

Algebraic structures, series, and enumeration

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Outline

Combinatorics and enumeration

Formal power series

Operads and composition

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Combinatorial sets

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such that for any $n \in \mathbb{N}$, $C(n) := \{x \in C : |x| = n\}$ is finite.

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5. **Define statistics** on C , that are maps $s : C \rightarrow \mathbb{N}$.

Examples

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A **descent** of $\sigma \in \mathfrak{S}$ is a position i such that $\sigma(i) > \sigma(i+1)$.

The map s sending a permutation to its number of descents is a statistics on \mathfrak{S} . For instance, $s(3**5**1**4**2) = 2$.

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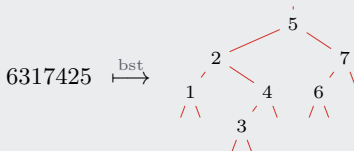
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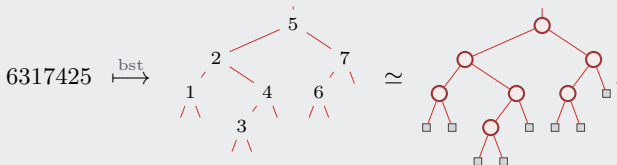
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Generating series

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$$\mathcal{G}_C(t) := \sum_{n \in \mathbb{N}} \#C(n) t^n = \sum_{x \in C} t^{|x|}.$$

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► $\mathcal{G}_{\text{BT}}(t) = 1 + t + 2t^2 + 5t^3 + 14t^4 + 42t^5 + \dots = \frac{1 - \sqrt{1 - 4t}}{2t}$

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$$\xrightarrow{\text{ev}}$$

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3. operations on C
 \leadsto extensions to **operations on formal power series**.

Outline

Formal power series

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The sum notation of \mathbf{f} is

$$\mathbf{f} = \sum_{x \in X} \langle x, \mathbf{f} \rangle x.$$

Formal series and algebraic structures

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Example

A binary product $\star : X \times X \rightarrow X$ leads to the (possibly partial) product

$$\mathbf{f} \bar{\star} \mathbf{g} := \sum_{x, y \in X} \langle x, \mathbf{f} \rangle \langle y, \mathbf{g} \rangle x \star y$$

on $\mathbb{K} \langle\langle X \rangle\rangle$.

Formal series and algebraic structures

Here are some examples of series on algebraic structures leading to well-known objects.

Structure on X	Sort of series
$(\mathbb{N}, +, 0)$	Usual series $\mathbb{K}[[t]]$
Free comm. monoid	Multivariate series $\mathbb{K}[[t_1, t_2, \dots]]$
Free monoid	Noncomm. series $\mathbb{K}\langle\langle t_1, t_2, \dots \rangle\rangle$ [Eilenberg, 1974]
Monoid	Series on monoids [Salomaa, Soittola, 1978]
Operads	Series on operads [Chapoton, 2002, 2008]

Generalizing the product of generating series

Let C be a combinatorial set.

A binary product \star on C is **graded** if $|x_1 \star x_2| = |x_1| + |x_2|$.

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The map ev is an algebra morphism between $(\mathbb{K} \langle\langle C \rangle\rangle, \bar{\star})$ and $(\mathbb{K} \langle\langle t \rangle\rangle, \cdot)$.

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Proof.

$$\begin{aligned} \text{ev}(\mathbf{f}_1 \bar{\star} \mathbf{f}_2) &= \sum_{x_1, x_2 \in C} \langle x_1, \mathbf{f}_1 \rangle \langle x_2, \mathbf{f}_2 \rangle \text{ev}(x_1 \star x_2) \\ &= \sum_{x_1, x_2 \in C} \langle x_1, \mathbf{f}_1 \rangle \langle x_2, \mathbf{f}_2 \rangle t^{|x_1| + |x_2|} \\ &= \left(\sum_{x_1 \in C} \langle x_1, \mathbf{f}_1 \rangle t^{|x_1|} \right) \left(\sum_{x_2 \in C} \langle x_2, \mathbf{f}_2 \rangle t^{|x_2|} \right) \\ &= \text{ev}(\mathbf{f}_1) \cdot \text{ev}(\mathbf{f}_2) \end{aligned}$$



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Therefore, $\bar{\star}$ offers a generalization of the usual product of generating series.

