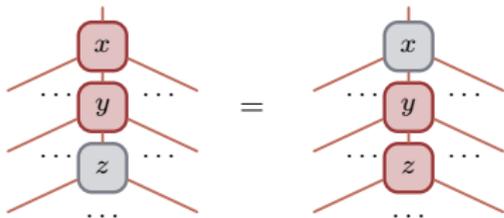
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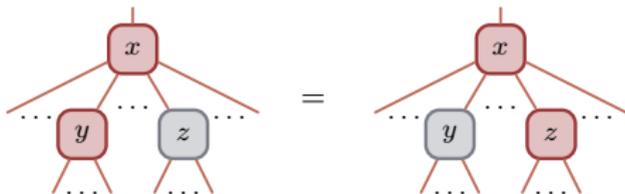


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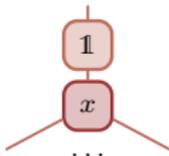


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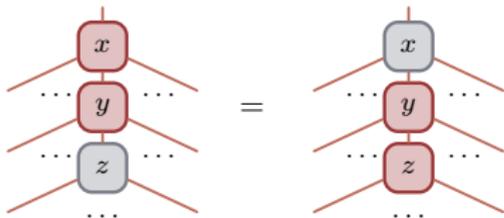
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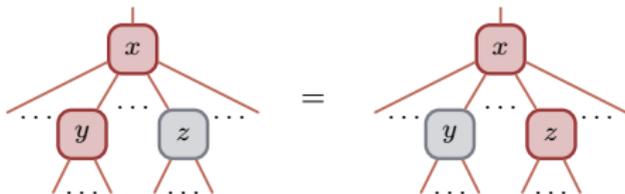


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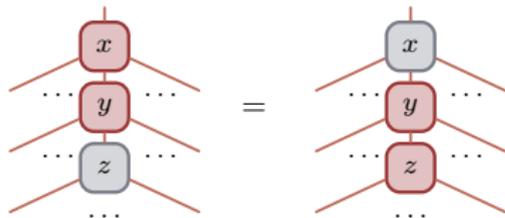
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$$(x \circ_i y) \circ_{i+j-1} z = x \circ_i (y \circ_j z)$$

$$x \in \mathcal{O}(n), y \in \mathcal{O}(m), z \in \mathcal{O}$$

$$i \in [n], j \in [m]$$

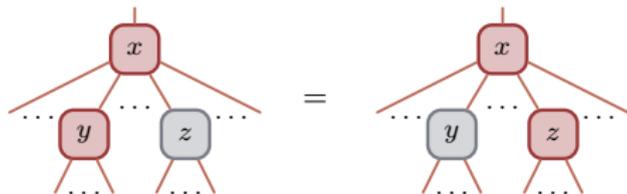


Commutativity:

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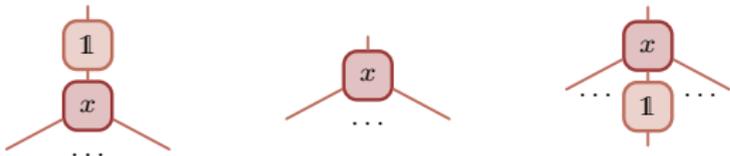


Unitality:

$$\mathbb{1} \circ_1 x \quad x \quad x \circ_i \mathbb{1}$$

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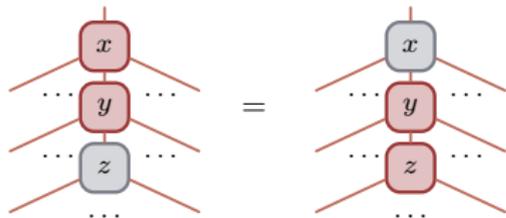
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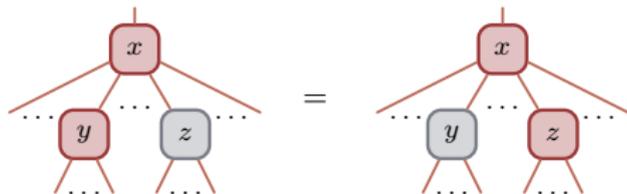


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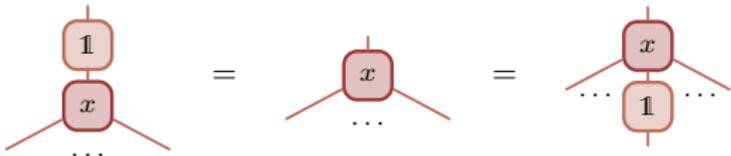


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Complete composition

Let \mathcal{O} be an operad.

The **complete composition map** of \mathcal{O} is the map

$$\circ : \mathcal{O}(n) \times \mathcal{O}(m_1) \times \cdots \times \mathcal{O}(m_n) \rightarrow \mathcal{O}(m_1 + \cdots + m_n)$$

defined by

$$x \circ [y_1, \dots, y_n] := (\dots ((x \circ_n y_n) \circ_{n-1} y_{n-1}) \dots) \circ_1 y_1.$$

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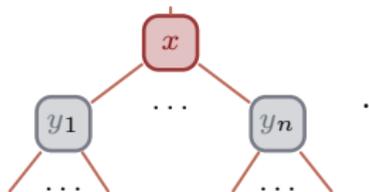
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Graphically, $x \circ [y_1, \dots, y_n]$ reads as



Formal power series on operads

Let \mathcal{O} be an operad.

The **characteristic series** of \mathcal{O} is the \mathcal{O} -series

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Goals

Use the partial and complete compositions of \mathcal{O} to obtain

- ▶ an expression for $\mathbf{f}_{\mathcal{O}}$;
- ▶ an expression for $\mathcal{G}_{\mathcal{O}}(t)$;
- ▶ new statistics on the objects of \mathcal{O} .

Free operads

Let $\mathfrak{G} := \sqcup_{n \geq 1} \mathfrak{G}(n)$ be a graded set.

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Example

Let $\mathfrak{G} := \mathfrak{G}(2) \sqcup \mathfrak{G}(3)$ with $\mathfrak{G}(2) := \{a, b\}$ and $\mathfrak{G}(3) := \{c\}$.

