# Canonical Ordering for Triangulations on the Cylinder, with Applications to Periodic Straight-line Drawings 

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GD'12, Redmond

Planar straight-line drawings

[Wagner'36]
[Fary'48] [Stein'51]

Planar straight-line drawings

[Wagner'36]
[Fary'48]
[Stein'51]

Classical algorithms:

[Tutte'63]
spring-embedding

[De Fraysseix,Pach,Pollack'88]
"FPP algo" incremental

[Schnyder'90] face-counting principle

Periodic straight-line drawings
On the cylinder ( $\Leftrightarrow$ annulus)

lifts to $x$-periodic drawing

[Kocay et al'01]
[Gortler et al'06]
[Gonçalves-Lévêque'12]
lifts to
$x$-periodic \& $y$-periodic drawing


1. Recall FPP algorithm 2. Extend to the cylinder 3. Get toroidal drawings [Castelli,Devillers,F'12]

Plane


Grid $(2 n-4) \times(n-2)$

## Cylinder



Torus


1. Recall FPP algorithm 2. Extend to the cylinder 3. Get toroidal drawings [Castelli,Devillers,F'12]

Plane


Torus

$\operatorname{Grid} \leq 2 n \times(1+n(2 c+1))$

Canonical ordering for planar triangulations $T$ a planar triangulation with marked bottom-edge $e$ Canonical ordering $=$ shelling order (from top to bottom)

## $T \backslash e$ has 7 vertices



At each step:

$\Downarrow$


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## Canonical ordering: primal tree \& dual tree

 primal tree: parent of each vertex $v=$ neighbour of $v$ of largest label dual tree: dual of primal tree (augmented with two outer edges)

## Canonical ordering: successive induced graphs



Notation: $G_{k}$ is the graph formed by $e$ and $\{1, \ldots, k\}$

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Incremental drawing algorithm

## [de Fraysseix, Pollack, Pach'89]



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## [de Fraysseix, Pollack, Pach'89]

1. $\triangle$

2. 



Grid size of $G_{k}: 2 k \times k$


## Reformulation of the shift-step

At each step: insert two vertical strips of width 1 using the dual tree $G_{k-1}$


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## $G_{k}$



Extension to the cylinder: canonical ordering


## annular representation

At each step:


Extension to the cylinder: canonical ordering


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Extension to the cylinder: canonical ordering


## annular representation

At each step:


Terminates if no chordal edge incident to bottom cycle

Extension to the cylinder: canonical ordering


## annular representation

At each step:


Underlying forest


## annular representation



At each step:


Underlying forest + dual forest

Extension to the cylinder: drawing algorithm
$G_{k-1}$


At each step: - insert two vertical strips of width 1 - insert next vertex as in the planar case

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## Execution on an example



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Execution on an example


## Grid size: $w=2 n$

Each edge has vertical extension at most $w$

$$
\Rightarrow h \leq n(2 d+1)
$$

with $d$ the graph-distance between the two boundaries

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Dealing with chordal edges at outer cycle initial spacing


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Toroidal ṭriangulations \& tambourịne,, Every toroidal triangulation admits a "tambourine" [Bonichon, Gavoille, Labourel'06]

## Torus



Cylinder


## Getting a drawing on the (flat) torus


resinsert edges in tambourine


## Getting a drawing on the (flat) torus


$\Delta h \leq 2 n+1$
resinsert edges in tambourine


$$
\begin{aligned}
& w \leq 2 n \\
& h \leq n(2 d+1)
\end{aligned}
$$

$$
d=2
$$

## Getting a drawing on the (flat) torus



Torus

$\Delta h \leq 2 n+1$
resinsert edges in tambourine


Let $c=$ length shortest non-contractible cycle, $c \leq \sqrt{2 n}$
Can choose tambourine so that $d<c \Rightarrow h=O\left(n^{3 / 2}\right)$

## Extensions and perspectives

- Extension to 3-connected maps on cylinder \& torus (cf Kant's extension of the FPP algorithm in the planar case)

weakly convex periodic straight-line drawing


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- Extension to 3-connected maps on cylinder \& torus (cf Kant's extension of the FPP algorithm in the planar case)

weakly convex periodic straight-line drawing
- Extension of our method to higher genus ?

Polygonal scheme

[Duncan, Goodrich, Kobourov'99]
[Chambers, Eppstein, Goodrich, Löffler'00] drawing in polynomial area

periodic drawing out of circle packing