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Let \mathcal{P} be the (non combinatorial) set of all paths where the size of a path $u_1 \dots u_n$ is $n - 1$ (number of steps).

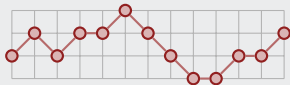
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Example

The path 1212232100112 is depicted as



and its size is 12.

Example: series of paths — a monoid of paths

Let \star be the binary product on \mathcal{P} defined by

$$u_1 \dots u_n \star v_1 \dots v_m :=$$

$$\uparrow_{\max(0, v_1 - u_n)} (u_1 \dots u_{n-1}) \max(u_n, v_1) \uparrow_{\max(0, u_n - v_1)} (v_2 \dots v_m)$$

where $\uparrow_i(w)$ is the word obtained by incrementing by i each letter of w .

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This product is associative, admits \bullet (the path 0) as unit, and is graded.

Hence, $(\mathcal{P}, \star, \bullet)$ is a graded monoid.

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The \star -Kleene star of a subset X of \mathcal{P} is

$$X^{\star\star} := \bigsqcup_{\ell \geq 0} \underbrace{X \star \dots \star X}_{\times \ell}.$$

Example: series of paths — Schröder paths

Let Schr be the subset of $\left\{ \begin{array}{c} \circ \\ \diagup \diagdown \\ \circ \end{array}, \begin{array}{c} \circ - \circ - \circ \\ \diagup \diagdown \end{array}, \begin{array}{c} \circ \\ \diagup \diagdown \\ \circ \end{array} \right\}^{**}$ restrained to the paths starting and finishing by 0. We call these **Schröder paths**.

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and, by using the product $\bar{\star}$ on \mathcal{P} -series, we obtain the relation

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By evaluation, we get the expression

$$\mathcal{G}_{\text{Schr}}(t) = 1 + t^2 \mathcal{G}_{\text{Schr}}(t) + t^2 \mathcal{G}_{\text{Schr}}(t)^2$$

for the generating series of the Schröder paths.

Example: series of paths — statistics

We can add formal parameters to refine the enumeration by taking into account of **statistics**. Then,

$$g_{\text{Schr}} = \circ + q_0 \circ \circ \circ \star g_{\text{Schr}} + q_1 \begin{array}{c} \circ \\ \diagup \diagdown \\ \circ \end{array} \star g_{\text{Schr}} \star \begin{array}{c} \circ \\ \diagdown \diagup \\ \circ \end{array} \star g_{\text{Schr}}$$

is the formal series of Schröder paths where the coefficient of a path u is $q_0^{\alpha_0} q_1^{\alpha_1}$ where α_0 (resp. α_1) is half of the number of $\circ \circ$ (resp. the number of $\begin{array}{c} \circ \\ \diagup \diagdown \\ \circ \end{array}$) appearing in it.

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Its evaluation is

$$\begin{aligned} \text{ev}(g_{\text{Schr}}) &= 1 + (q_0 + q_1) t^2 + (q_0^2 + 3q_0 q_1 + 2q_1^2) t^4 \\ &\quad + (q_0^3 + 6q_0^2 q_1 + 10q_0 q_1^2 + 5q_1^3) t^6 \\ &\quad + (q_0^4 + 10q_0^3 q_1 + 30q_0^2 q_1^2 + 35q_0 q_1^3 + 14q_1^4) t^8 + \dots \end{aligned}$$

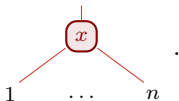
Its specialization $\text{ev}(g_{\text{Schr}})|_{q_1:=1}$ is Triangle **A088617** of OEIS and $\text{ev}(g_{\text{Schr}})|_{q_0:=1}$ is Triangle **A060693** of OEIS.

Outline

Operads and composition

Operators

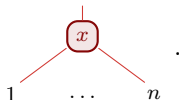
An **operator** is an entity having $n \geq 1$ inputs and a single output:



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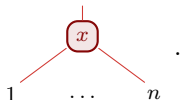
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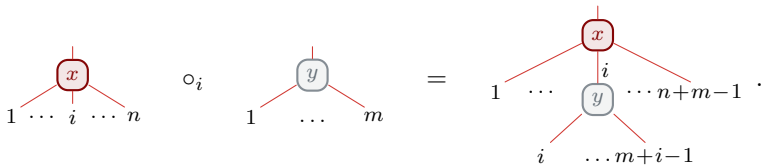


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Composing two operators x and y consists in

1. selecting an input of x specified by its position i ;
2. grafting the output of y onto this input.