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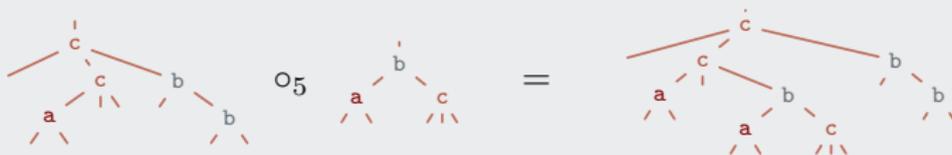
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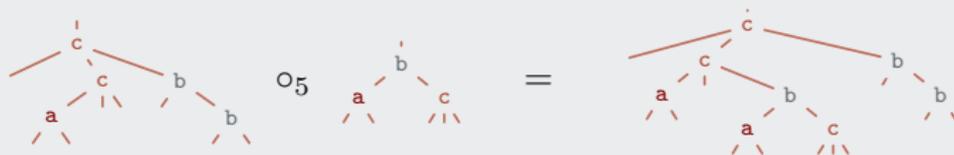
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- ▶ The unit is the tree with exactly one leaf.

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- ▶ **binary** when \mathfrak{G} is concentrated in arity 2;
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Poincaré-Birkhoff-Witt bases

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Let As be the **associative operad**, defined as the operad having the presentation (\mathfrak{G}, \equiv) where $\mathfrak{G} := \mathfrak{G}(2) := \{a\}$ and \equiv is the smallest operad congruence satisfying

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A PBW basis of As is the set of the syntax trees on \mathfrak{G} avoiding $a \circ_1 a$.

Contents

Operads and enumeration

Operad of compositions

The operad of compositions is the operad \mathbf{Comp} such that

- ▶ $\mathbf{Comp}(n)$ is the set of the ribbon diagrams of compositions of n .

Example



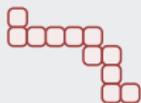
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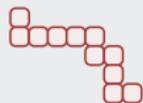


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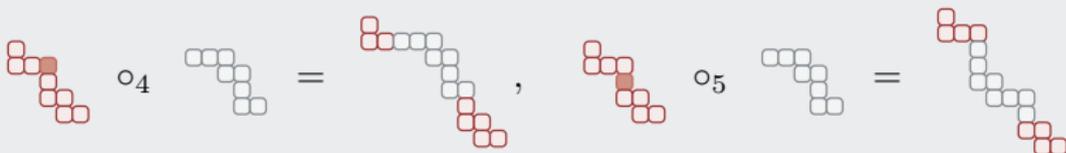
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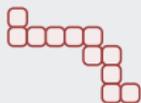


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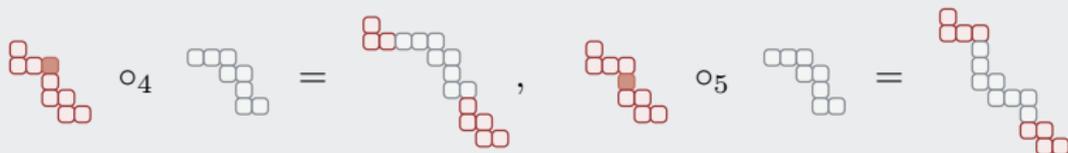
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Example



- ▶ The unit is \square .

Presentation of Comp

Proposition

The operad **Comp** admits the presentation (\mathfrak{G}, \equiv) where

$$\mathfrak{G} := \left\{ \begin{array}{c} \square \square \\ \square \end{array}, \begin{array}{c} \square \\ \square \end{array} \right\}$$

and \equiv is the smallest operad congruence satisfying

$$\begin{array}{c} \square \square \\ \square \end{array} \circ_1 \begin{array}{c} \square \square \\ \square \end{array} \equiv \begin{array}{c} \square \square \\ \square \end{array} \circ_2 \begin{array}{c} \square \square \\ \square \end{array},$$

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This presentation is binary and quadratic.

PBW basis of Comp and enumeration

Proposition

The set of the syntax trees on \mathfrak{G} avoiding

$$\begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array}, \quad \begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array}, \quad \begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array}, \quad \begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array}$$

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This leads to

$$f_{\text{Comp}} = \square + \begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array} \circ [\square, f_{\text{Comp}}] + \begin{array}{c} \square \square \\ \circ_1 \\ \square \square \end{array} \circ [\square, f_{\text{Comp}}].$$

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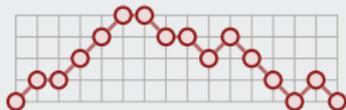
$$\mathcal{G}_{\text{Comp}}(t) = t + 2t\mathcal{G}_{\text{Comp}}(t) = t + 2t^2 + 4t^3 + 8t^4 + 16t^5 + \dots$$

Operad of Motzkin paths

The operad of Motzkin paths is the operad **Motz** such that

- ▶ **Motz**(n) is the set of the Motzkin paths with $n - 1$ steps.

Example



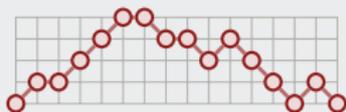
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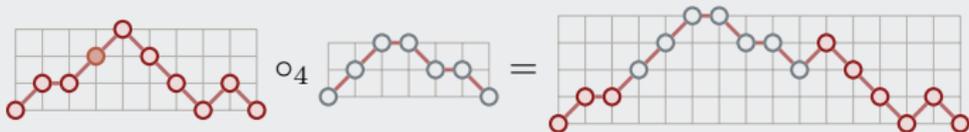
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Exemple

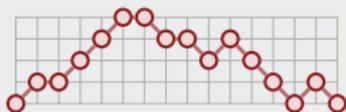


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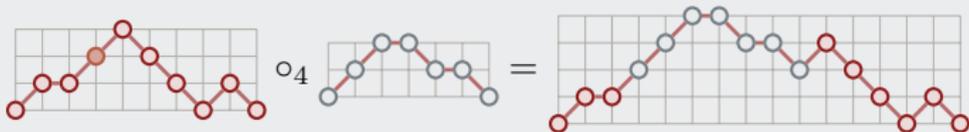
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Presentation of Motz

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$$\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \circ_1 \circ\circ \equiv \circ\circ \circ_2 \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array},$$

$$\circ\circ \circ_1 \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \equiv \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \circ_3 \circ\circ,$$

$$\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \circ_1 \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \equiv \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \circ_3 \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array}.$$

This presentation is not binary but is quadratic.

PBW basis of Motz and enumeration

Proposition

The set of the syntax trees on \mathfrak{S} avoiding



forms a PBW basis of **Motz**.

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This leads to

$$f_{\text{Motz}} = \circ + \circ\circ \circ [\circ, f_{\text{Motz}}] + \circ\circ \circ [\circ, f_{\text{Motz}}, f_{\text{Motz}}].$$

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Hence,

$$\mathcal{G}_{\text{Motz}}(t) = t + t\mathcal{G}_{\text{Motz}}(t) + t\mathcal{G}_{\text{Motz}}(t)^2.$$

Statistics on Motz

Let us add some parameters in the expression of f_{Motz} to obtain

$$g_{\text{Motz}} = \circ + q_0 \circ \circ \circ [\circ, g_{\text{Motz}}] + q_1 \circ \circ \circ [\circ, g_{\text{Motz}}, g_{\text{Motz}}].$$

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Let us add some parameters in the expression of f_{Motz} to obtain

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We obtain

$$\begin{aligned} \mathcal{G}_{\text{Motz}}(t, q_0, q_1) := \text{ev}(g_{\text{Motz}}) &= t + q_0 t^2 + (q_0^2 + q_1) t^3 \\ &\quad + (q_0^3 + 3q_0 q_1) t^4 \\ &\quad + (q_0^4 + 6q_0^2 q_1 + 2q_1^2) t^5 \\ &\quad + (q_0^5 + 10q_0^3 q_1 + 10q_0 q_1^2) t^6 \\ &\quad + (q_0^6 + 15q_0^4 q_1 + 30q_0^2 q_1^2 + 5q_1^3) t^7 \\ &\quad + (q_0^7 + 21q_0^5 q_1 + 70q_0^3 q_1^2 + 35q_0 q_1^3) t^8 \\ &\quad + (q_0^8 + 28q_0^6 q_1 + 140q_0^4 q_1^2 + 140q_0^2 q_1^3 + 14q_1^4) t^9 + \dots \end{aligned}$$

Statistics on Motz

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This counts Motzkin paths following the number of horizontal steps (parameter q_0) and rising steps (parameter q_1).

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This counts Motzkin paths following the number of horizontal steps (parameter q_0) and rising steps (parameter q_1).

For instance, $\mathcal{G}_{\text{Motz}}(t, 1, q_1)$ is Triangle **A055151**.

Operad of Schröder trees

The operad of Schröder trees is the operad Schr such that

- ▶ $\text{Schr}(n)$ is the set of the Schröder trees with $n + 1$ leaves.

Example



is a Schröder tree of arity 9.

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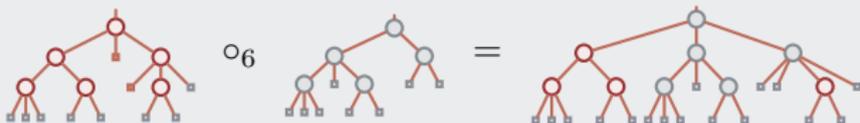
Example



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- ▶ The partial composition $\mathfrak{t} \circ_i \mathfrak{s}$ is obtained by inserting a copy of \mathfrak{s} between the i th and the $i + 1$ st leaves of \mathfrak{t} .

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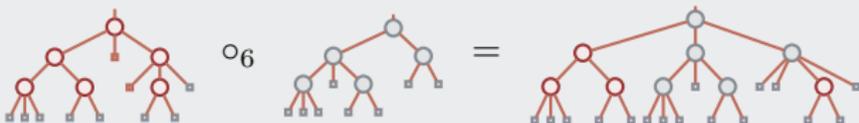
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Example



- ▶ The unit is .

Presentation of Schr

Proposition

The operad **Schr** admits the presentation (\mathfrak{G}, \equiv) where

$$\mathfrak{G} := \left\{ \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \right\}$$

and \equiv is the smallest operad congruence satisfying

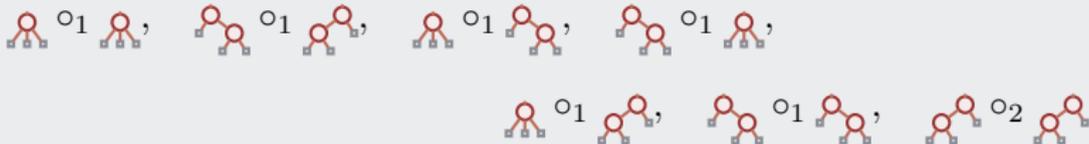
$$\begin{array}{ll} \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, & \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, \\ \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, & \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, \\ \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}, & \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \equiv \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array} \circ_2 \begin{array}{c} \circ \\ \swarrow \quad \searrow \\ \square \quad \square \end{array}. \end{array}$$

This presentation is binary and quadratic.

PBW basis of Schr and enumeration

Proposition

The set of the syntax trees on \mathfrak{S} avoiding

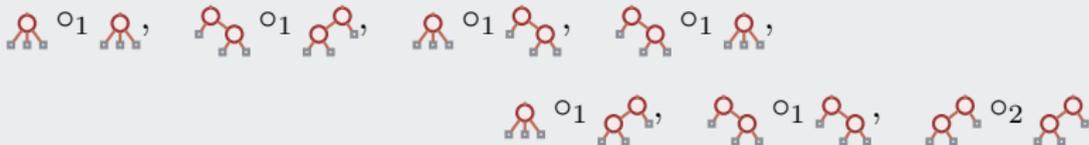


forms a PBW basis of **Schr**.