This produces a new operator $x \circ_i y$ of arity $n + m - 1$:



Operads

Operads are algebraic structures formalizing the notion of operators and their composition.

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A (nonsymmetric set-theoretic) **operad** is a triple $(\mathcal{O}, \circ_i, \mathbb{1})$ where

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$$\mathcal{O} := \bigsqcup_{n \geq 1} \mathcal{O}(n);$$

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2. \circ_i is a map, called **partial composition map**,

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This data has to satisfy some axioms.

Operad axioms

Associativity:

$$(x \circ_i y) \circ_{i+j-1} z = x \circ_i (y \circ_j z)$$

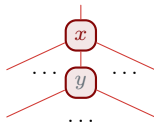
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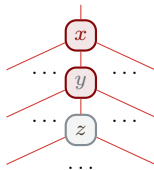


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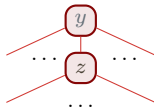
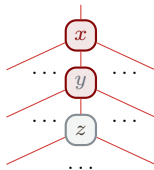


Operad axioms

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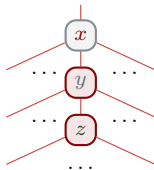
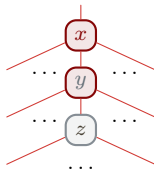


Operad axioms

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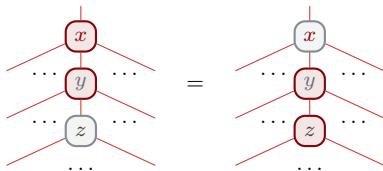


Operad axioms

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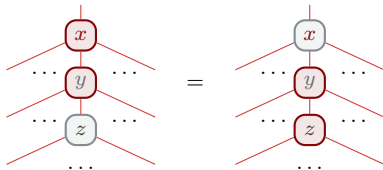
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Operad axioms

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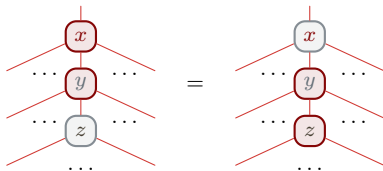
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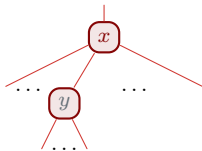
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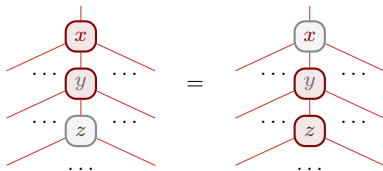


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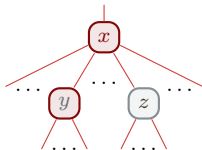
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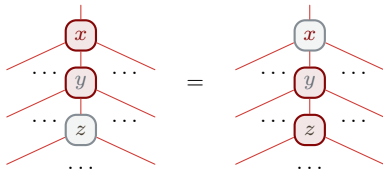


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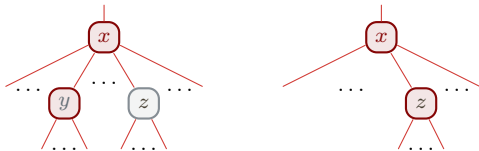
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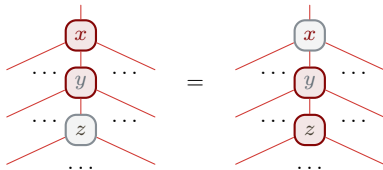


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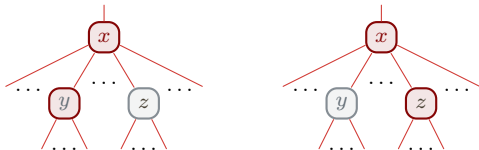
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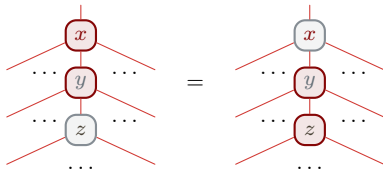