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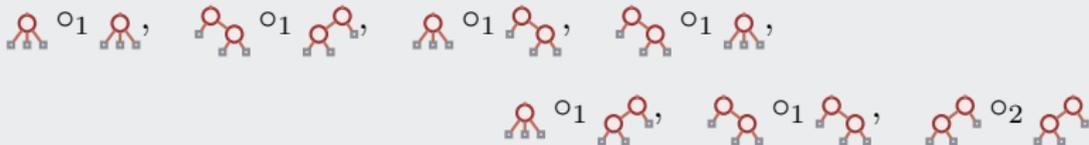
This leads to

$$\begin{aligned}
 f_{\text{Schr}} = & \circlearrowleft + \circlearrowright \circ \left[\circlearrowleft, f_{\text{Schr}} \right] + \circlearrowright \circ \left[\circlearrowright, f_{\text{Schr}} \right] \\
 & + \circlearrowright \circ \left[f_{\text{Schr}}, \circlearrowleft + \circlearrowright \circ \left[\circlearrowleft, f_{\text{Schr}} \right] + \circlearrowright \circ \left[\circlearrowright, f_{\text{Schr}} \right] \right].
 \end{aligned}$$

PBW basis of Schr and enumeration

Proposition

The set of the syntax trees on \mathfrak{S} avoiding



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This leads to

$$\begin{aligned}
 f_{\text{Schr}} = & \text{root node} + \text{root node} \circ \left[\text{root node}, f_{\text{Schr}} \right] + \text{root node} \circ \left[\text{root node}, f_{\text{Schr}} \right] \\
 & + \text{root node} \circ \left[f_{\text{Schr}}, \text{root node} + \text{root node} \circ \left[\text{root node}, f_{\text{Schr}} \right] + \text{root node} \circ \left[\text{root node}, f_{\text{Schr}} \right] \right].
 \end{aligned}$$

Hence,

$$\mathcal{G}_{\text{Schr}}(t) = t + 3t\mathcal{G}_{\text{Schr}}(t) + 2t\mathcal{G}_{\text{Schr}}(t)^2.$$

Statistics on Schr

Let us add some parameters in the expression of f_{Schr} to obtain

$$\begin{aligned} \mathfrak{g}_{\text{Schr}} = & \text{[tree]} + q_0 \text{[tree]} \circ \left[\text{[tree]}, \mathfrak{g}_{\text{Schr}} \right] + q_1 \text{[tree]} \circ \left[\text{[tree]}, \mathfrak{g}_{\text{Schr}} \right] \\ & + q_2 \text{[tree]} \circ \left[\mathfrak{g}_{\text{Schr}}, \text{[tree]} + \text{[tree]} \circ \left[\text{[tree]}, \mathfrak{g}_{\text{Schr}} \right] + \text{[tree]} \circ \left[\text{[tree]}, \mathfrak{g}_{\text{Schr}} \right] \right]. \end{aligned}$$

Statistics on Schr

Let us add some parameters in the expression of f_{Schr} to obtain

$$\begin{aligned} \mathbf{g}_{\text{Schr}} = & \text{Diagram}_1 + q_0 \text{Diagram}_2 \circ [\text{Diagram}_1, \mathbf{g}_{\text{Schr}}] + q_1 \text{Diagram}_3 \circ [\text{Diagram}_1, \mathbf{g}_{\text{Schr}}] \\ & + q_2 \text{Diagram}_4 \circ [\mathbf{g}_{\text{Schr}}, \text{Diagram}_1 + \text{Diagram}_2 \circ [\text{Diagram}_1, \mathbf{g}_{\text{Schr}}] + \text{Diagram}_3 \circ [\text{Diagram}_1, \mathbf{g}_{\text{Schr}}]] . \end{aligned}$$

We obtain

$$\begin{aligned} \mathcal{G}_{\text{Schr}}(t, q_0, q_1, q_2) := \text{ev}(\mathbf{g}_{\text{Schr}}) = & t + (q_0 + q_1 + q_2)t^2 \\ & + (q_0^2 + 2q_0q_1 + q_1^2 + 2q_0q_2 + 2q_1q_2 + q_2^2 + 2q_2)t^3 \\ & + (q_0^3 + 3q_0^2q_1 + 3q_0q_1^2 + q_1^3 + 3q_0^2q_2 + 6q_0q_1q_2 \\ & + 3q_1^2q_2 + 3q_0q_2^2 + 3q_1q_2^2 + q_2^3 + 6q_0q_2 + 6q_1q_2 + 6q_2^2)t^4 + \dots . \end{aligned}$$

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We obtain

$$\begin{aligned} \mathcal{G}_{\text{Schr}}(t, q_0, q_1, q_2) := \text{ev}(\mathbf{g}_{\text{Schr}}) = & t + (q_0 + q_1 + q_2)t^2 \\ & + (q_0^2 + 2q_0q_1 + q_1^2 + 2q_0q_2 + 2q_1q_2 + q_2^2 + 2q_2)t^3 \\ & + (q_0^3 + 3q_0^2q_1 + 3q_0q_1^2 + q_1^3 + 3q_0^2q_2 + 6q_0q_1q_2 \\ & + 3q_1^2q_2 + 3q_0q_2^2 + 3q_1q_2^2 + q_2^3 + 6q_0q_2 + 6q_1q_2 + 6q_2^2)t^4 + \dots \end{aligned}$$

The specialization $\mathcal{G}_{\text{Schr}}(t, 1, 1, q_2)$ is Triangle **A114656** and the specialization $\mathcal{G}_{\text{Schr}}(t, q_0, 1, 1)$ is Triangle **A098473**.

Operad of k -trees

The operad of k -trees is the operad $\mathbf{FCat}^{(k)}$, $k \in \mathbb{N}$, such that

- ▶ $\mathbf{FCat}^{(k)}(n)$ is the set of the planar rooted trees with n internal nodes, all having $k + 1$ children.

Example



is a 2-tree of arity 4.

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Example



- ▶ The unit is .