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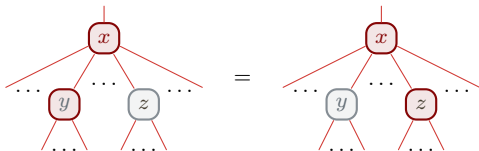
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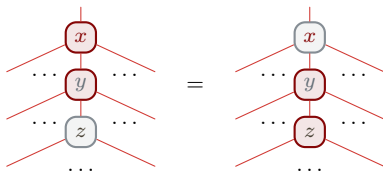


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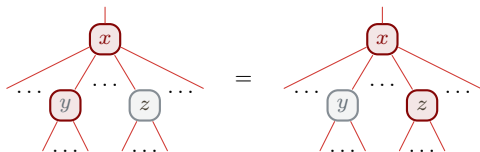
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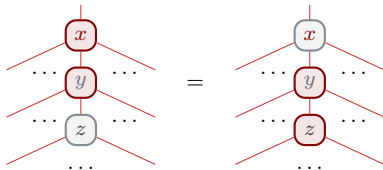
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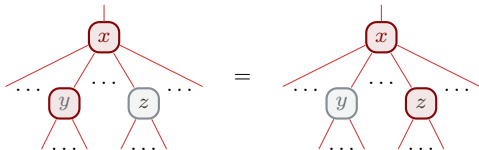
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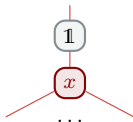
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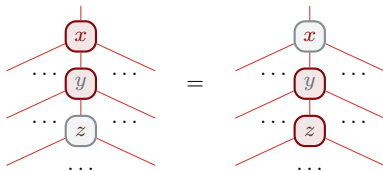


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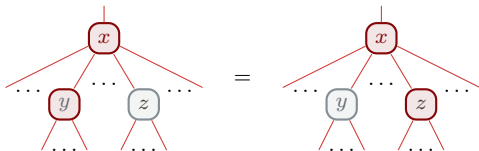
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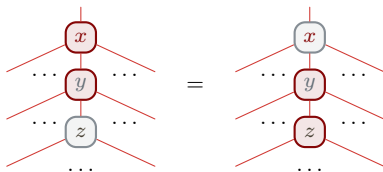


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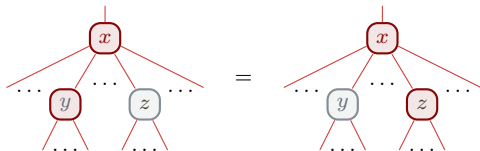
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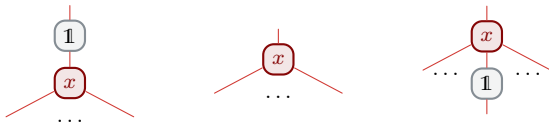
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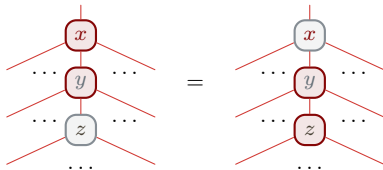


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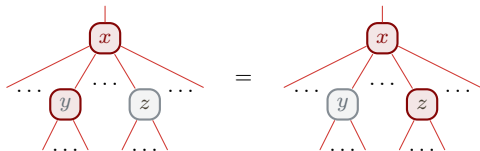
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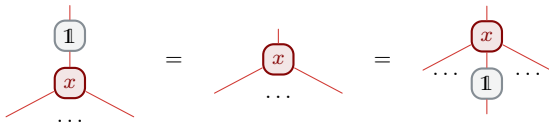
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Complete composition

Let \mathcal{O} be an operad

The **complete composition map** of \mathcal{O} is the map

$$\circ : \mathcal{O}(n) \times \mathcal{O}(m_1) \times \cdots \times \mathcal{O}(m_n) \rightarrow \mathcal{O}(m_1 + \cdots + m_n)$$

defined by

$$x \circ [y_1, \dots, y_n] := (\dots ((x \circ_n y_n) \circ_{n-1} y_{n-1}) \dots) \circ_1 y_1.$$

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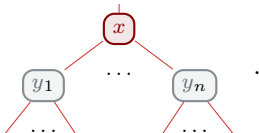
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Then, $x \circ [y_1, \dots, y_n]$ is the operator

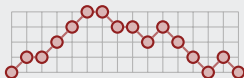


An operad of Motzkin paths

Let **Motz** be an operad wherein:

- **Motz**(n) is the set of all Motzkin paths with n points.

Example



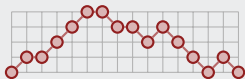
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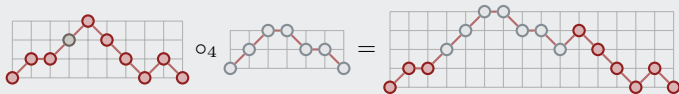
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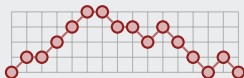


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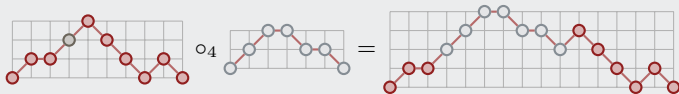
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on \mathcal{O} -series, satisfying for any $\mathbf{f}, \mathbf{g}_1, \dots, \mathbf{g}_n \in \mathbb{K} \langle \langle \mathcal{O} \rangle \rangle$ and $x \in \mathcal{O}$,

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Generalizing the composition of generating series

Proposition

The map ev is a monoid morphism between $(\mathbb{K} \langle\langle \mathcal{O} \rangle\rangle, \odot, \mathbf{1})$ and $(t\mathbb{K} \langle\langle t \rangle\rangle, \circ, t)$.

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$$\begin{aligned}\text{ev}(\mathbf{f} \odot \mathbf{g}) &= \sum_{\substack{y \in \mathcal{O} \\ z_1, \dots, z_{|y|} \in \mathcal{O}}} \langle y, \mathbf{f} \rangle \left(\prod_{1 \leq i \leq |y|} \langle z_i, \mathbf{g} \rangle \right) \text{ev}(y \circ [z_1, \dots, z_{|y|}]) \\ &= \sum_{\substack{y \in \mathcal{O} \\ z_1, \dots, z_{|y|} \in \mathcal{O}}} \langle y, \mathbf{f} \rangle \left(\prod_{1 \leq i \leq |y|} \langle z_i, \mathbf{g} \rangle \right) t^{|z_1| + \dots + |z_{|y|}|} \\ &= \left(\sum_{y \in \mathcal{O}} \langle y, \mathbf{f} \rangle t^{|y|} \right) \circ \left(\sum_{z \in \mathcal{O}} \langle z, \mathbf{g} \rangle t^{|z|} \right) \\ &= \text{ev}(\mathbf{f}) \circ \text{ev}(\mathbf{g})\end{aligned}$$



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Therefore, \odot offers a generalization of the usual composition of generating series.

Example — Motzkin paths

Proposition [G., 2015]

The set

$$\left\{ \text{---} \circ \text{---} \circ \text{---}, \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ \circ \quad \circ \end{array} \right\}$$

is the unique minimal generating set of the operad **Motz**.

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From this (and a reasoning about unambiguous decompositions of Motzkin paths), we obtain that the characteristic series of **Motz** satisfies

$$\mathbf{f}_{\text{Motz}} = \circ + \circ - \circ \bar{\circ} [\circ, \mathbf{f}_{\text{Motz}}] + \begin{array}{c} \circ - \circ \\ \circ - \circ - \circ \end{array} \bar{\circ} [\circ, \mathbf{f}_{\text{Motz}}, \mathbf{f}_{\text{Motz}}].$$

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Now, the evaluation $\mathcal{G}_{\text{Motz}}(t)$ of \mathbf{f}_{Motz} satisfies

$$\mathcal{G}_{\text{Motz}}(t) = t + t\mathcal{G}_{\text{Motz}}(t) + t\mathcal{G}_{\text{Motz}}(t)^2$$

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2. define operations \star on C ;
3. extends these operations to operations $\bar{\star}$ on C -series;
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We have in passing defined generalizations of natural operations on generating series:

Operation on C	Operation on $\mathbb{K} \langle\langle C \rangle\rangle$
—	Addition $+$
Binary graded product \star	Analog $\bar{\star}$ of the multiplication
Operadic product \circ	Analogs $\bar{\circ}$ and \odot of the composition