Presentation of $\mathbf{FCat}^{(k)}$

Proposition

The operad $\mathbf{FCat}^{(k)}$ admits the presentation (\mathfrak{G}, \equiv) where

$$\mathfrak{G} := \left\{ \begin{array}{c} \text{Diagram 1} \\ \text{Diagram 2} \end{array} : 0 \leq a \leq k+1 \right\}$$

and \equiv is the smallest operad congruence satisfying

$$\begin{array}{c} \text{Diagram 1} \\ \text{Diagram 2} \end{array} \circ_1 \begin{array}{c} \text{Diagram 3} \\ \text{Diagram 4} \end{array} \equiv \begin{array}{c} \text{Diagram 5} \\ \text{Diagram 6} \end{array} \circ_2 \begin{array}{c} \text{Diagram 7} \\ \text{Diagram 8} \end{array}, \quad 0 \leq a, b, \quad a + b \leq k.$$

This presentation is binary and quadratic.

PBW basis of $\mathcal{FCat}^{(k)}$ and enumeration

Proposition

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This leads to

$$\mathbf{f}_{\mathbf{FCat}^{(k)}} = \text{tree} + \sum_{0 \leq a \leq k} \mathbf{f}_{\mathbf{FCat}^{(k)}}^{(a)},$$

where the $\mathbf{f}_{\mathbf{FCat}^{(k)}}^{(a)}$ are the $\mathbf{FCat}^{(k)}$ -series satisfying

$$\mathbf{f}_{\mathbf{FCat}^{(k)}}^{(a)} = \text{tree} \circ \left[\text{tree} + \sum_{a+1 \leq b \leq k} \mathbf{f}_{\mathbf{FCat}^{(k)}}^{(b)}, \mathbf{f}_{\mathbf{FCat}^{(k)}} \right].$$

Enumeration in $\mathbf{FCat}^{(k)}$

By setting

$$\mathcal{G}_{\mathbf{FCat}^{(k)}}(t) := \text{ev}(\mathbf{f}_{\mathbf{FCat}^{(k)}})$$

and

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and

$$\mathcal{G}_{\mathbf{FCat}^{(k)}}^{(a)}(t) = t \mathcal{G}_{\mathbf{FCat}^{(k)}}(t) + \mathcal{G}_{\mathbf{FCat}^{(k)}}(t) \left(\sum_{a+1 \leq b \leq k} \mathcal{G}_{\mathbf{FCat}^{(k)}}^{(b)}(t) \right).$$

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This leads to

$$\mathcal{G}_{\mathbf{FCat}^{(k)}}^{(a)}(t) = t \mathcal{G}_{\mathbf{FCat}^{(k)}}(t) (1 + \mathcal{G}_{\mathbf{FCat}^{(k)}}(t))^{k-a}$$

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and

$$\mathcal{G}_{\mathbf{FCat}^{(k)}}(t) = t + t \mathcal{G}_{\mathbf{FCat}^{(k)}}(t) \sum_{0 \leq a \leq k} (1 + \mathcal{G}_{\mathbf{FCat}^{(k)}}(t))^{k-a}.$$

Statistics on $\text{FCat}^{(k)}$

Let us add some parameters in the expressions of $f_{\text{FCat}^{(k)}}$ to obtain

$$g_{\text{FCat}^{(k)}} = \text{Diagram} + \sum_{0 \leq a \leq k} g_{\text{FCat}^{(k)}}^{(a)},$$

where the $g_{\text{FCat}^{(k)}}^{(a)}$ are the $\text{FCat}^{(k)}$ -series satisfying

$$g_{\text{FCat}^{(k)}}^{(a)} = q_a \text{Diagram} \circ \left[\text{Diagram} + \sum_{a+1 \leq b \leq k} g_{\text{FCat}^{(k)}}^{(b)}, g_{\text{FCat}^{(k)}} \right].$$

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$$\mathbf{g}_{\text{FCat}^{(k)}}^{(a)} = q_a \text{Diagram} \circ \left[\text{Diagram} + \sum_{a+1 \leq b \leq k} \mathbf{g}_{\text{FCat}^{(k)}}^{(b)}, \mathbf{g}_{\text{FCat}^{(k)}} \right].$$

We obtain, when $k := 1$,

$$\begin{aligned} \mathcal{G}(t, q_0, q_1) := \text{ev}(\mathbf{g}_{\text{FCat}^{(1)}}) &= t + (q_0 + q_1)t^2 + (q_0^2 + 3q_0q_1 + q_1^2)t^3 \\ &\quad + (q_0^3 + 6q_0^2q_1 + 6q_0q_1^2 + q_1^3)t^4 \\ &\quad + (q_0^4 + 10q_0^3q_1 + 20q_0^2q_1^2 + 10q_0q_1^3 + q_1^4)t^5 \\ &\quad + (q_0^5 + 15q_0^4q_1 + 50q_0^3q_1^2 + 50q_0^2q_1^3 + 15q_0q_1^4 + q_1^5)t^6 \\ &\quad + (q_0^6 + 21q_0^5q_1 + 105q_0^4q_1^2 + 175q_0^3q_1^3 + 105q_0^2q_1^4 + 21q_0q_1^5 + q_1^6)t^7 + \dots \end{aligned}$$

Statistics on $\text{FCat}^{(k)}$

Let us add some parameters in the expressions of $f_{\text{FCat}^{(k)}}$ to obtain

$$g_{\text{FCat}^{(k)}} = \text{Diagram} + \sum_{0 \leq a \leq k} g_{\text{FCat}^{(k)}}^{(a)},$$

where the $g_{\text{FCat}^{(k)}}^{(a)}$ are the $\text{FCat}^{(k)}$ -series satisfying

$$g_{\text{FCat}^{(k)}}^{(a)} = q_a \text{Diagram} \circ \left[\text{Diagram} + \sum_{a+1 \leq b \leq k} g_{\text{FCat}^{(k)}}^{(b)}, g_{\text{FCat}^{(k)}} \right].$$

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This is Triangle **A001263**, known as the triangle of Narayana numbers.

General result

Theorem

Let \mathcal{O} be an operad admitting a presentation (\mathfrak{G}, \equiv) and \mathcal{B} be a PBW basis of \mathcal{O} for (\mathfrak{G}, \equiv) . Let the \mathcal{O} -series \mathbf{f} satisfying

$$\mathbf{f} = \mathbb{1} + \sum_{x \in \mathfrak{G}} \mathbf{f}_x,$$

where for any syntax tree \mathbf{t} on \mathfrak{G} of arity n , $\mathbf{f}_{\mathbf{t}}$ satisfies

$$\mathbf{f}_{\mathbf{t}} = \mathbf{t} \circ \left[\mathbf{f} - \sum_{\mathbf{t} \circ_1 \mathfrak{s} \notin \mathcal{B}} \mathbf{f}_{\mathfrak{s}}, \dots, \mathbf{f} - \sum_{\mathbf{t} \circ_n \mathfrak{s} \notin \mathcal{B}} \mathbf{f}_{\mathfrak{s}} \right].$$

Then, \mathbf{f} is the characteristic series of \mathcal{O} .

Contents

Constructions of operads

From monoids to operads

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$$u \circ_i v := u_1 \dots u_{i-1} (u_i * v_1) \dots (u_i * v_m) u_{i+1} \dots u_n.$$

Example

In $\mathbb{T}(\mathbb{N}, +, 0)$,

$$2100213 \circ_5 3001 = 2100522313.$$

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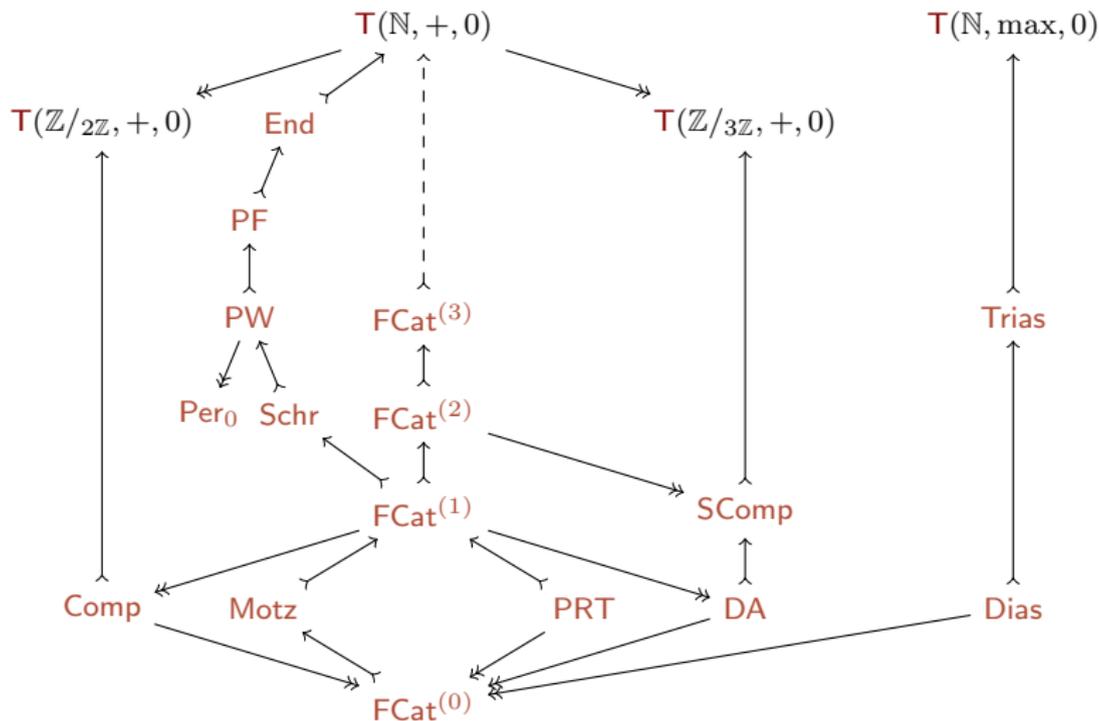
Theorem

For any monoid \mathcal{M} , $\mathbb{T}\mathcal{M}$ is an operad.

From monoids to operads – main examples

Monoid	Operad	Generators	First dimensions	Combinatorial objects
$(\mathbb{N}, +, 0)$	End	–	1, 4, 27, 256, 3125	Endofunctions
	PF	–	1, 3, 16, 125, 1296	Parking functions
	PW	–	1, 3, 13, 75, 541	Packed words
	Per_0	–	1, 2, 6, 24, 120	Permutations
	PRT	01	1, 1, 2, 5, 14, 42	Planar rooted trees
	$\text{FCat}^{(k)}$	$00, 01, \dots, 0k$	Fuß-Catalan numbers	k -trees
	Schr	$00, 01, 10$	1, 3, 11, 45, 197	Schröder trees
	Motz	$00, 010$	1, 1, 2, 4, 9, 21, 51	Motzkin words
$(\mathbb{Z}/2\mathbb{Z}, +, 0)$	Comp	$00, 01$	1, 2, 4, 8, 16, 32	Compositions
$(\mathbb{Z}/3\mathbb{Z}, +, 0)$	DA	$00, 01$	1, 2, 5, 13, 35, 96	Directed animals
	SComp	$00, 01, 02$	1, 3, 27, 81, 243	Seg. compositions
$(\mathbb{N}, \max, 0)$	Dias	$01, 10$	1, 2, 3, 4, 5	Bin. words with exact. one 0
	Trias	$00, 01, 10$	1, 3, 7, 15, 31	Bin. words with at least one 0

From monoids to operads – diagram



From posets to operads

Let (\mathcal{P}, \preceq) be a finite poset. We denote by \uparrow the binary (partial operation) min on \mathcal{P} with respect to the order relation \preceq .

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and \equiv is the smallest operad congruence satisfying

$$\star_a \circ_1 \star_b \equiv \star_{a \uparrow b} \circ_1 \star_{a \uparrow b} \equiv \star_{a \uparrow b} \circ_2 \star_{a \uparrow b} \equiv \star_a \circ_2 \star_b, \quad a \preceq b \text{ or } b \preceq a.$$

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Let (\mathcal{P}, \preceq) be a finite poset. We denote by \uparrow the binary (partial operation) min on \mathcal{P} with respect to the order relation \preceq .

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Example

Let the poset $\mathcal{P} := \begin{matrix} \textcircled{1} & \textcircled{2} \\ & \textcircled{3} \textcircled{4} \end{matrix}$.

Then, $\mathbf{P}\mathcal{P}$ is generated by $\{\star_1, \star_2, \star_3, \star_4\}$ and these generators are subjected to the relations

$$\star_1 \circ_1 \star_1 \equiv \star_1 \circ_1 \star_3 \equiv \star_3 \circ_1 \star_1 \equiv \star_3 \circ_2 \star_1 \equiv \star_1 \circ_2 \star_3 \equiv \star_1 \circ_2 \star_1,$$

$$\star_2 \circ_1 \star_2 \equiv \star_2 \circ_1 \star_3 \equiv \star_2 \circ_1 \star_4 \equiv \star_3 \circ_1 \star_2 \equiv \star_4 \circ_1 \star_2$$

$$\equiv \star_4 \circ_2 \star_2 \equiv \star_3 \circ_2 \star_2 \equiv \star_2 \circ_2 \star_4 \equiv \star_2 \circ_2 \star_3 \equiv \star_2 \circ_2 \star_2,$$

$$\star_3 \circ_1 \star_3 \equiv \star_3 \circ_2 \star_3, \quad \star_4 \circ_1 \star_4 \equiv \star_4 \circ_2 \star_4.$$

From posets to operads — PBW basis

Theorem

When \mathcal{P} is a finite poset avoiding the pattern



the operad $\mathbf{P}\mathcal{P}$ is a Koszul operad and the set of all syntax trees on \mathfrak{G} avoiding

$$\star_a \circ_1 \star_b, \quad a \preceq b \text{ or } b \preceq a$$

and

$$\star_a \circ_2 \star_b, \quad a \neq b \text{ and } (a \preceq b \text{ or } b \preceq a).$$

is a PBW basis of $\mathbf{P}\mathcal{P}$.

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is a PBW basis of $\mathbf{P}\mathcal{P}$.

When the premises of this theorem hold, $\mathbf{P}\mathcal{P}$ can be realized in terms of an operad of Schröder trees labeled on \mathcal{P} .

From posets to operads — example

More precisely, the elements of $\mathbf{PP}(n)$ are Schröder trees with n leaves such that the internal nodes are labeled on \mathcal{P} and the label of a node and any of its child are incomparable in \mathcal{P} .

From posets to operads — example

More precisely, the elements of $\mathbf{PP}(n)$ are Schröder trees with n leaves such that the internal nodes are labeled on \mathcal{P} and the label of a node and any of its child are incomparable in \mathcal{P} .

Example

Let

$$\mathcal{P} := \begin{array}{c} \textcircled{1} \\ \textcircled{2} \quad \textcircled{3} \\ \textcircled{4} \\ \textcircled{5} \\ \textcircled{6} \end{array} .$$

Then, in \mathbf{PP} ,

$$\begin{array}{c} \textcircled{6} \\ \textcircled{4} \quad \square \\ \square \quad \square \end{array} \circ_3 \begin{array}{c} \textcircled{2} \\ \textcircled{6} \\ \square \quad \square \end{array} = \begin{array}{c} \textcircled{6} \\ \textcircled{4} \quad \textcircled{2} \\ \square \quad \square \quad \textcircled{6} \\ \square \quad \square \end{array} ,$$

$$\begin{array}{c} \textcircled{1} \\ \textcircled{4} \\ \square \quad \square \end{array} \circ_1 \begin{array}{c} \textcircled{2} \\ \textcircled{3} \quad \textcircled{3} \\ \square \quad \square \quad \square \quad \square \end{array} = \begin{array}{c} \textcircled{1} \\ \square \quad \square \quad \square \quad \textcircled{4} \\ \square \quad \square \end{array} .$$

From unitary magmas to operads — cliques

Let $(\mathcal{M}, *, \mathbb{1}_{\mathcal{M}})$ be a unitary magma.

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Let $(\mathcal{M}, *, \mathbb{1}_{\mathcal{M}})$ be a unitary magma.

An \mathcal{M} -clique \mathfrak{p} is a clique on $[n + 1]$ where each edge (x, y) is labeled by an element $\mathfrak{p}(x, y)$ of \mathcal{M} . The arity $|\mathfrak{p}|$ of \mathfrak{p} is n .

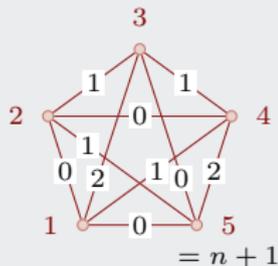
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Example

Let \mathcal{M} be the unitary magma $(\mathbb{Z}/3\mathbb{Z}, +, 0)$. Here is an \mathcal{M} -clique:



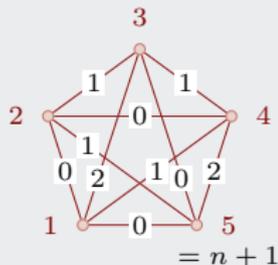
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Example

Let \mathcal{M} be the unitary magma $(\mathbb{Z}/3\mathbb{Z}, +, 0)$. Here is an \mathcal{M} -clique:



An edge (x, y) of \mathbf{p} is **solid** if $\mathbf{p}(x, y) \neq \mathbb{1}_{\mathcal{M}}$.

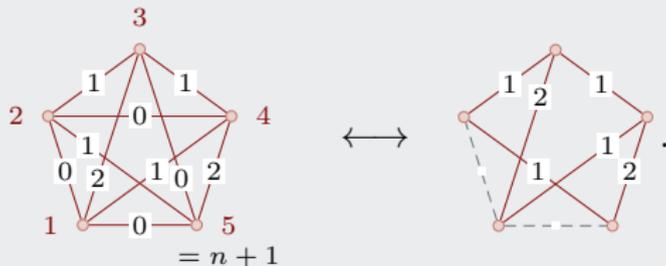
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Let \mathcal{M} be the unitary magma $(\mathbb{Z}/3\mathbb{Z}, +, 0)$. Here is an \mathcal{M} -clique:



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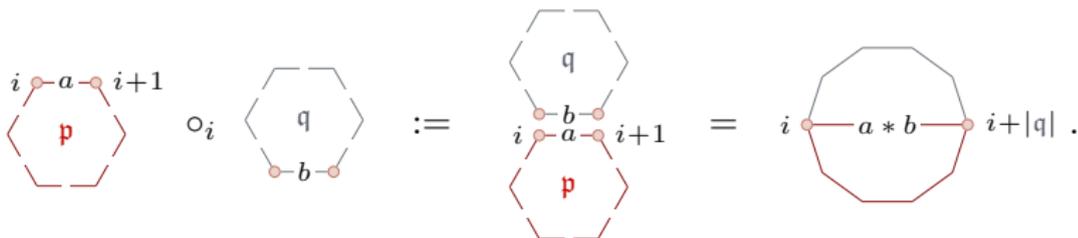
- ▶ $\mathbf{C}\mathcal{M}(n)$ is the set of all \mathcal{M} -cliques of arity n . By convention, $\mathbf{C}\mathcal{M}(1)$ is the singleton containing $\circ - \circ$.

From unitary magmas to operads

Let $(\mathcal{M}, *, \mathbb{1}_{\mathcal{M}})$ be a unitary magma.

We define $(\mathbf{CM}, \circ_i, \mathbb{1})$ as the triple such that

- ▶ $\mathbf{CM}(n)$ is the set of all \mathcal{M} -cliques of arity n . By convention, $\mathbf{CM}(1)$ is the singleton containing $\circ - \circ$.
- ▶ For any \mathcal{M} -cliques \mathfrak{p} and \mathfrak{q} ,

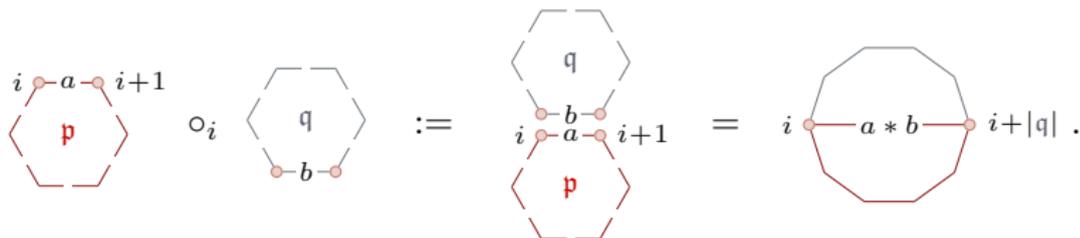


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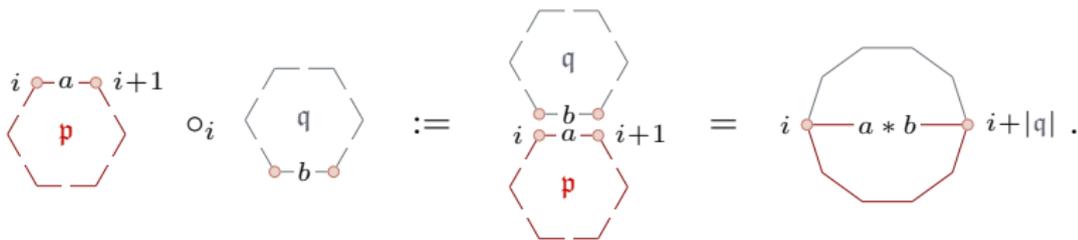
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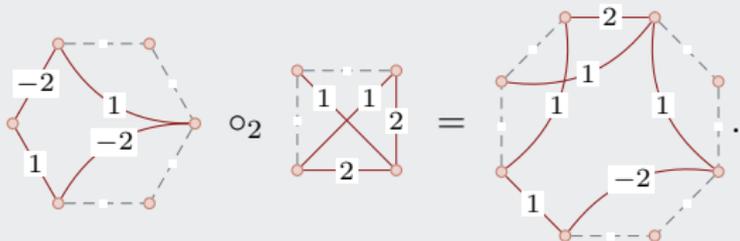
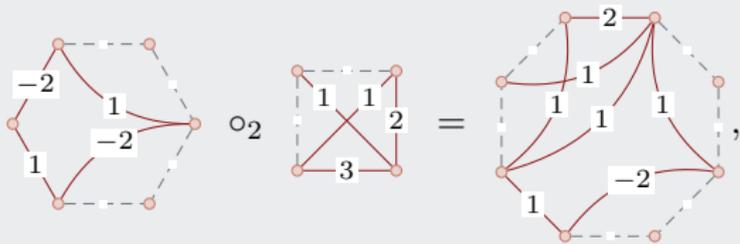
Theorem

For any unitary magma \mathcal{M} , \mathbf{CM} is an operad.

From unitary magmas to operads — example

Exemple

In $\mathbf{C}(\mathbb{Z}, +, 0)$,



From unitary magmas to operads — main examples

The construction \mathbf{C} leads to alternative definitions of

- ▶ the operad \mathbf{NCT} of noncrossing trees [Chapoton, 2007];
- ▶ the operad \mathbf{FF}_4 of some formal fractions [Chapoton, Hivert, Novelli, 2016];
- ▶ the operad \mathbf{BNC} of bicolored noncrossing configurations [Chapoton, Giraudo, 2014];
- ▶ the operad \mathbf{MT} of multi-tildes [Luque, Mignot, Nicart, 2013];
- ▶ the operad \mathbf{DMT} of double multi-tildes [Giraudo, Luque, Mignot, Nicart, 2016];
- ▶ the dipterous operad [Loday, Ronco, 2003];
- ▶ the gravity operad [Getzler, 1994].